

SECURITY SYSTEM OF A RELATIONAL DATA BASE SYSTEM

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By
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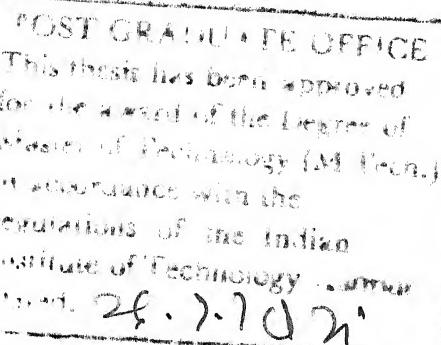
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CERTIFICATE

CERTIFIED that the work entitled SECURITY SYSTEM OF A RELATIONAL DATA BASE SYSTEM has been carried out under my supervision by Sri S. K. Goel and it has not been submitted elsewhere for a degree.

Kanpur
July 19, 1978

R. Sankar
Professor of Computer Sciences
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Credit for expert typing goes to Mr. H.K. Nathani, his patience is gratefully acknowledged.

Kanpur
July 19, 1978

- Satish K. Goel

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ABSTRACT

This thesis deals with the design and implementation of a security mechanism for a relational data base. It also deals with removing a relation from the data base and reorganization of the data base.

It is a part of the project to implement a relational data base system on TDC-316. This thesis alongwith (Ref. 5) completes the implementation of functions falling in the domain of the data base administrator, i.e., controlling the overall view of the data, reorganization of the data base and ensuring security of the data base.

1. INTRODUCTION

One of the most important aspects of a data base system is to provide for adequate security of the data base. Since a data base generally contains the lump sum of data useful for any organization and various users keep interacting with the same, retrieving and updating the data placed in a central place, it is highly important that there be a mechanism for stopping illegal use of the data. Hence appropriate checks must be placed on every user of the data base. Moreover, there must be ways of checking out users who tend to use the system in an unauthorised manner and also indicate the same whenever asked to do so, so that defaulting users may be asked for proper explanations.

The description of a relational data base system and the various terms related to it as well as hardware and software tools available are very well described in a separate thesis (Reference 5). A data base as we know, is the collection of data useful for many applications and placed in a central place. The various applications extract from the data base their own views of data and application program are independent of the way data is organised in a data base. It is the responsibility of the data base management system to take care of any changes in the physical organization of the data base, to make the application programs immune to any changes taking place in the physical data base organization.

In a relational data base system, all the data is stored in terms of relations, i.e., flat files. A flat file is many similar occurrences of a record which consists of a number of fields.

The security system for a data base can be of varied complexity. The following discussion seeks to illustrate this point: Suppose that our data base contains PERSON-DETAILS as one of the relations as shown in Figure 1-1. A few security implementations (of increasing complexity)

Personnel No. P#	Name N	Rank R	Pay P	Confidential Report CR

Figure 1-1.

could be as follows:

1. The user has access to all records of the relation if he is an authorised user, otherwise he has no access to the relation, i.e., either a user has authority to use a relation in entirety or not at all.
2. The user may see name, rank and pay of any person but may not see confidential report of any one. Thus the authority here is of field level i.e., either the user has some particular authority over a particular field all records or he has no authority over that field for all records.

3. The user may see the name and rank of any person but may see his pay and confidential report only if his rank is higher than that of the person under consideration. Moreover a user may change the name if the record pertains to his own personnel number. Here we see that the user has some authority over a field for a particular record depending upon his qualification in relation to the other field values for that record. Here the users identity must reveal to the system his personnel number, rank etc., to be able to enforce this type of security scheme.

Moreover there can be other access parameters like location of the accessor, the time and day of the access and the maximum frequency with which an accessor is allowed to access the system etc.

This present thesis describes a security scheme of the field level of a relation as in No. 2 above. The five access privileges associated with a field are as below :

1. READ : In this case, the particular field value can be read for all records for processing.
2. OUTPUT : In this case the particular field value can be printed on a listing device for any record.
3. DELETE : In this case a record of a relation may be deleted if the accessor has DELETE right for all its fields.
4. MODIFY : In this case, the particular field value may be modified for any record.
5. INSERT : In this case, any no. of records can be inserted in a relation if the user has INSERT right over all its fields.

The access in the present system does not depend upon any parameters like location, time and day, frequency etc., because of hardware limitations (because the system has only one access location and there is no timer in the system).

In the last chapter, programs are described for removing a relation from the data base and a discussion on how to reorganise the data base.

2. PLANNING OF THE SECURITY SYSTEM

The previous chapter describes how a security system can be of varied complexity. The most primitive security scheme could be to allow or disallow any user from accessing the data base. Once a user is put through, he shall be able to perform any operations on the data in the data base. But such a scheme is not suitable for any practical data base system because a data base essentially contains data for various applications and we shall like any user to have only limited access to it, to perform the functions falling only in his application area. As the complexity of the security mechanism increases, there will be a corresponding increase in its cost in terms of software development effort and more prominently in terms of the fraction of total run time taken by the security mechanism to respond to any query.

The present security mechanism is as described below:

We know that in a Relational Data Base System, each relation consists of some fields, each identified by a field identification number. A field may be common to many relations, but its field identification number is the same for its occurrence in every relation. In the type of security system provided here, we shall build an authority vector or access rights vector for any user from the security codes supplied by him to the system. The i th element of this vector shall give the rights of the user related to the field with identification No. i . Since TDC-316 is a byte addressable

machine with 8 bit byte, it was decided that each element of the authority vector shall be stored in a byte, whose 8 bits shall be associated with the following rights. Figure 2-1 depicts this association.

1. READ
2. OUTPUT
3. DELETE
4. MODIFY
5. INSERT
6. }
7. } UNUSED
8. }

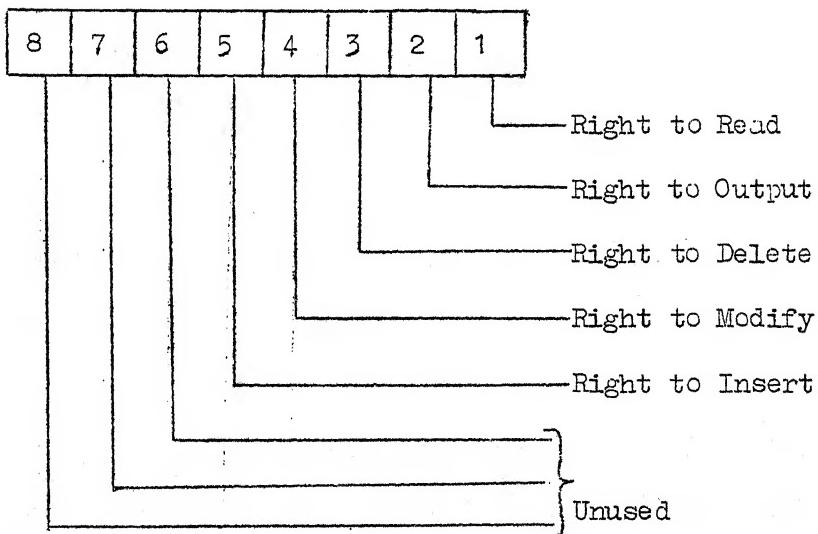


Figure 2-1: Rights associated with byte bits.

If a user has a particular right over a field, the corresponding bit shall be set, else it will be clear. Since the core memory available is very limited and all security codes and authority augmenting vectors cannot be profitably accommodated in core, it was decided to store the security codes and authority augmenting vectors in some area of the disk.

The present disk pack attached to TDC-316 has 203 cylinders. Each cylinder has 10 tracks (0 to 9) and each track is divided into 10 sectors (1 to 10). Each sector has a capacity to store 256 bytes (8 bits each).

Since the upper limit on the number of fields is 256 in the system designed, hence each authority augmenting vector can be accommodated in a sector of the disk.

In our system, we shall have two types of security codes:

1. FIRST SECURITY CODE: This is a code which helps the user to identify himself to the system. This helps the user to get access to the system. This is also called identification number of the user and is unique for each user.
2. AUTHORITY AUGMENTING SECURITY CODE: This code is used for augmenting the authority of the user. Each "Authority Augmenting Code" has an authority augmenting vector associated with it. A user may give more than one authority augmenting code to the system. The authority vector of a user is the logical OR of all those authority augmenting vectors for which he has given the corresponding authority augmenting security codes as input to the system.

Though authority vectors could be associated directly with "First Security Codes", i.e., user's identification numbers, but having two types of security codes as provided is useful on two counts -

1. Even if an unauthorised user breaks through the system by trying various first security codes till he succeeds in getting the right one (even though this is highly improbable), it is just impossible for him to give the correct authority augmenting security code in the first trial. Once he gives the wrong authority augmenting security code (after having succeeded in breaking the cardon of first security codes), the corresponding "First security code" is deleted from the group of valid first security codes and hence his access is again restricted to the system. It is in a way like a two level '': ': protection, where any mistake at the second level puts the user out of the first level.

2. In general, many users (at the same level of management) shall have similar authority vectors. Therefore, we can merge their authority vectors into one and give them same authority augmenting security code whereas they have different first security codes. Since an authority augmenting vector takes much larger storage space (one sector) compared to that taken for storing only security code, we shall save in storage space on disk by merging authority augmenting vectors.

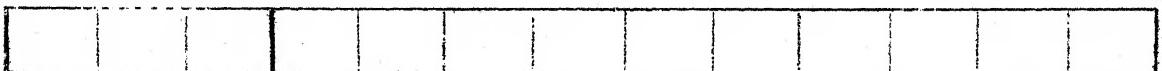
DATA STRUCTURES USED

1. FIRST SECURITY CODES

Since there are many first security codes (one for each user), therefore, to know whether a particular security code belongs to the set of valid codes, either we have to match it with each member of the set to see if it tallies with any one of these, or there should be

some information within the code itself, which indicates its position in the set. So we don't have to search through the entire set for a match but we only need to see whether security code at that particular position matches or not. It was decided that the first three digits of the code shall indicate its position within the array of codes. The length of the rest of the code was chosen to be 10 digits, thus giving a 13 digit security code as shown in Figure 2-2.

The internal storage of each security code is 5 bytes in the format shown in Figure 2-3. The lower four bytes contain the binary equivalent of the number generated from the lower 10 digits of the security code. Bits of the 5th byte contain information telling whether this security code exists or not, whether it was illegally used and if yes in what manner as indicated in Figure 2-3.



First three digits tell the position of the security code in the security codes table.

These lower 10 digits are converted into a number for internal storage taking four bytes.

Figure 2-2. Format of Security Codes.

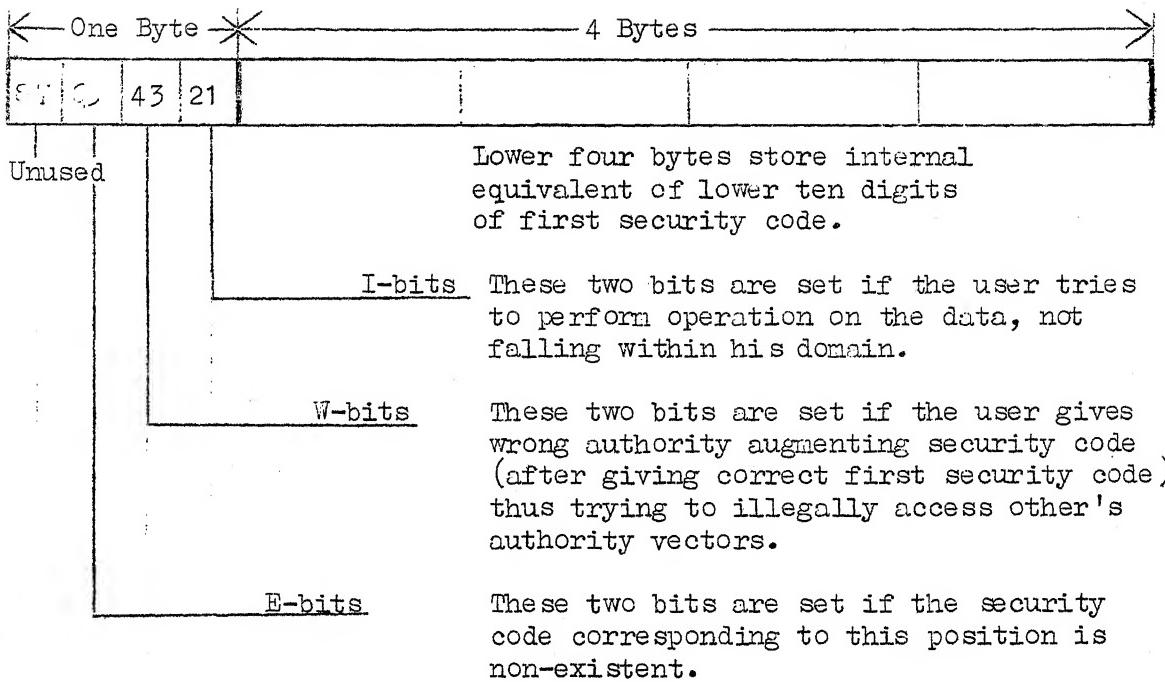


Figure 2-3: Format for internal storage of first security codes.

Since one security code takes five bytes for internal storage, so the number of security codes that can be accommodated in one sector of the disk is $\lfloor 256/5 \rfloor = 51$.

To keep track of the sectors which store first security codes, we have a table shown in Figure 2-4 called "FIRST Security Codes" TABLE DIRECTORY (FSECDC). As can be seen from the figure, each row of this table takes 14 bytes. Since this table itself is placed in a sector of the disk, so maximum number of rows of this table is $\lfloor 256/14 \rfloor = 18$. Since one row contains information about one sector (which may contain maximum of 51 security codes), so maximum number of First-Security-Codes that can be issued is equal to $18 \times 51 = 918$, which is a fairly large number and sufficient for all practical applications.

The columns of "First Security Codes' Table Directory" have the following meanings:

14 Bytes		
Cylinder No. CYLDRF	Sector & Surface No. SECTR	10 digits Random No. RDND
1	2 Bytes	2 Bytes
2		
3		
4		
:		
15		
16		
17		
18		

The ith row of the table contains relevant information for security codes whose first 3 digit number lie between ix51-50 and (ix51)

Figure 2-4: First Security Codes' Table Directory (FSECDC)

CYLDRF: This column contains the cylinder number of the disk where the sector containing information about the corresponding group of codes is stored. A non-existing entry in this table is indicated by a -ve value of this field for that particular row.

SECTR: This column contains the surface number and sector number of the sector of the disk containing information about the corresponding group of codes. These two values are packed into one word in the form shown in Figure 2-5, which is the same as that used by the disk controller.

15	9 8	5 4	0
Unused	Surface No. (0-9)	Sector No. (1-10)	

Figure 2-5: Format for putting sector No. & Surface No. in one word.

RDND: This column stores a 10-digit random number for each group of 51 security codes (i.e., one sector). This is used to generate another number from the lower 10 digits of the security code using the routine (WORBLE). The idea is to make the internal number stored on the disk different from the corresponding first security code so that it may be extremely difficult, if not impossible, for any person to compute the first security code of some one else, after he has somehow (though it shall not be possible through normal channels as the disk shall also be protected) read the internally stored security code from the disk. Since this routine WORBLE is changable at the discretion of the DBA, it will be very difficult to know, what transformation this routine is applying on the security code to generate another 10 digit number which is converted into its binary value for internal storage, by the routine GIVSEC.

The only information that is maintained in core, regarding the first security codes, is the address of the sector containing "First Security Codes' Table Directory". There are two variables (one word each) for keeping track of this information:

ISECYL : contains cylinder number of the disk where "First Security Codes' Table Directory" (FSECDC) is stored.

ISECTR : contains sector and surface number of the disk where FSECDC is stored (in the format shown in Figure 2-5).

ISECYL = 196
ISECTR = 35

- 001 003
15 9 8 5 4 0

FSECDG

Sector of disk in
cylinder No. 196,
surface No. 1,
sector No. 3

Cyl. No. 196
Surface No. 1
Sector No. 7.

Cyl. 197
Surface 0
Sector 1

Cyln. 197
Surface 0
Sector 5

1	48	-
2	:	:
15	48	-
16	0	1 5 9 6 8 2 3 4 5 7
17	48	-
	:	:
	:	:
27	48	-
28	03	2 3 0 5 8 2 6 4 1 7
29	0	0 6 8 3 0 0 5 0 2 8
30	48	-
	:	:
	:	:
43	48	-
44	0	3 2 5 8 8 6 6 9 7 3
45	48	-
	:	-
	:	-
51	48	-

	1	48	-
	.	.	.
	.	.	.
Sec.	15	48	-
code	16	0 123540723	
016...	17	12 086532432	
	18	48	-
	.	.	.
	.	.	.
	.	.	.
	.	.	.
	.	.	.
	.	.	.
	.	.	.
Sec.		.	.
code		.	.
=044		.	.
	51	48	-

One byte Four bytes

Sec.
Code
539..

Figure 2-6: (An Instance) of First Security Codes' Tables.

Figure 2-6 shows, how an instance of these tables may look at any moment. This also depicts how the mapping of a security code to its internal stored equivalent is done starting from ISECYL and ISECTR. In the Figure, -1 in the CYLDRF field indicates a non-existing entry. We see that in the figure only two entries exist and contents of these sectors show that only 7-security codes exist, out of which two have been used illegally and hence are unusable.

2. AUGMENTING SECURITY CODES

The augmenting security code also has a length of 13 digits with the first 3 digits giving information regarding the position of a code in the table of authority augmenting security codes.

The internal storage for each security code is 10 bytes as shown in the format in Figure 2-7.

2 Bytes	4 Bytes	4 Bytes
0	1 2	5 6

These two bytes contain the cylinder No.

These 4 bytes contain a scrambled no. generated from the 9-bit surface and sector No. using routine VORBLE.

These four bytes contain the internal equivalent of the lower 10 digits of the authority augmenting security code.

Figure 2-7: Format for Internal Storage of Authority Augmenting Security Code.

Again 4 bytes are taken up by the binary equivalent of the internal number generated from the lower 10 digits of the security code, as in the case of first security codes. The remaining six bytes are taken by two fields given below:

Cylinder No. (2 Bytes): We know that there is one authority augmenting vector stored on the disk corresponding to each existing authority augmenting security code. This field gives the cylinder number of the disk which contains the corresponding authority augmenting vector in one of its sectors. A -ve value of this field indicates that the security code corresponding to this position of the table has not been defined by the DBA.

Sector and Surface No. (4-Bytes): This field stores the surface and sector number of the disk (with cylinder number indicated by cylinder number field) where the corresponding authority augmenting vector is stored. Though 2 bytes are enough for this field (in the format shown in Figure 2-5), the reason for providing four bytes is as follows -

As the contents of authority augmenting vectors are of crucial importance, we would like it to be impossible for any intruder to be able to know where his authority augmenting vector is stored and to change the contents of any authority augmenting vector to give him access to the areas prohibited to him. Therefore, instead of storing the surface and sector number of the authority augmenting vector directly, we scramble it to produce another artificial number of 4 bytes from this 9-bit information using a routine (VORBLE) and put its value on the disk. There shall be a routine (DVORBL) to change this four byte value back to the same 9 bits of information. These two routines (VORBLE) and (DVORBL) are changable at the discretion of the DBA but should be such that for each 9-bit input to routine (VORBLE), the 4-byte output produced by it should be such that when fed to the routine (DVORBL), it produces as output the same pattern of 9 bits which was given as input to routine (VORBLE).

One thing to be noted is that there is no information stored regarding wrongful use of any authority augmenting security code. This information is purposeful only for the first security codes so that the users trying to use the system in any illegal manner may be marked and asked for explanations.

With 10 bytes of internal storage being needed for each Authority Augmenting Security Code, the number of codes about which information can be maintained in a sector of the disk is $\lfloor 256/10 \rfloor = 25$. Again to keep track of the sectors which store authority augmenting security codes, we have an AUGMENTING SECURITY CODES' TABLES DIRECTORY (ASECDC) as shown in Figure 2-8. This has exactly the same format as FSECDC shown in Figure 2-4 and the fields have exactly similar explanations. It is also accommodated in a sector of the disk and has maximum of 18 entries. So maximum number of authority augmenting security codes that can be issued is $18 \times 25 = 450$.

Again the only information that is maintained in core, regarding authority augmenting security codes, is the address of the sector containing Augmenting Security Codes' Table Directory. The two variables (one word each) holding this information are as below -

ASECYL: contains the cylinder number of the disk where the Augmenting Security Code's Table Directory (ASECDC) is stored.

ASECTR: contains the surface and sector number of the disk where ASECDC is stored (in the format of Figure 2-5).

14 bytes		
Cylinder No. CYLDRA	Sector & Surface No. SECTRA	10-digit Random No. RDNDA
2 bytes	2 Bytes	10 Bytes
1		
2		
3		
4		
..		
15		
16		
17		
18		

Figure 2-8: Augmenting Security Codes' Table Directory
(ASECDC)

Figure 2-9 shows an instance of authority augmenting security codes tables with only 3 authority augmenting security codes existing.

Since the DBA (data base administrator) has also to approach the system through a "first security code" (no separate provision being made for him), to identify the first security code of the DBA, we have a variable DBSC.1 (one word) which stores the first three digits (converted to equivalent binary form) of DBA's, first security code.

ALGORITHM FOR SECURITY CHECKING

When a user indicates his interest in using the data base system, only a small part of the programs get loaded from disk into main memory. This part of the program asks the user to give his security codes and then if found valid, it loads the appropriate programs depending upon whether the user is a Data Base Administrator or some other user. If user is the Data Base Administrator, then the programs loaded are as below:

ASECYL = 200

ASECTR = 03

-	0	3		
9	8	5	4	0

ASECDC

Sector of the Disk in Cylinder	CYLDRA	SECTRA	RDNDA
1	-1	-	-
2	-1	-	-
3	197	08	5678432109
4	-1	-	-
5	-1	-	-
6	-1	-	-
7	-1	-	-
8	-1	-	-
9	-1	-	-
10	-1	-	-
11	-1	-	-
12	-1	-	-
13	-1	-	-
14	-1	-	-
15	-1	-	-
16	-1	-	-
17	-1	-	-
18	-1	-	-

Cylinder No. 197,
Surface No. 0,
Sector No. 8

Cylinder No.	Sector & Surface No. (Scrambled value)	Security Code
1	-1	-
2	-1	-
3	198	-
4	-1	Scrambled value of 69(2,5)
...
16	-1	-
17	196	Scrambled value of 100(3,4)
18	-1	-
19	-1	-
20	199	Scrambled value of 04(0,4)
...
25	-1	-

Cylinder No. 198
Surface No. 2
Sector No. 5

Authority
vector for
Sec. code
= 105.....

Cylinder No. 196
Surface No. 3
Sector No. 4

Auth.
Aug.
Vec.
for Sec.
Code =
119....

Cylinder No. 195
Surface No. 0
Sector No. 4

Auth.
vec.
of sec.
code =
122....

Figure 2-9: An Instance of Authority Augmenting Security Codes' Tables.

1. Program for deleting any security codes.
1. 2. Programs for changing any security codes or authority vectors.
3. Program for inserting new security codes.
2. 4. Program for printing of the security codes tables.
3. 5. Program for initializing the security codes tables.
4. 6. Program for building of the data base (setting up new relations)
5. 7. Program for removing a relation from the data base.

Otherwise the programs loaded are as below -

1. Program for retrieving any information from the data base.
2. Program for updating information in the data base.
3. Other routines and data used by above two programs.

The flowchart of the program which asks for security codes and loads appropriate programs is shown in Figure 2-10. The following paragraph supplements the information given in the flowchart.

The initial data is the value of ISECYL, ISECTR, ASECYL, ASECTR, DBSC.1, data which tells from where in disk various programs like update, retrieve, security programs, build programs etc. are to be loaded and to which areas of main memory they go. (Since the object program are in absolute machine languageform). Only the assembly language equivalent of flowchart (Figure 2-10) is incorporated under KDM with code "DB". This program does rest of the job itself.

The program for retrieval and updating are developed under the separate theses (Reference 1 and 7). The present thesis along with (Ref. 5) where programs for building the data base are described completes the functions carried out by the DBA. Flowchart in Figure 2-11

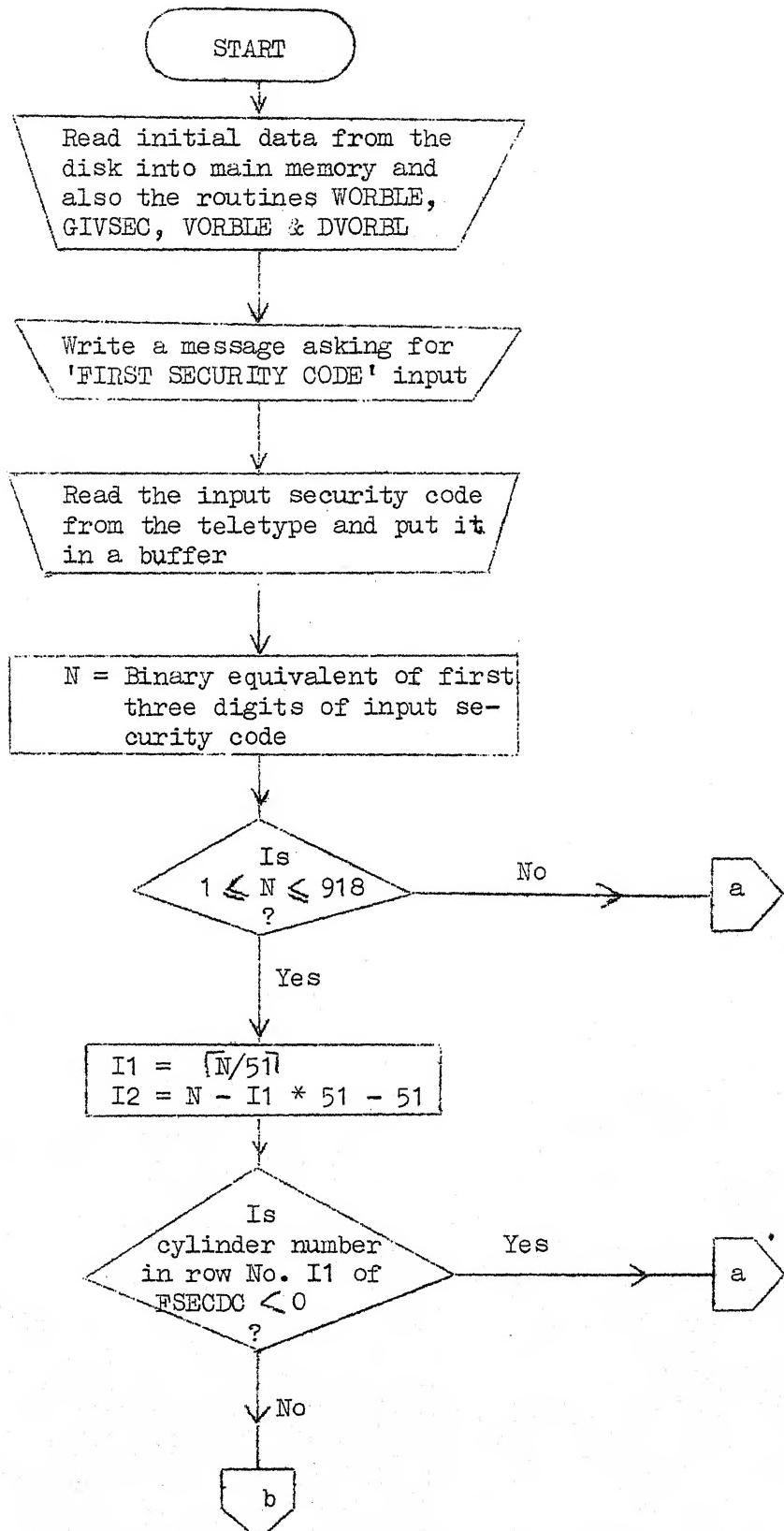


Figure 2-10: Continued on next page.

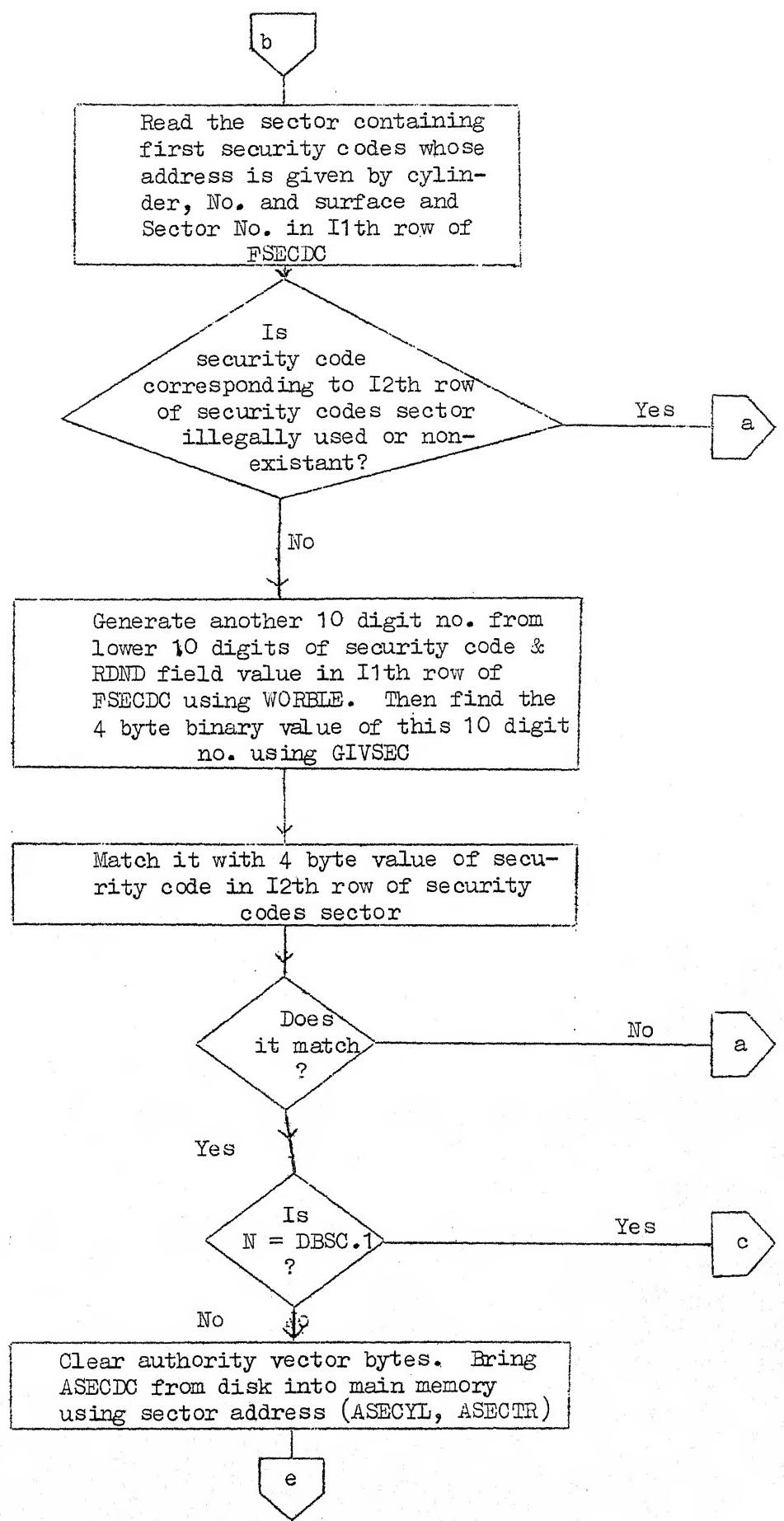


Figure 2-10: Continued on next page.

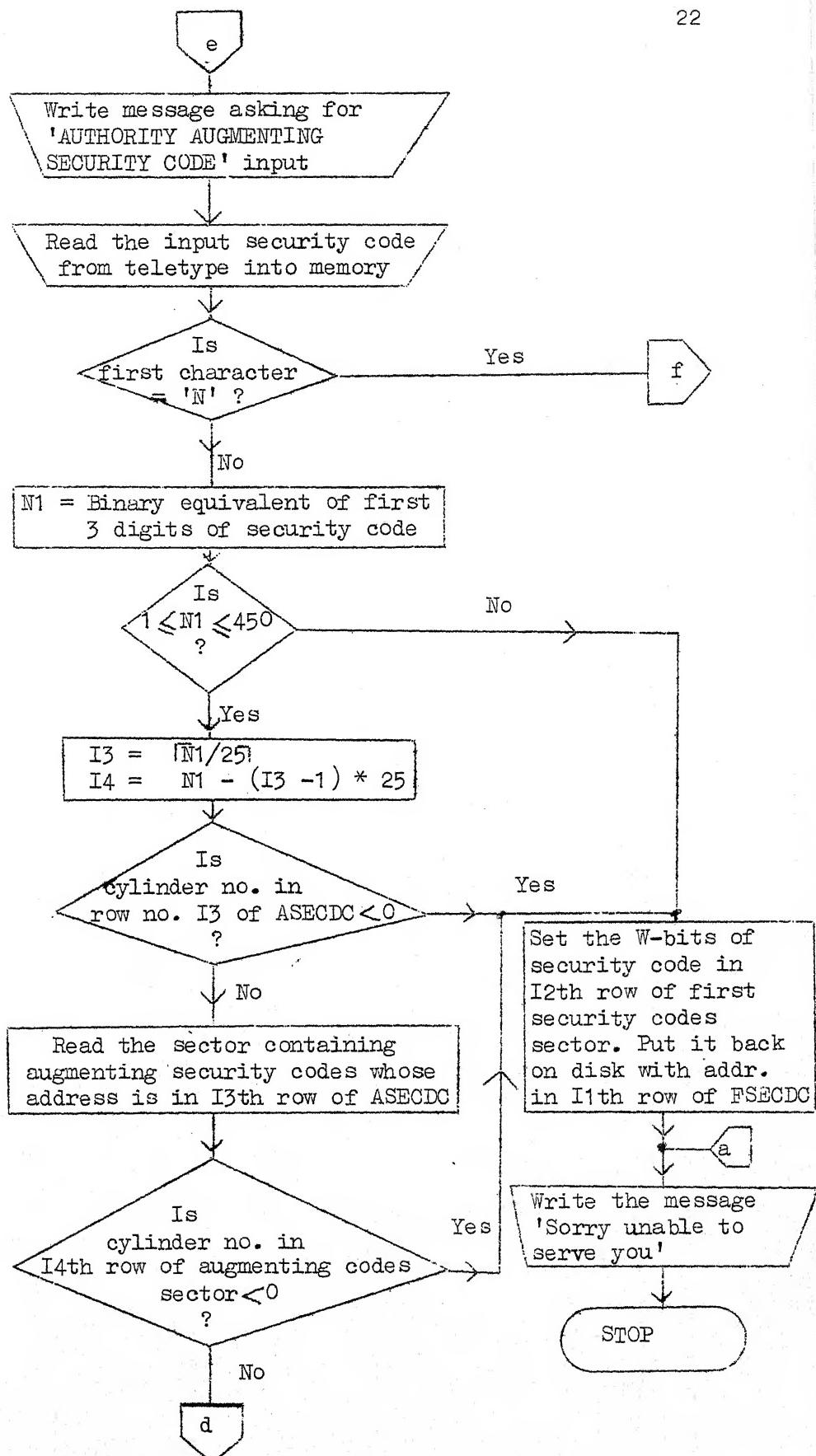


Figure 2-10: Continued on next page.

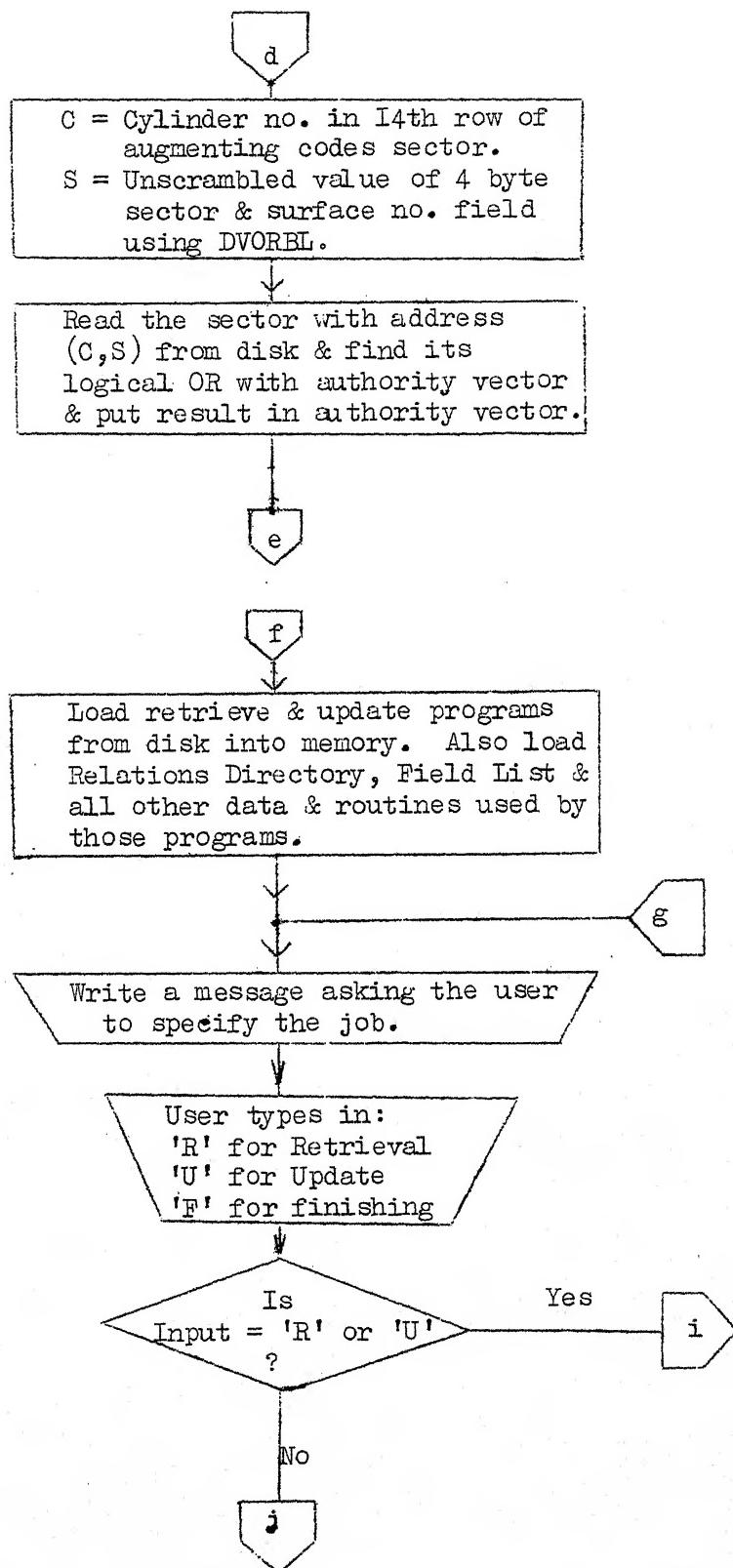


Figure 2-10: Continued on next page.

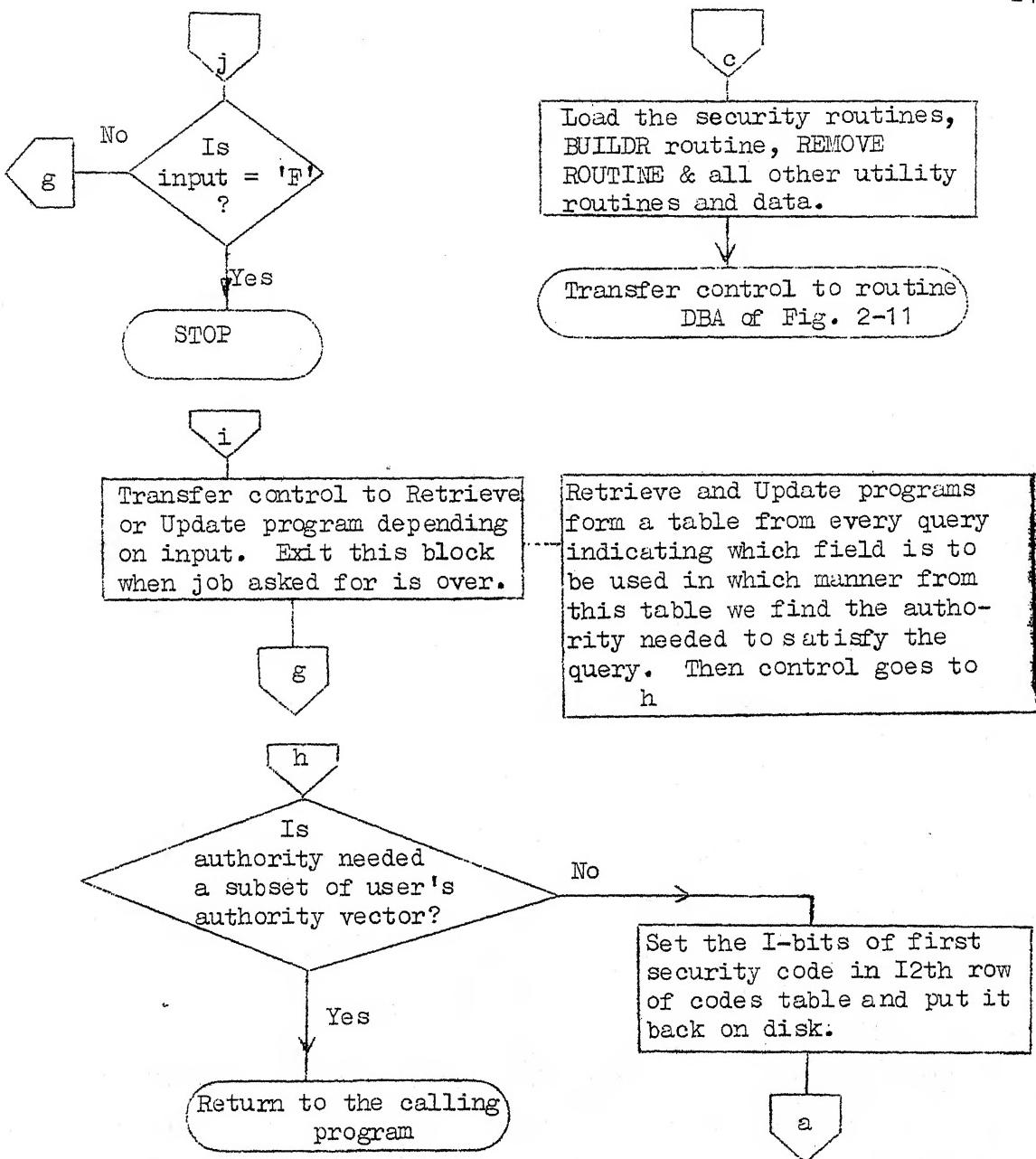


Figure 2-10: Flowchart for checking Security Codes:

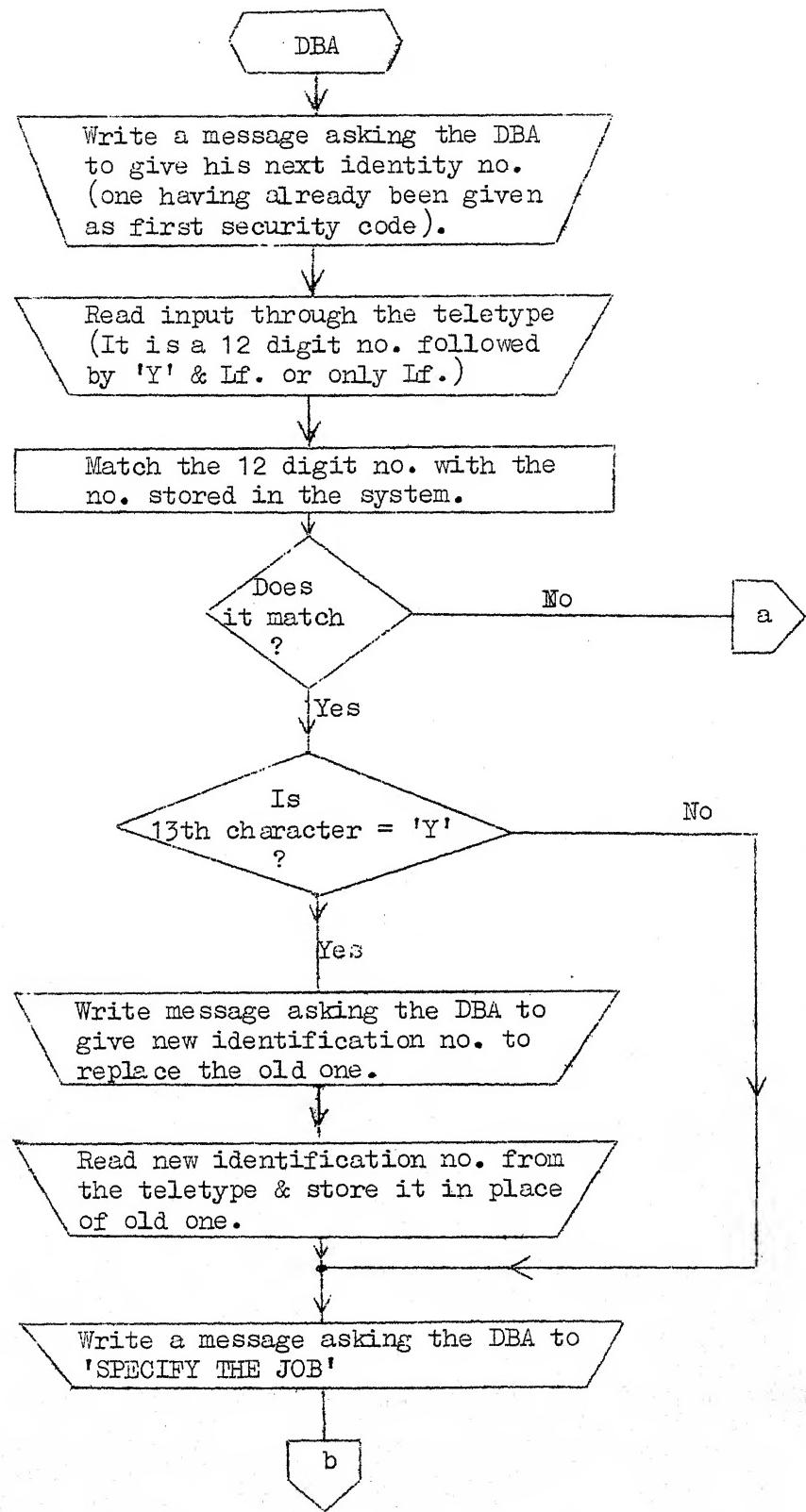


Figure 2-11. Continued on next page.

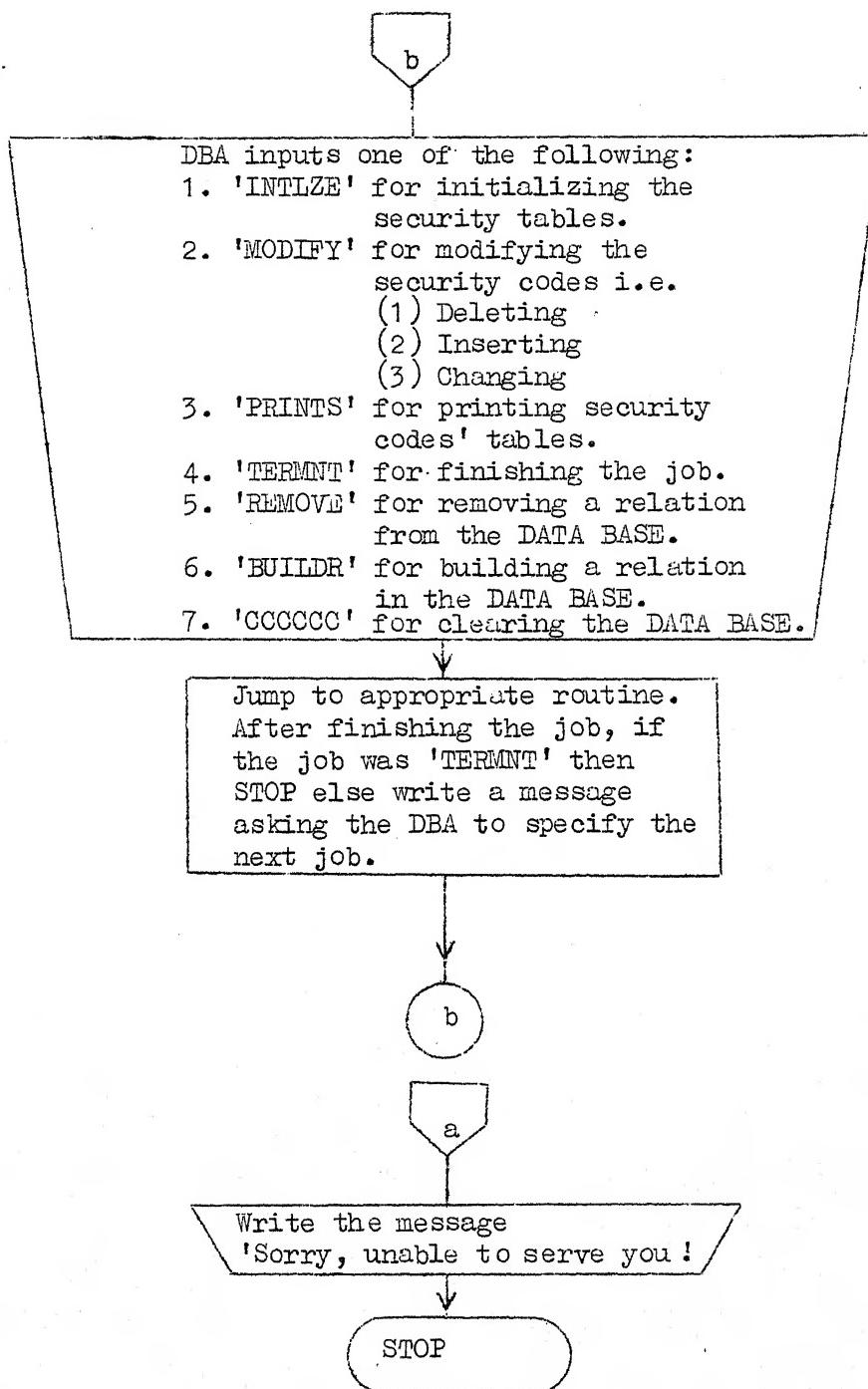


Figure 2-11: Flow chart for processing DBA's request.

shows the path that the DBA has to go through to carry out his task. The flowcharts for these tasks are detailed in later chapters in the following sequence.

Chapter 3 - Initialising the Security Tables

Chapter 4 - Modifying the Security Tables.

Chapter 5 - Printing the Security Tables.

Chapter 6 - Removing a Relation.

3. INITIALISING THE SECURITY TABLES

As described in the previous chapter, there are two types of security codes -

1. First Security Codes.
2. Authority Augmenting Security Codes.

Both these security codes involve two types of tables for internal storage as seen in last chapter which are -

1. Codes' Table Directory
2. Codes Table .

Codes' Table Directory gives pointer to the sector of the Codes Table in appropriate sequence depending upon the group of codes that the sector contains. Initialisation routine clears all the authority augmenting security codes from the tables and also clears all the first security codes except the first security code of the DBA which is retained to allow him continued access, for adding new security codes to initialized tables.

Initializati on is essential when the DBA wants to change any of the routines WORBLE, VORBLE and DVORBLE etc., because now the internally stored security codes and surface and sector number in ASECDC will be computed using new routines and therefore the tables must be initialised and filled all over again. Also in this case, when the routines are changed, internally stored first security code of the DBA has also to be changed to the value produced by these new routine while acting on DBA's first security code and this is done at the same time when routines are changed.

Since all the security codes and authority augmenting vectors are stored on the disk, some space needs to be assigned on the disk for the purpose of storing these. Because maximum of $1+18+1+18+450 = 488$ sectors of security space may be needed, we reserve 5 cylinders (500 sectors) on the disk and call this reserved space as "Disk Security Space". To manage the disk security space, we store the status of each sector (free or full) in a bit and a 496 bit vector (an array of 31 words starting from the address BITMP) reflects the status of disk security space at any time. The routine AIOCTS allocates a sector of the disk security space, makes its status bit 1 and return the address of this sector to the calling program. The routine RELSEC releases a sector of disk security space with sector address given by the calling program by clearing its status bit. The program for these two routines appear in the Program Listings.

INITIALIZATION ALGORITHM

This algorithm in the form of a flowchart is shown in Figure 3-1. The equivalent program appears in the Program Listings. It deletes all security codes except DBA's first security code from the security tables. It also changes the disk storage addresses of FSECDC and ASECDC by allocating them new sectors and releasing the sectors previously held by them, thus making the positions of these tables changeable. After initialisation, new security codes can be added using the MODIFY routine in the manner explained in the next chapter.

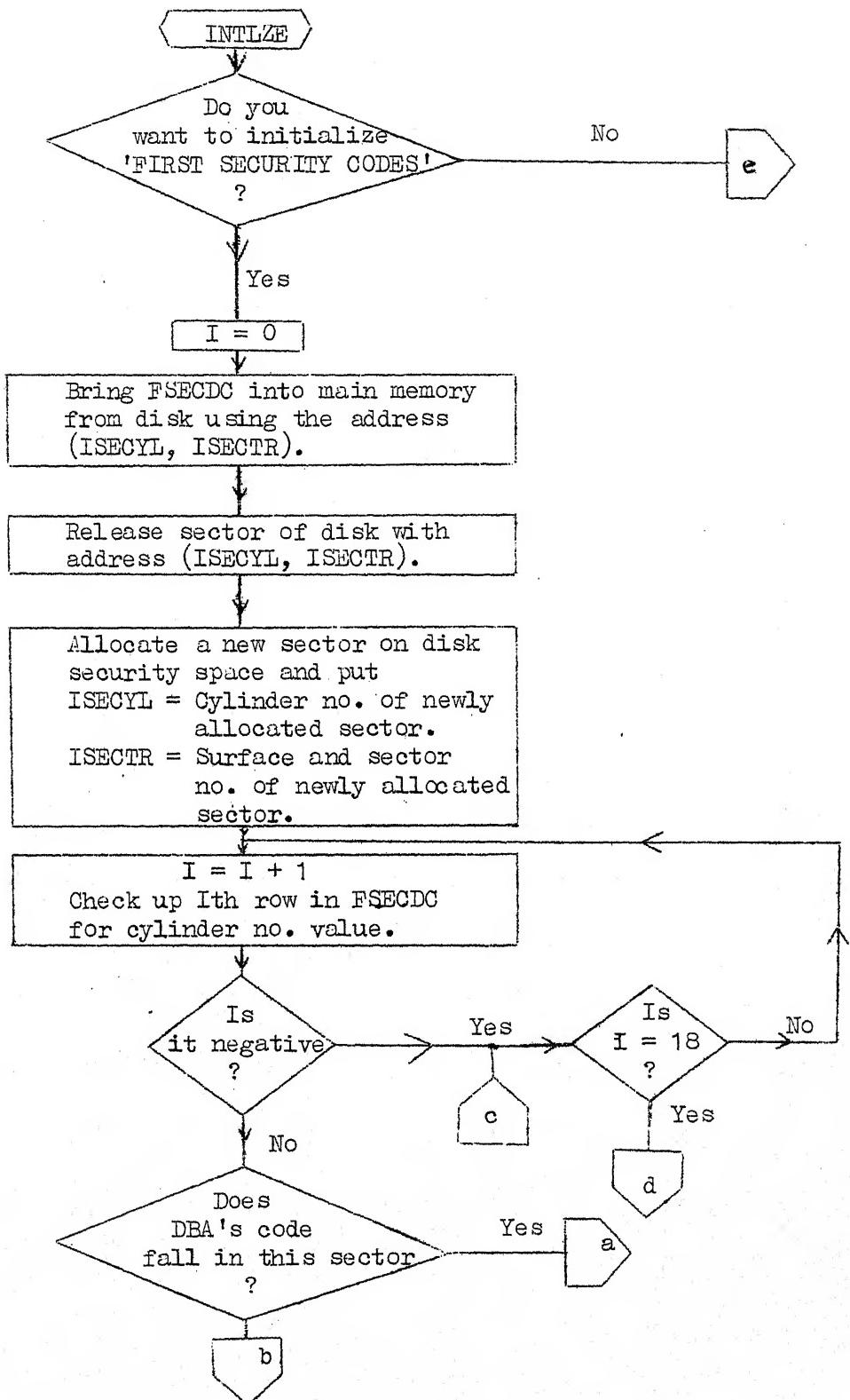


Figure 3-1: Continued on next page.

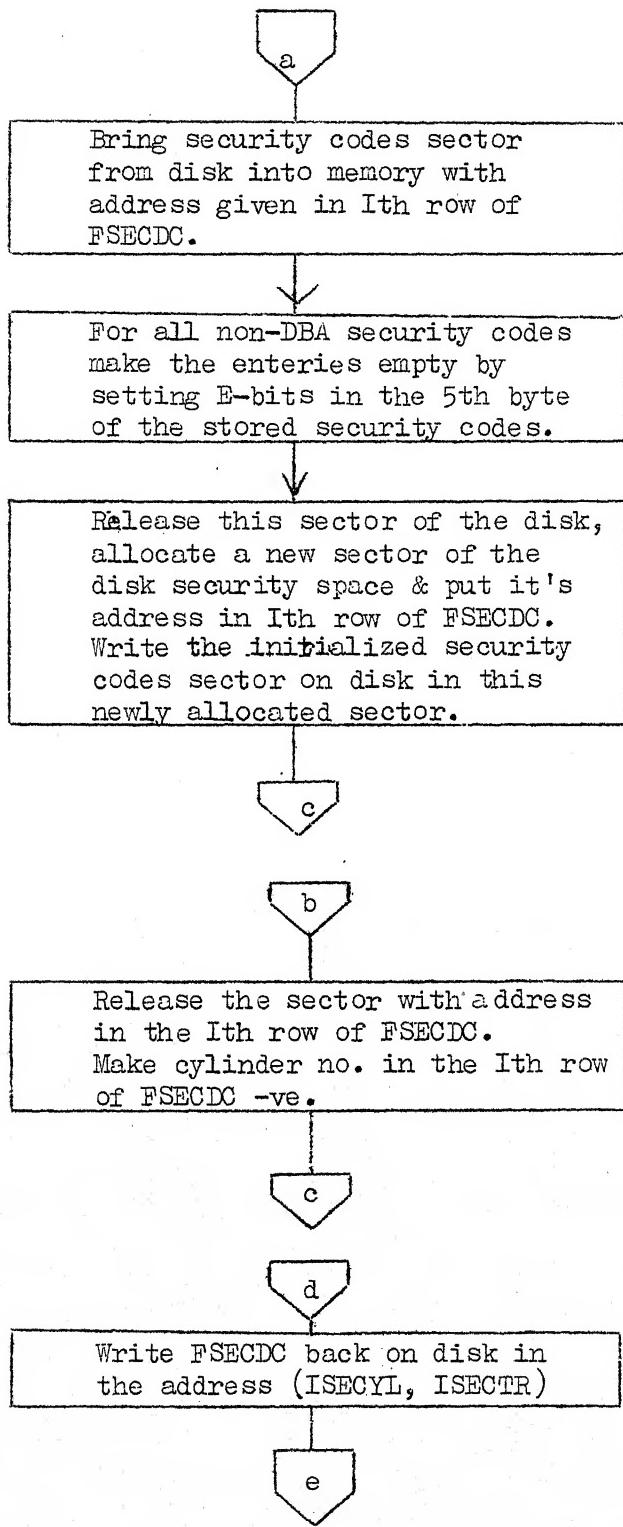


Figure 3-1: Continued on next page.

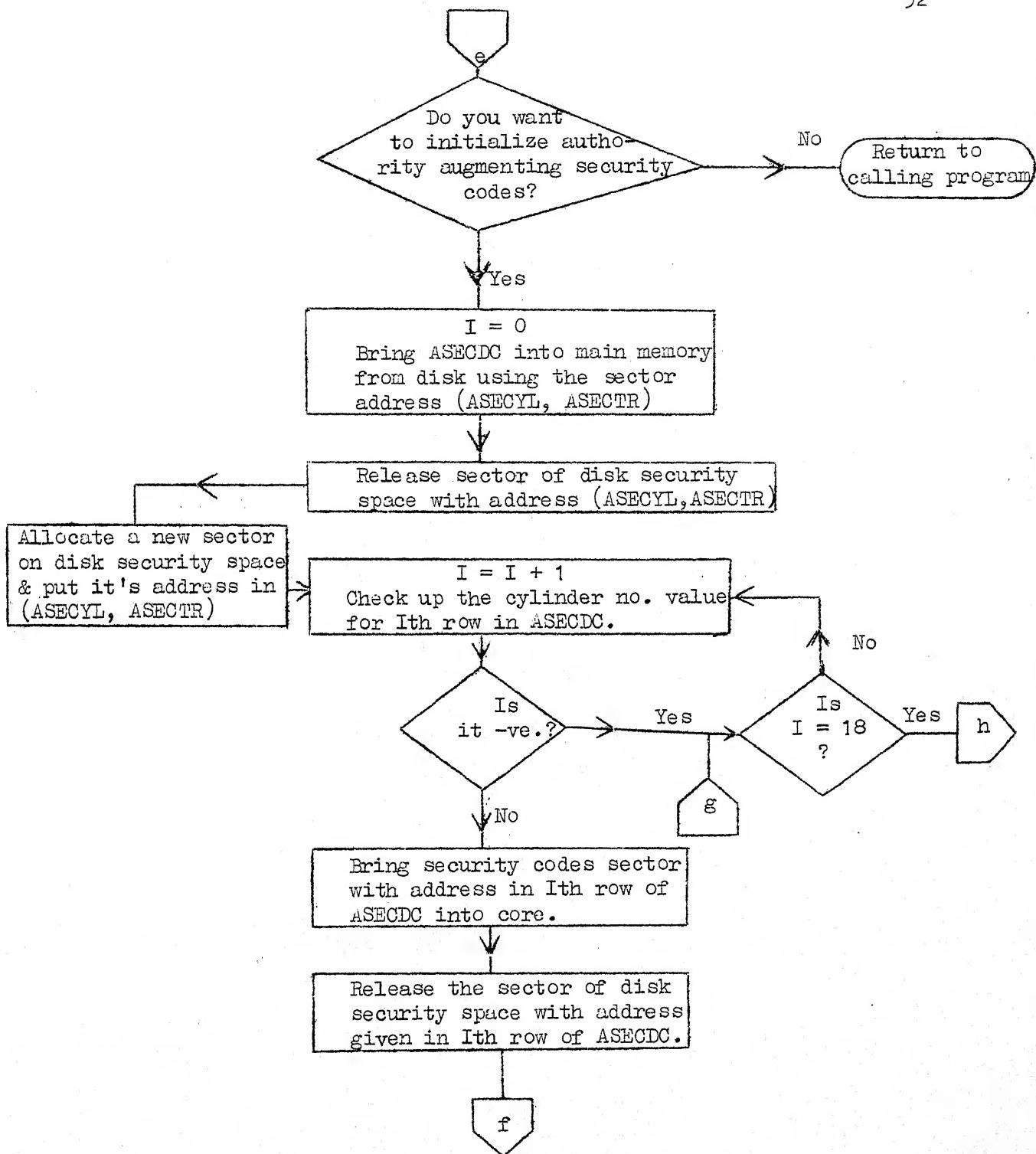


Figure 3-1: Continued on next page.

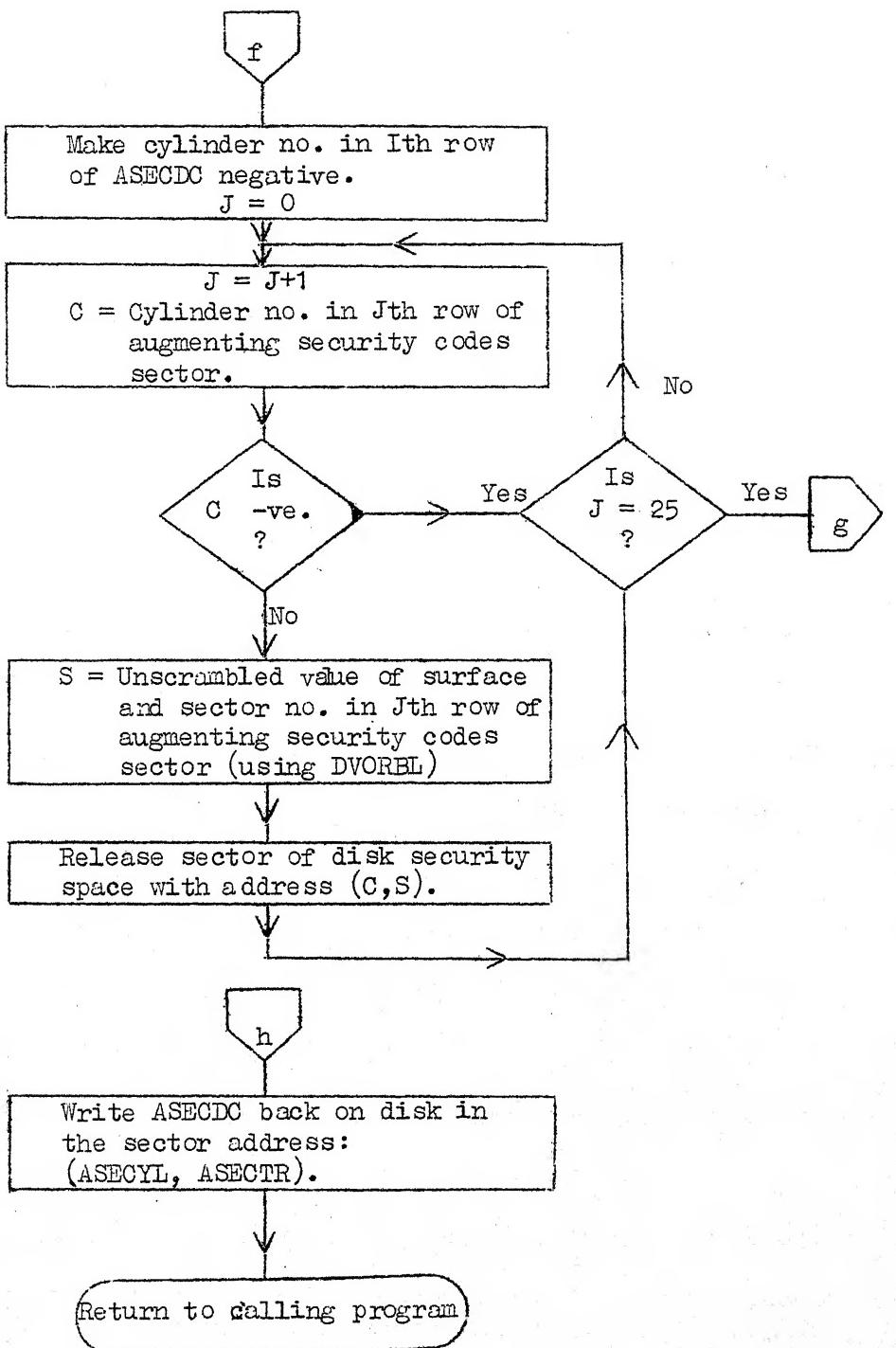


Figure 3-1: Flowchart for initializing the security codes.

4. MODIFYING THE SECURITY TABLES

Modification of the First Security ^{Code's} Tables can be done in any of the following ways:

1. Deleting a first security code
2. Inserting a new first security code or replacing an existing one.

Modification of the authority augmenting security code can be done in any of the following ways:

1. Deleting the security code (and hence also its associated authority augmenting vector).
2. Inserting a security code and its associated authority augmenting vector.
3. Changing an already existing authority augmenting security code while leaving its associated augmenting vector intact.
4. Changing the authority augmenting vector while leaving its associated security code (already existing) intact.

Deletion of any security code (of either type) can be done by specifying only its first three digits i.e., its position. The syntax of the language statements for deleting a security code is described below in terms of its fields:

First field (one ch) = 'D' indicating deletion of security code.
2nd field (one ch) = '1' if first security code is to be deleted
= '2' if authority augmenting security code is to be deleted
3rd field (3 chars) = '3 digit No.' indicating the position of the security code.
4th field (one char.) = "." It is statement terminator.

For inserting a new security code or changing the value of an existing security code or authority vector, the syntax of the language statements is described below in terms of its fields -

First field (one character) = 'E' indicating insertion of new security code or authority augmenting vector or both.

Second field (one character) =

'1' if first security code is to be inserted or changed.

'2' if value of an existing authority augmenting security code is to be change leaving its associated authority augmenting vector intact.

'3' if a new authority augmenting security code with its associated vector is to be inserted.

'4' if the vector associated with an existing authority augmenting security code is to be changed while leaving the security code intact.

Third field =

'13 digit security code' (for second field = 1,2 or 3)

'3 digit No.' indicating position of security code (for second field = 4)

Fourth field =

Authority vector with syntax explained below (for second field = 3 or 4)

Null (for second field = 1 or 2)

Fifth field (one character) = '.' It is statement terminator.

The syntax of the authority vector (fourth field above) is expressed by the regular expression $(\#N, (A,)^*A)^*$ where -

N represents byte No. of the vector from where to start in decimal.

A represents contents of corresponding bytes of authority vector in octal.

e.g., "#10,011,023,054 #55, 111, 321, 12" corresponds to the following contents of authority vector bytes.

<u>Byte No.</u>	<u>Contents in Octal</u>
1 to 9	000
10	011
11	023
12	054
13 to 54	000
55	111
56	321
57	012
58 to 256	000

Now we shall see examples for all types of statements alongwith their explanations:

Example 1: "D1 123!"

This statement shall delete first security code whose first three digits are 123.

Example 2: "D2 421." : This statement shall delete authority augmenting security code whose first three digits are 421.

Example 3: "E1 159111111111." : This statement inserts the first security code "159111111111" in the set of valid first security codes. If a first security code starting with "159" digits already exists, the new security code takes its place and previous code is removed.

Example 4: "E2 126111111111." : This statement changes an already existing authority augmenting security code while leaving the associated vector intact. If no authority augmenting code with positional digits "126" exists, this statement is ignored and an error message is printed.

Example 5: "E3 021111111111 ≠ 10,011,230 ≠ 20, 11." : This statement introduces an authority augmenting security code "021111111111" in the list of valid codes with the associated authority augmenting vector as below:

<u>Byte No.</u>	<u>Contents in Octal</u>
1 to 9	000
10	011
11	230
12 to 19	000
20	011
21 to 256	000

If an augmenting security code with positional digits "021" already exists, then the value of this security code and its associated vector are replaced by new values of both.

Example 6: "E4 111 #10,011,230#20,11." : This statement changes the authority vector associated with the authority augmenting security code whose positional digits are "111" from present value to that shown in Example 5 above, while leaving the security code value intact. If no security code with positional digits "111" exists already, this statement is ignored and an error message is printed.

Any number of these modification statements can be given at a time through the input device which can be a card reader ("CA"), high speed paper tape reader ("PA") or keyboard ("KY"). More than one statement can be put on one card and a single statement can be continued on many cards. All columns (1-80) are usable. There can be arbitrary number of blanks between various fields of a statement and between statements. To indicate that no more statements follow, we put the character "F" after period of the last statement.

MODIFICATION ALGORITHM

Assumptions made -

1. GET routine is available. This routine when called puts the next character of the input in the variable "CH ". When calling this routine for the first time in the program, we clear the variable FLAG of this routine to indicate fresh input. Also the variable INPUT is loaded with the device number of the input device before calling this routine.
2. GET.3(N) routine is available which reads the 3 characters (digits) of input and puts their binary value in N.
3. GET.10(N1) routine is available which reads ten characters from input and puts them in a buffer starting from location N1.
4. WORBLE (N1,N2) routine is available which generates a ten digit number from two 10 digit number inputs from buffers N1 and N2 and puts the generated result in buffer N1.
5. GIVSEC (N1,N2) routine is available which takes 10 byte (digit) input from the buffer N1 and puts its 4 bytes binary equivalent in buffer N2.
6. AUTHSC routine is available which reads the authority vector part of the input modification statement and builds its equivalent authority vector in a specified area of the core.

The program for all the six routines above appear in the Program Listings.

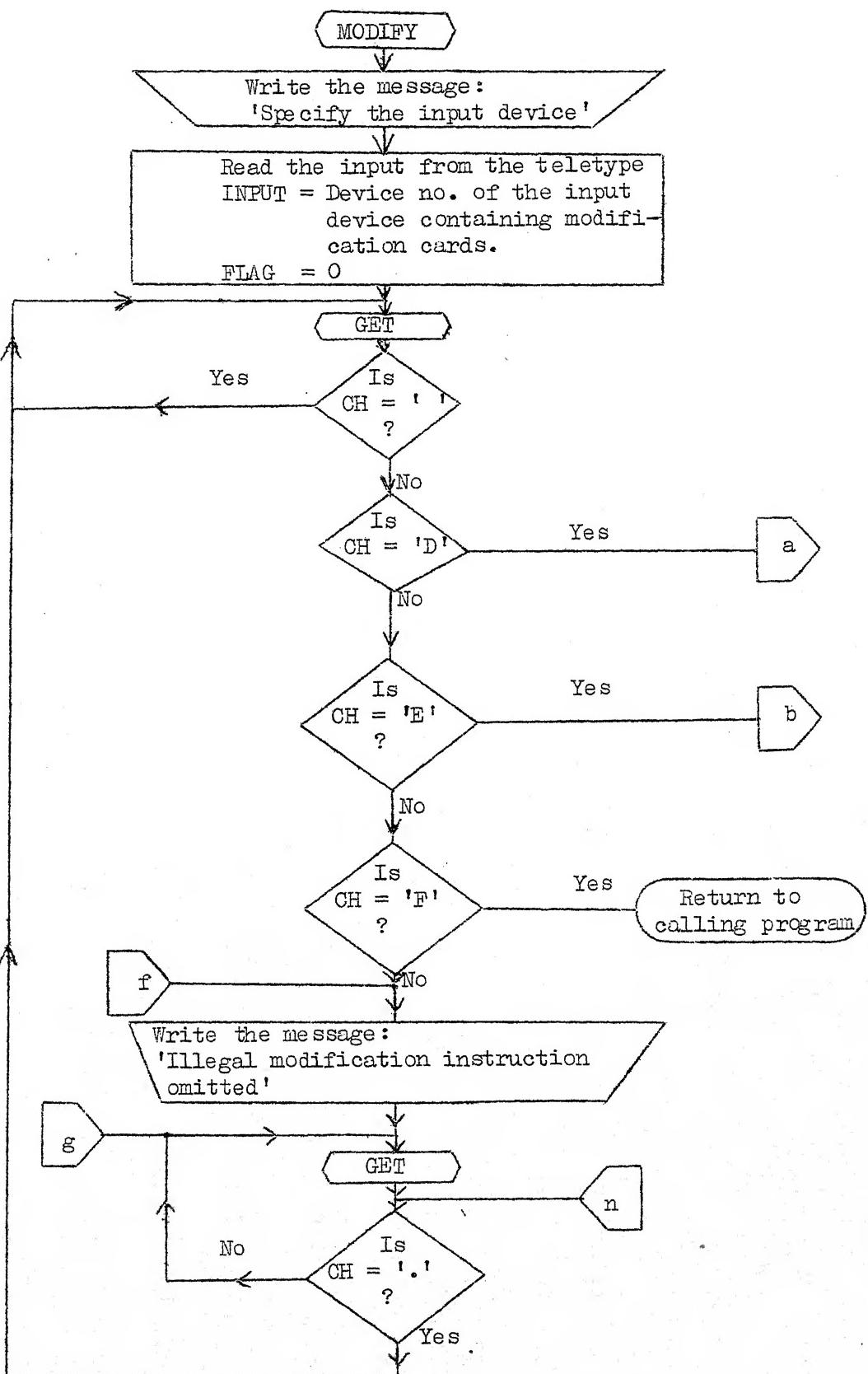


Figure 4-1: Continued on next page.

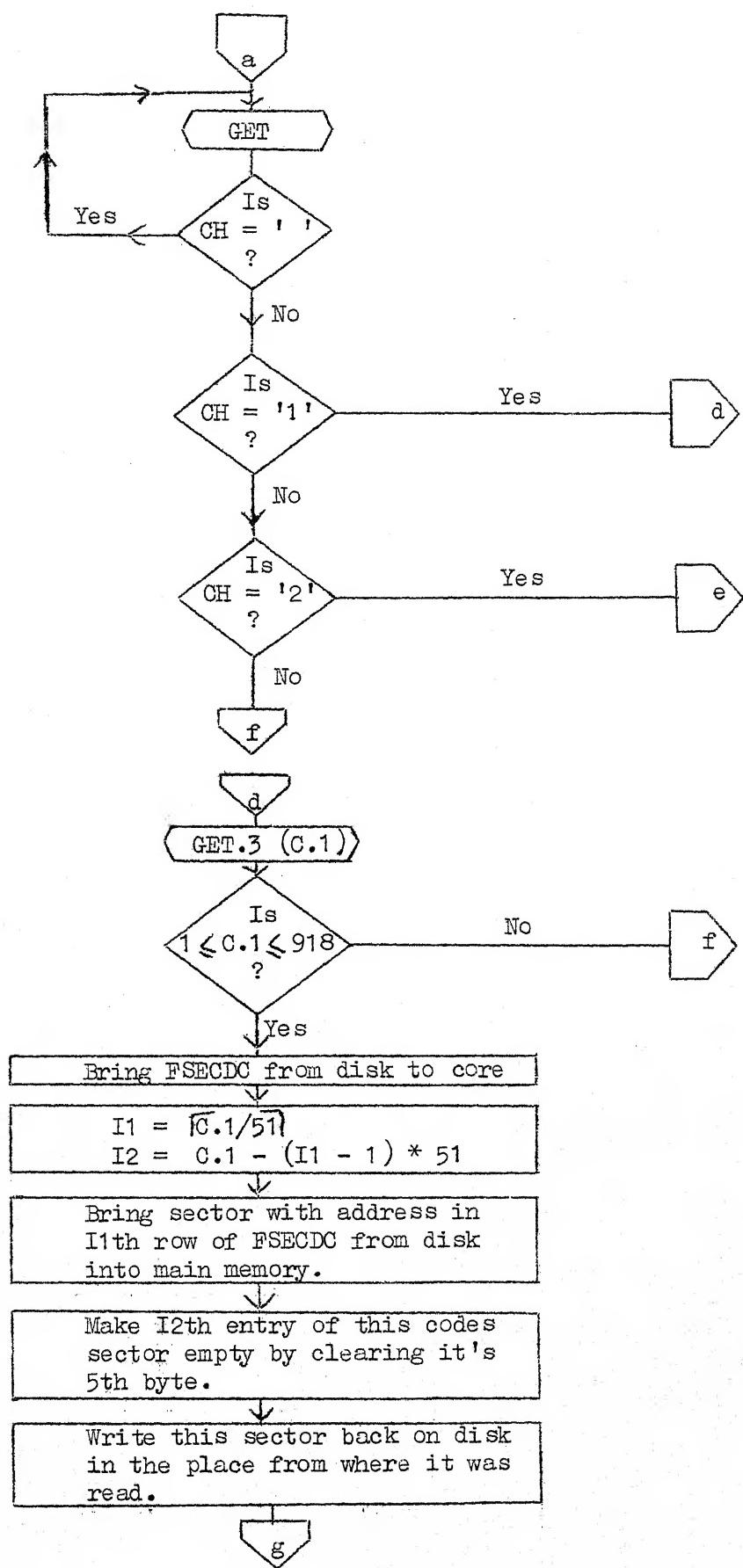


Figure 4+1: Continued on next page.

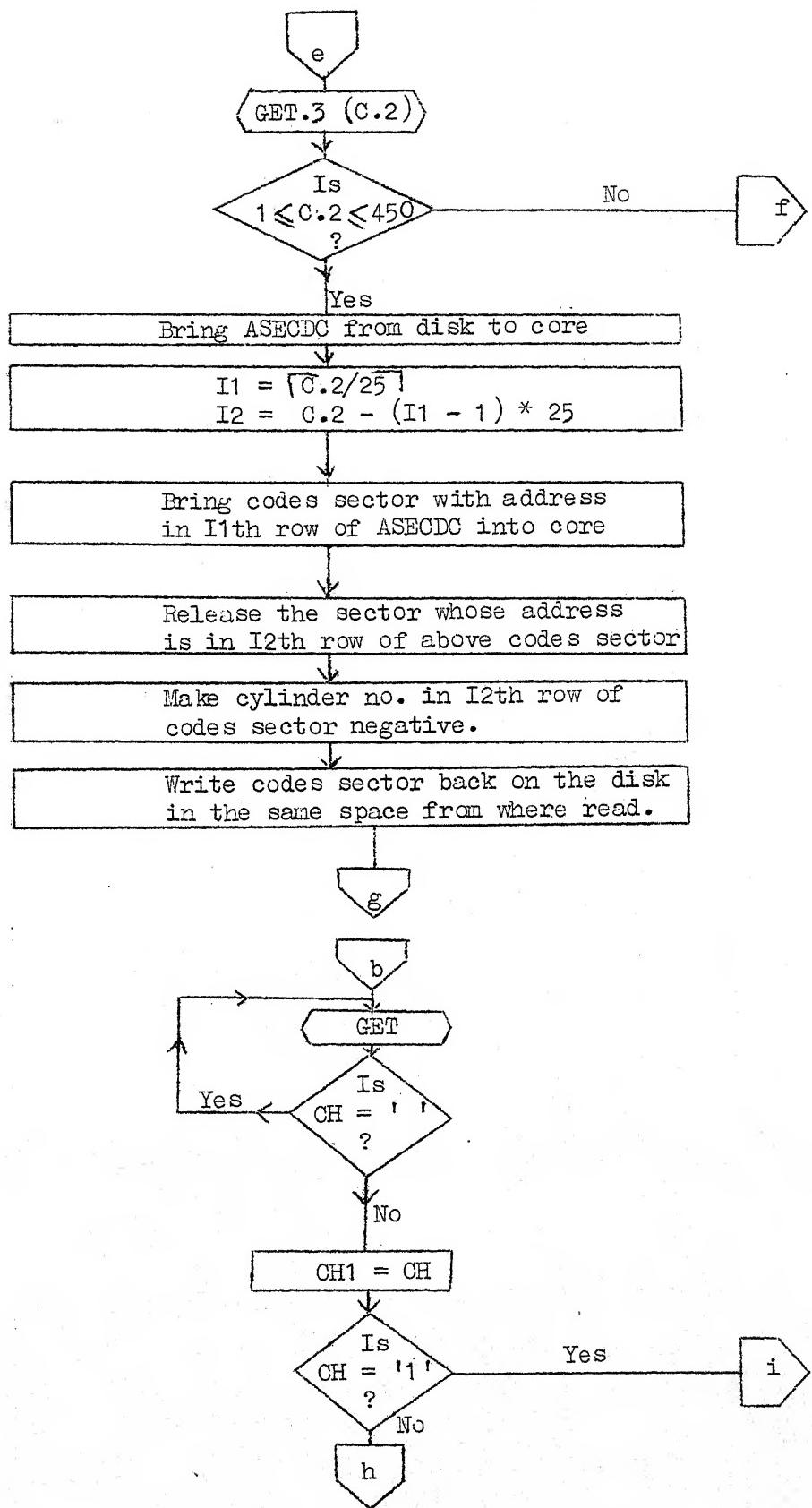


Figure 4-1: Continued on next page.

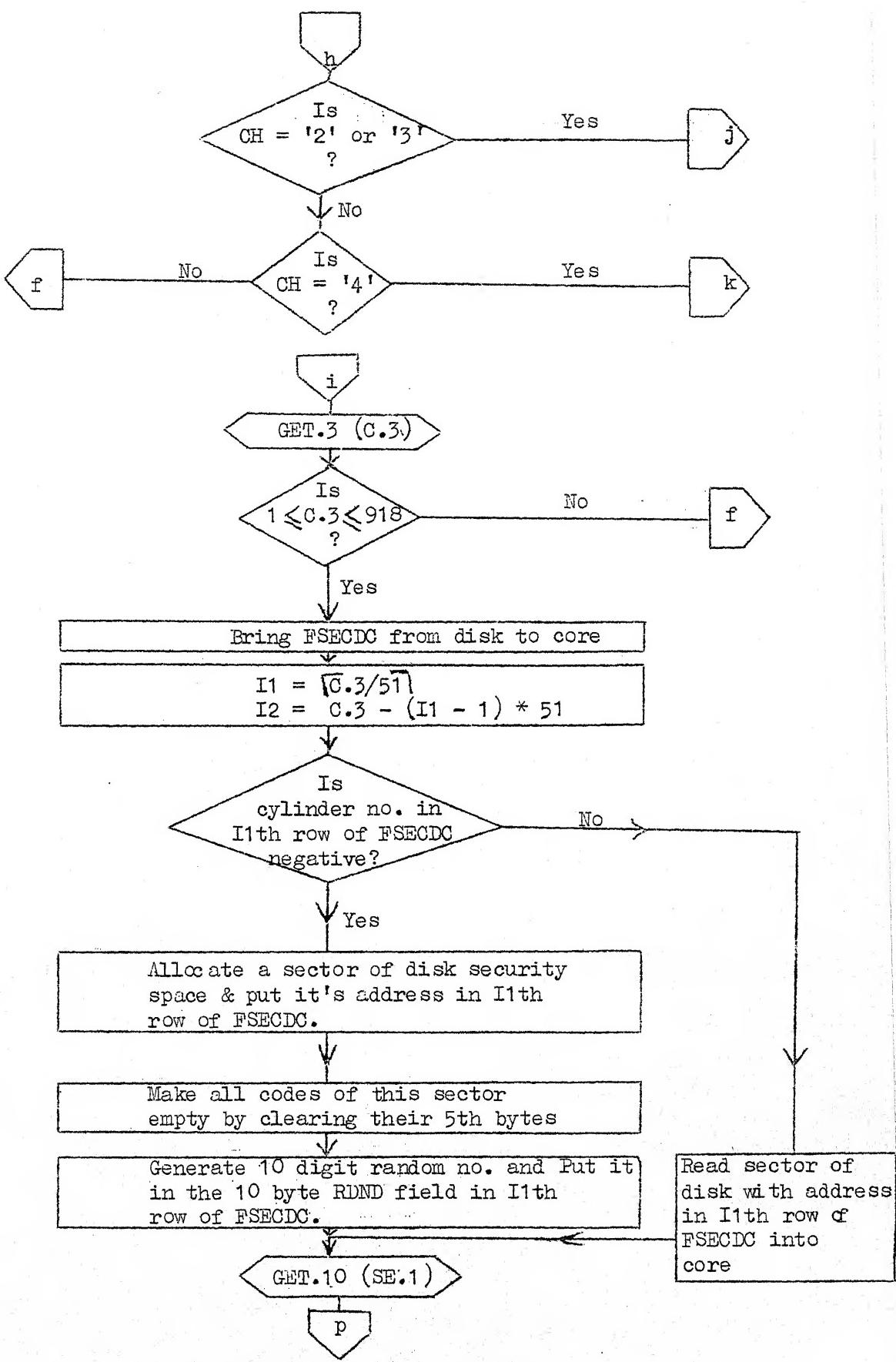


Figure 4-1: Continued on next page.

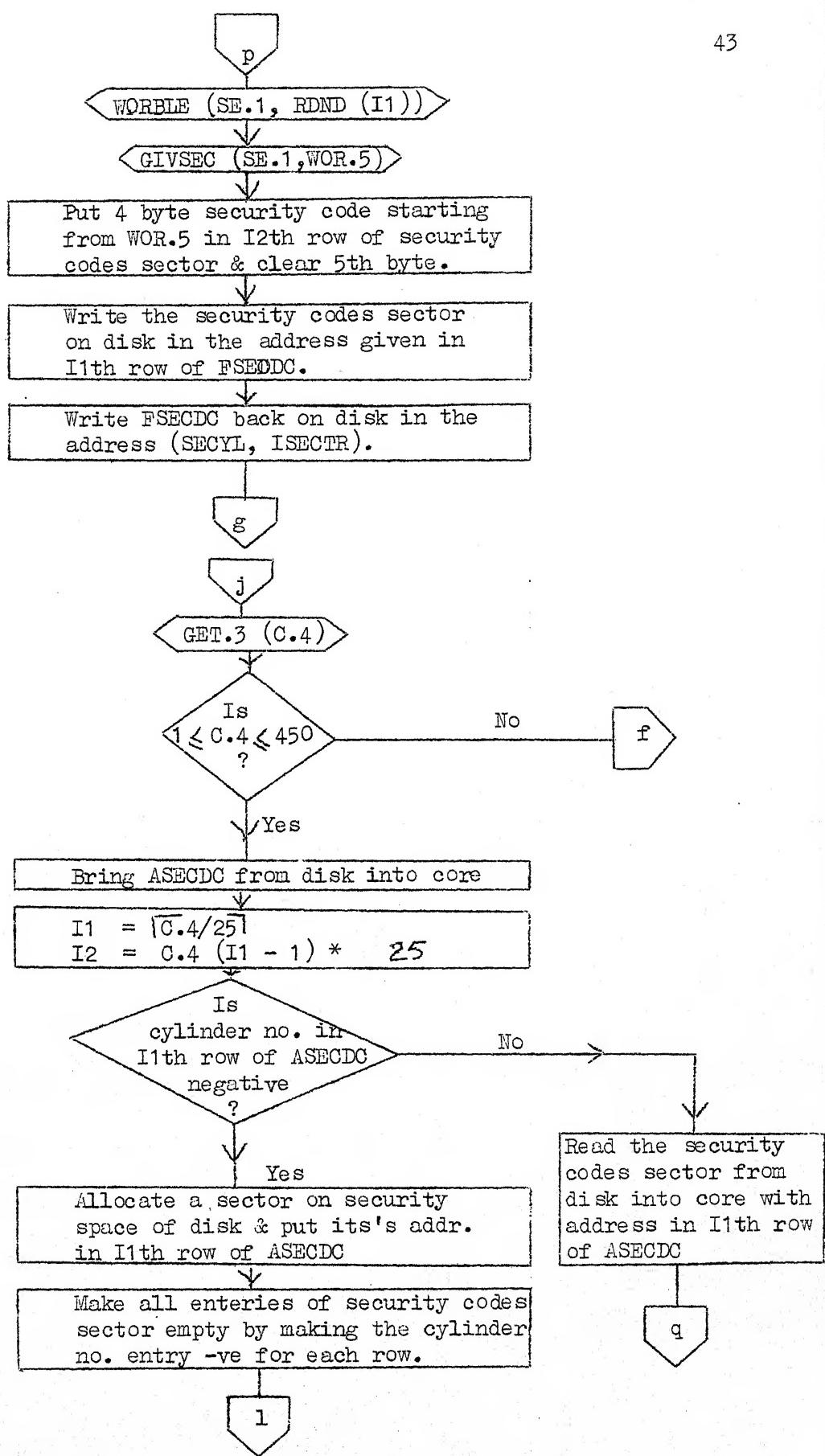


Figure 4-1: Continued on next page.

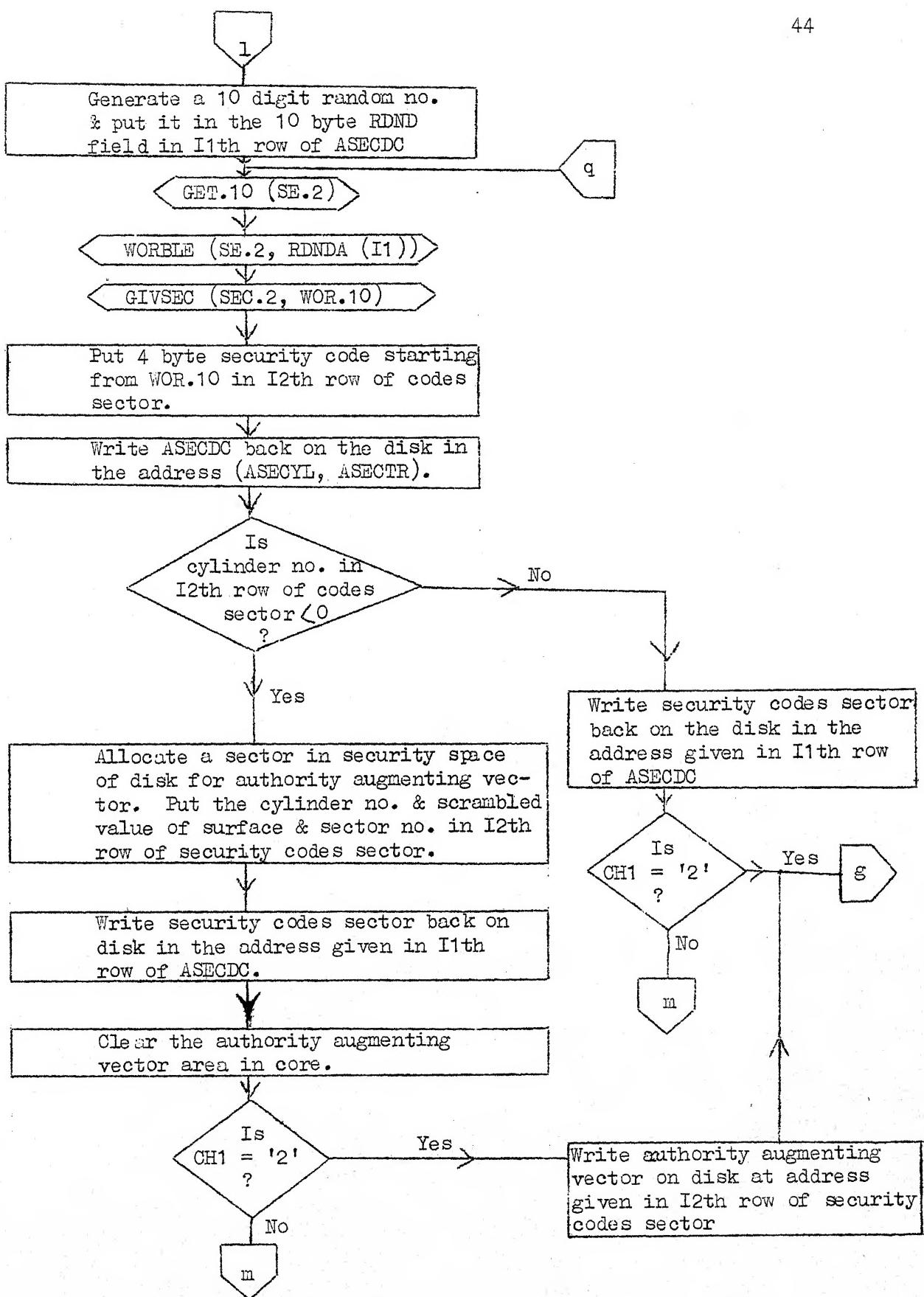


Figure 4-1: Continued on next page.

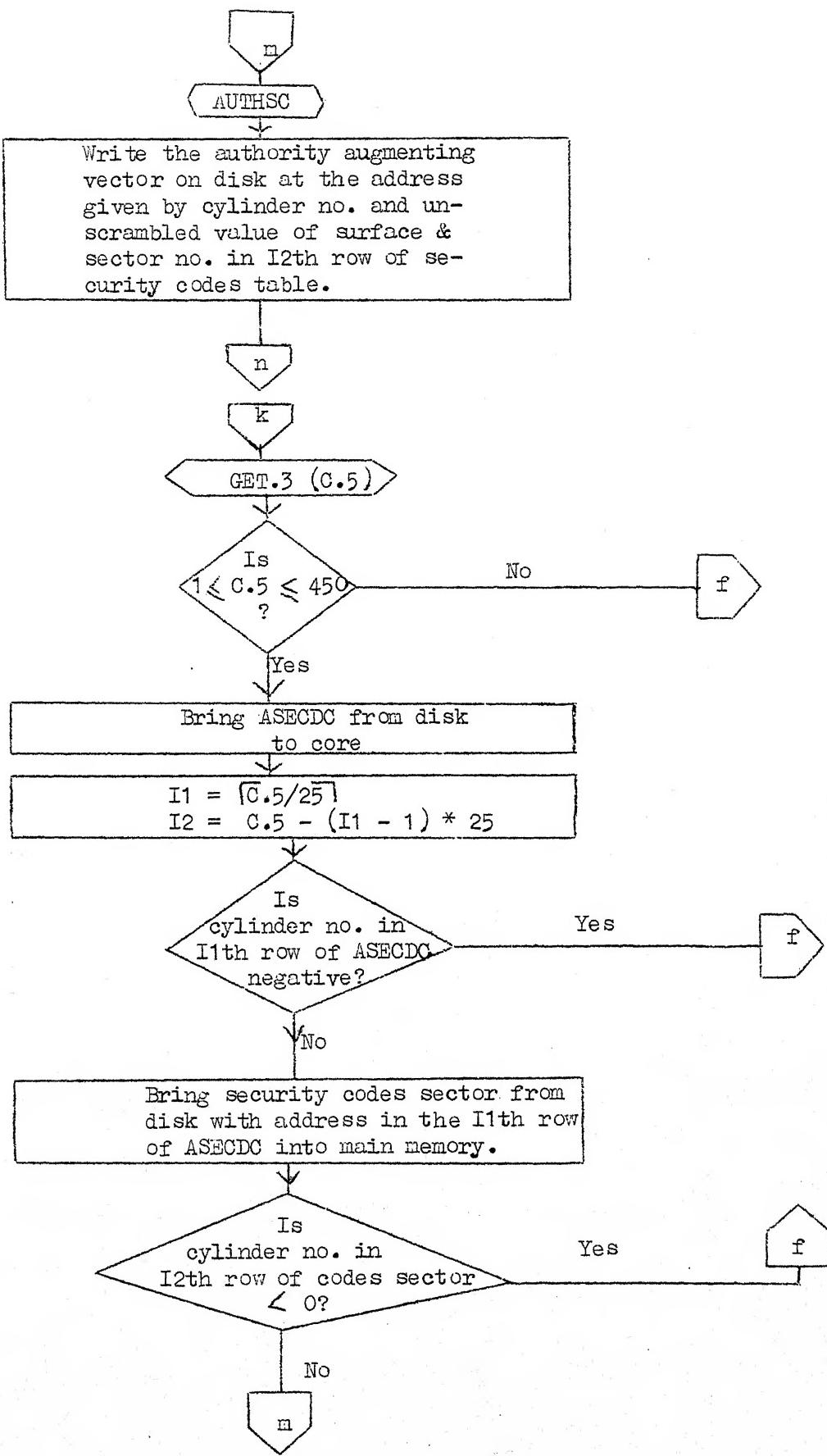


Figure 4-1: Flowchart for modifying the security codes' tables.

The modification algorithm in the form of a flowchart is shown in Figure 4-1. The program for this algorithm appears in the Program Listings appearing at the end of the thesis.

After modifying the security tables, the DBA can print the tables to satisfy himself that proper modification has been carried out. This he can do using the PRINTS routine which is given in the next chapter.

5. PRINTING THE SECURITY TABLES

ALGORITHM FOR PRINTING

The data base administrator may be interested in knowing the status of security tables at any time. The "PRINTS" routine gives the DBA the facility of printing FIRST Security codes' tabless and Authority Augmenting Security Codes Tables. First the DBA is asked to specify the output device where the tables are to be printed. The DBA gives "PR" for line printer and "TT" for teletype. Thereafter, system asks the DBA if he wants to print "First Security Codes' Table s". DBA gives "Y" "If", if he wants to print the tables. Otherwise he gives "N" "If". Similarly the DBA is asked if he want s to print "Authority Augmenting Security Codes' Table s". Again he gives "Y" "If". for printing the tables and "N" "If." for not printing the same.

In case of First Security Codes, first the system prints "First Security CodesTable Directory", i.e., FSECDC in the format of Figure 2-4. Then the system prints each existing first security code along with its status. Status can be any of the following -

- (i) A blank status field indicates that the code has been used properly.
- (ii) "ILLGLE" in status field indicates that the user tried to give wrong authority augmenting code to access the system illegally.
- (iii) "MORASK" in status field indicates that the user tried to ask the system for more than permitted to him, by giving such retrieval or update commands as were not permitted to him.

In case of Authority Augmenting Security Codes, the system first prints "Augmenting Security Codes Table Directory" i.e., ASECDC in the format of Figure 2-8. Thereafter, the system prints each existing authority augmenting code alongwith its associated authority augmenting vector. It also prints the sector address where this authority vector is stored on the disk. Since at any time, the data base shall generally contain much lesser number of fields then the maximum 256 permitted, the PRINTS routine asks the DBA to give a three digit input number (say N) specifying numberof fields in the system. Then it prints only the first N fields of each authority vector. Moreover, if the DBA does not want to print any authority vectors but only the authority augmenting security codes, he may give N=0.

The flowchart of the algorithm for printing security codes appears in Figure 5-1. The equivalent program of this algorithm appears in the program listings.

TERMINATING THE JOB

Before terminating DBA's job, we must write all the variable data as outlined in Flowchart 5-2 back on the disk. So at the end of any operation that the DBA may perform on the data base, he must perform the job "TERMNT". With this, all useful information in core gets stored back on the disk and may be fetched for later usage.

Having described all the programs that go into building the security system for the data base, the only function left is removing a relation from the data base which is the subject matter for the next chapter.

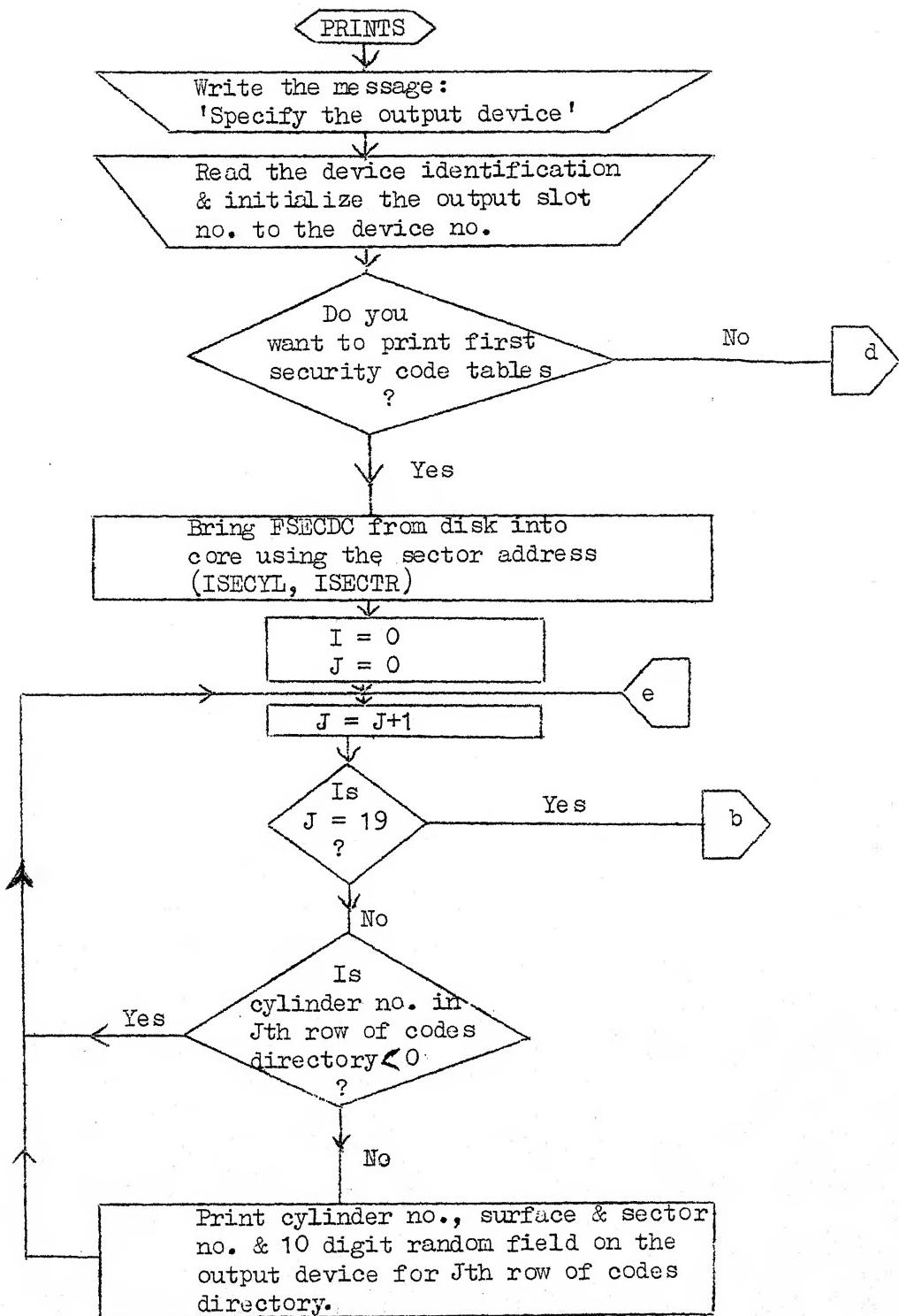


Figure 5-1: Continued on next page.

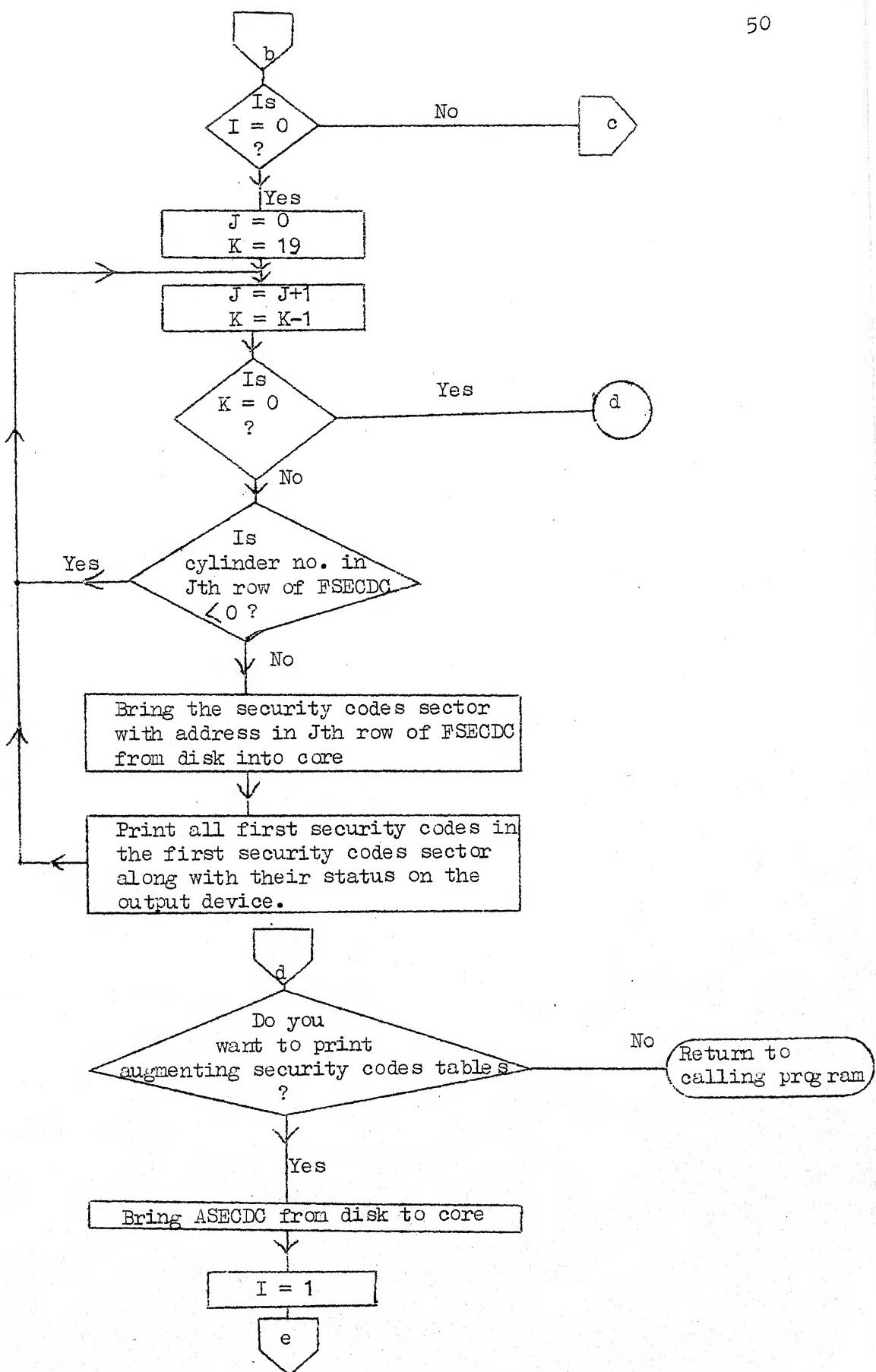


Figure 5-1: Continued on next page.

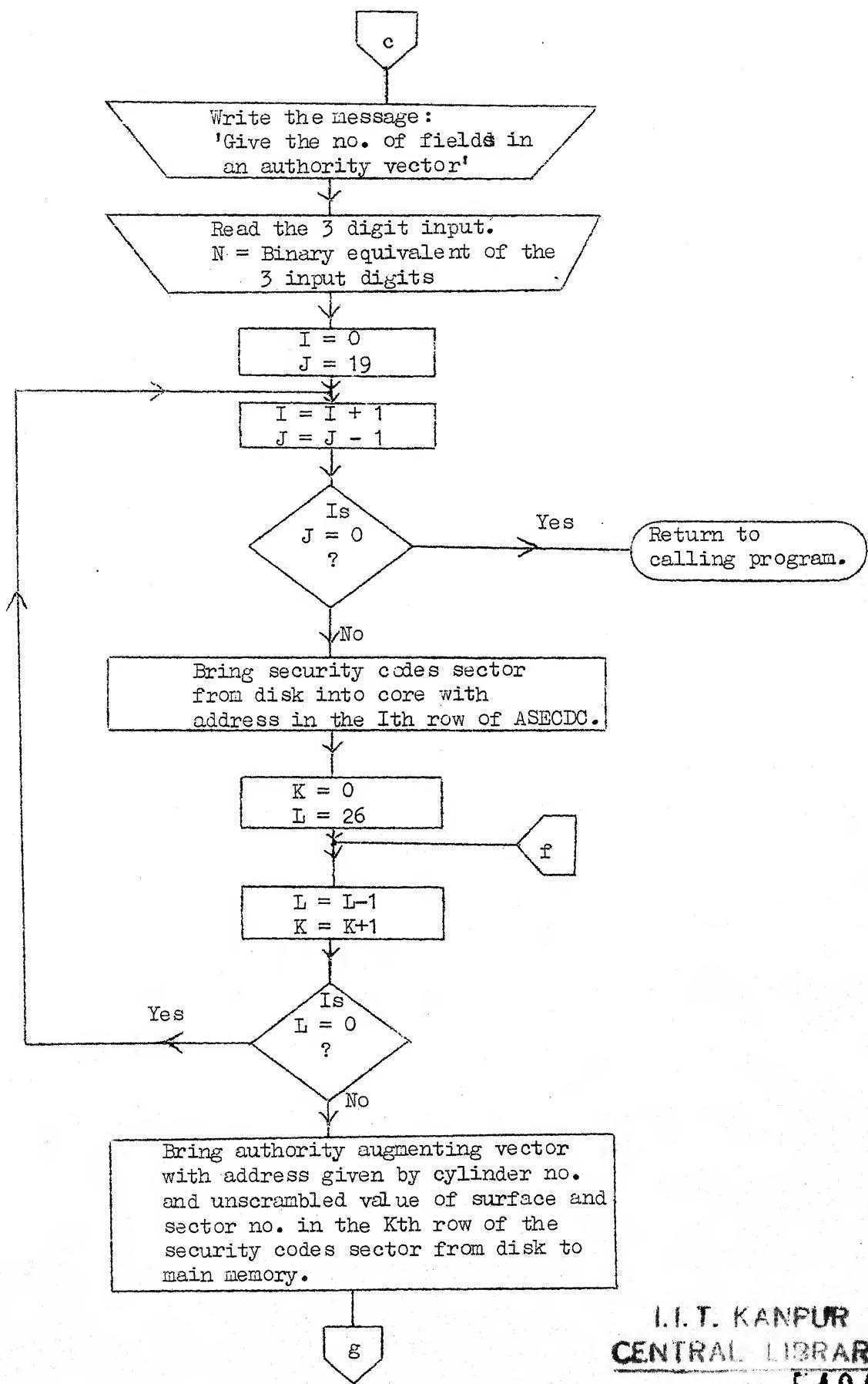


Figure 5-1: Continued on next page.

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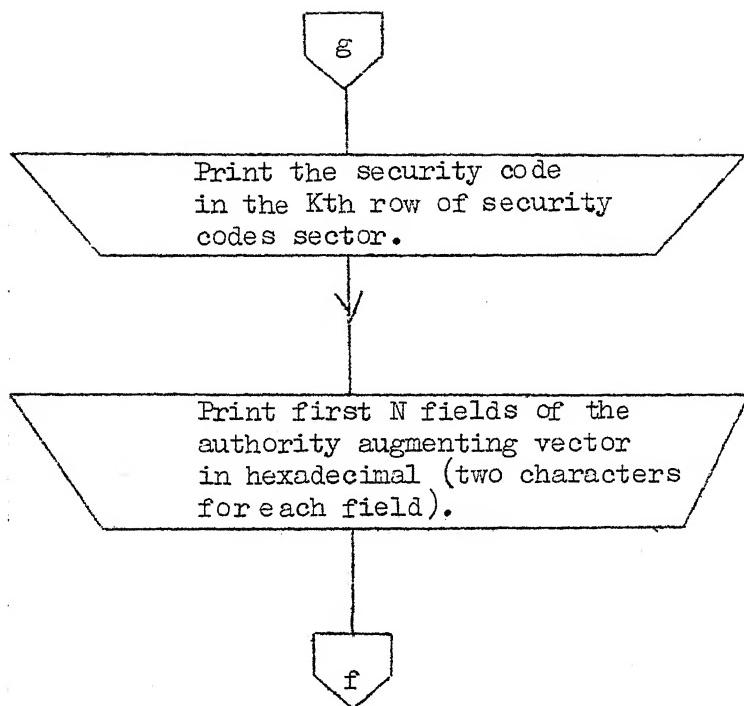


Figure 5-1: Flowchart for printing the security codes tables.

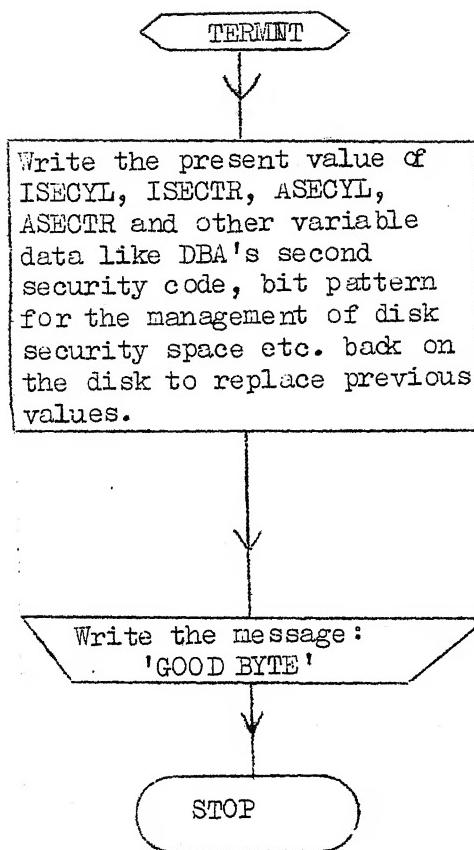


Figure 5-2: Flowchart for terminating DBA's job.

6. REMOVING A RELATION

As a data base grows older with time, it may become necessary to add more relation to the data base which may pertain to some data which has currently acquired importance, or to remove some relations which pertain to the data which has lost its relevance, or to re-organise the data base by combining some relations into one or breaking one relation into many or any combination of these in order to improve the efficiency of the system in the light of the past queries.

To illustrate by example let us say that we have a data base in IIT Kanpur where we store one relation for each year of students. Now as a new batch enters, a new relation shall have to be added to the system. But the relation corresponding to the outgoing batch may be dumped on some tape file (for any possible emergency need) and then removed from the data base.

Adding a relation is done by BUILDR whose programs appear in a separate thesis (Ref. 5). The present chapter discusses removing a relation. Since all the data structures used here appear in the thesis describing BUILDR programs, it is essential to go through it before studying this algorithm.

REMOVAL ALGORITHM

Removing of a relation from the data base involves the following tasks -

1. Removing this relation from the Relations Directory RELEDIR by clearing the RELID field for this relation, so that this relation identification number can be used to define new relation.

2. Removing the field records corresponding to this relation from the field list "FDLIST" and closing this gap by moving the field records lying below this gap upwards.
3. Setting the cylinders occupied by the relation on the disk free, by clearing their status bits in the bit vector (BITMAP).
4. Clearing the entries corresponding to this relation from the primary index table and moving the other entries up to close the gap.
5. For all existing relations whose primary index entries have been moved up, reflecting this movement in the relations directory RELEDIR by subtracting the length of the gap (i.e., length of upward movement) from their PMINX field values.
6. Writing these new tables values of RELEDIR, FDLIST, primary index etc. back on the disk.

All these steps are reflected in the flowchart of removal algorithm shown in Figure 6-1. The flowchart uses a routine RELCYL(N) which releases cylinder No. N of the disk and returns control ^{to} the calling program. Flowchart of this routine has been drawn in Figure 6-2. Programs corresponding to these two flowcharts appear in the Program Listings.

A FEW WORDS ON REORGANIZATION

Reorganization can be of two kinds:

1. Reorganizing the data base by merging two or more relation or breaking one relation into many or a combination of both to build new relations.
This can be done as follows:

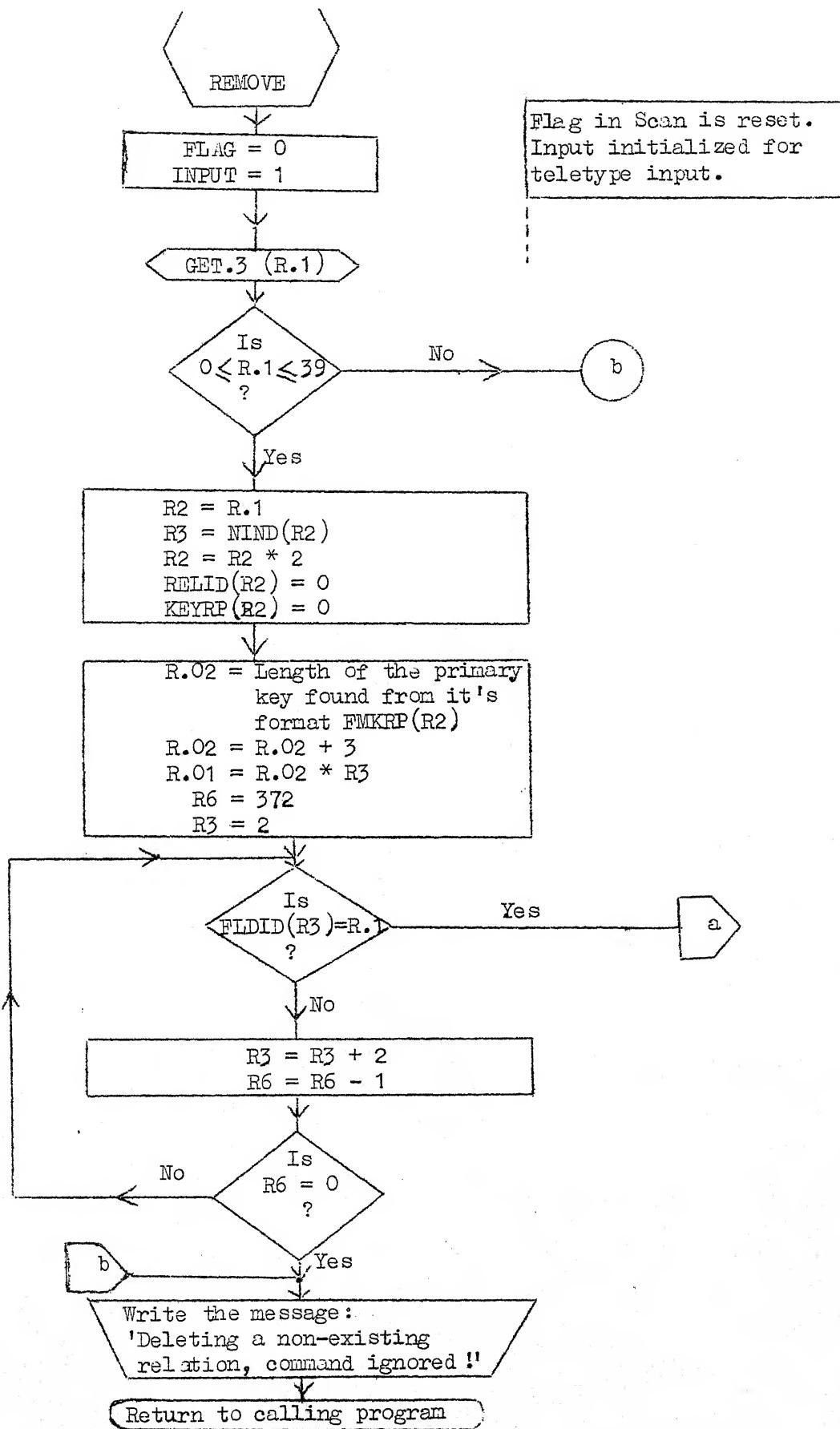


Figure 6-1: Continued on next page.

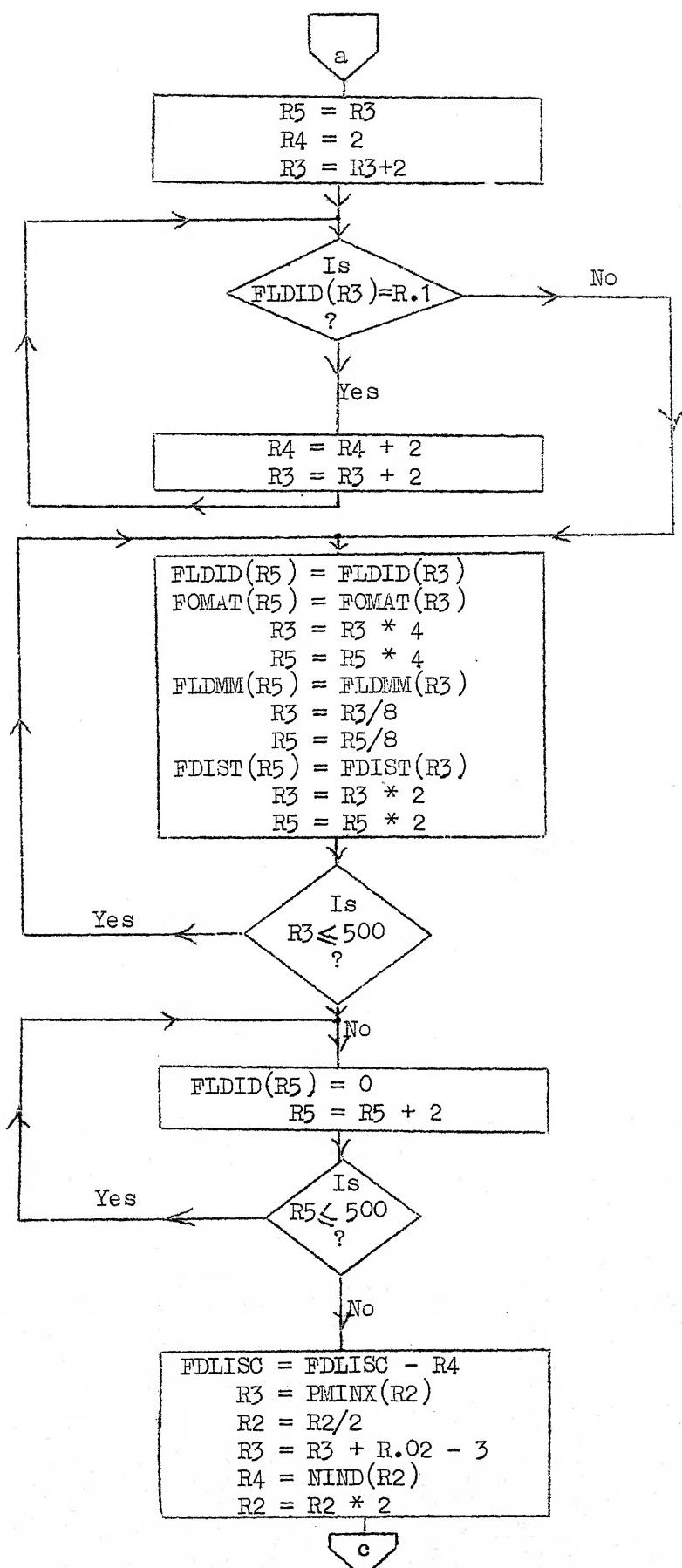


Figure 6-1: Continued on next page.

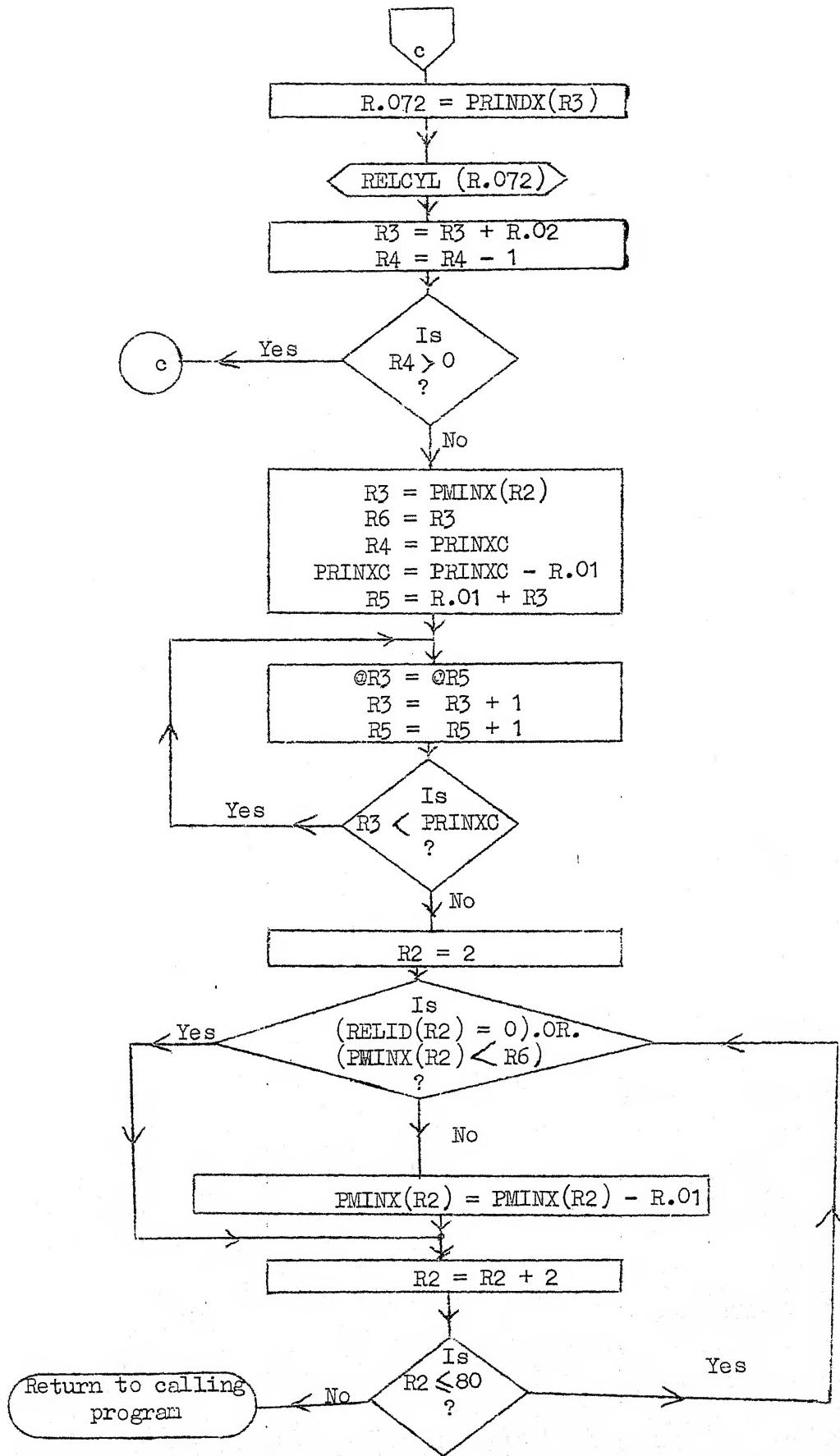


Figure 6-1: Flowchart for removing a relation from the data base.

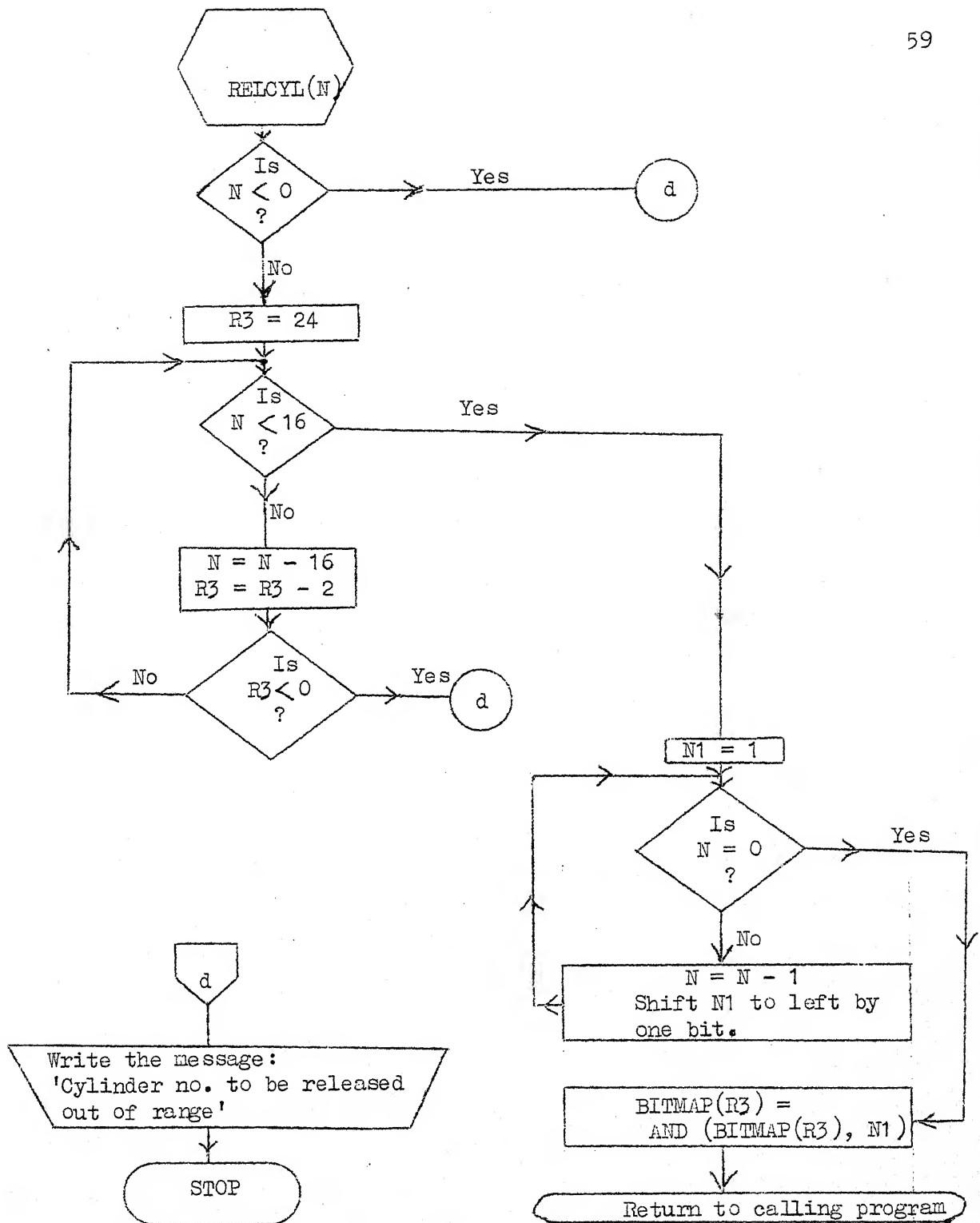


Figure 6-2: Subroutine for releasing a cylinder of the disk.

- (i) Retrieve the new relations to be created one by one by giving appropriate retrieval commands with output going on the disk (in a prefixed area).
 - (ii) Build these new relations by using BUILDR routine of (Ref. 5) by giving the specification of the new relations and reading the data from the disk.
 - (iii) Remove the older relations which have become redundant as a result of the creation of these new relations.
2. Reorganizing the data base by reorganising the relations such that all data from their overflow area is brought back into regular area and overflow area is cleared. This can be done for each relation to be reorganised as follows:

- (i) Retreieve the relation with output going to disk by giving appropriate retrieval command.
- (ii) Remove this routine from the data base using REMOVE routine.
- (iii) Build this relation using BUILDR routine by giving the specification of this relation and reading the data from the disk.

Hence using above steps the DBA can reorganise the data base in any manner he likes.

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1. Agarwala, S., "Updating of a Relational Data Base System on TDC-316", M.Tech. Thesis, Computer Science Programme, Indian Institute of Technology, Kanpur, July 1978.
2. Chamberlin, D.D., "Relational Data-Base Management Systems", ACM Computing Surveys, March 1976, Vol. 8, No. 1, pp. 43-66.
3. Codd, E.F., "A Relational Model of Data for Large Shared Data Banks", Comm. of ACM, Vol. 13, No. 6, June 1970, pp. 370-397.
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5. Ghonekwar, D.K., "An Implementation of a Relational Data Base Model", M.Tech. Thesis, Computer Science Program, Indian Institute of Technology, Kanpur, July 1977.
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8. Sibley, E.H., Fry, James, P., "Evaluation of Data-Base Management Systems", ACM Computing Survey, March 1976, Vol. 8, No. 1, pp. 7-72.
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Appendix 1

INITIAL LOADING OF THE SYSTEM

We have seen that in our system all the programs and security tables reside on the disk. Therefore to load the system initially i.e., to generate the system, the following steps must be followed.

1. Load the "initial loading program" from paper tape into core.

Execute it with starting address = 130000. This will put the routines WORBLE, DVORBL, VORBLE, GIVSEC on the preassigned sectors of the disk and will also initialise the security codes' tables so as to contain only one security code (that of DBA), with its value being 0291234512345 and also DBA's identification number being 000000000000 for values shown in the listings.

2. Load the paper tape for "program to load other programs" as shown in the listings into core.

3. Load paper tapes of all security programs into core except the START program (which is incorporated under KDM) with address ranges as below:

1. PRINTS & TERMNT	70000 to 76000
2. DBA, AIOCTS, RELSEC	76000 to 100240
3. Utility Routines	100400 to 102200
4. INTLIZE	102400 to 104000
5. MODIFY	104100 to 107500

4. Load the paper tapes of REMOVE and PUT routines into core with address range as below:

1. PUT	66000 to 66200
2. REMOVE	66400 to 67700

5. Load the programs developed for BUILDING of a relation developed as part of (Ref. 5) into core with address range as below:
 1. SCANNER 27600 to 46400
 2. STATIC Routine 1 46400 to 50430
 3. BUILD 50440 to 54340
 4. INDATA 54400 to 56230
 5. STATIC Routine 2 56300 to 57700
 6. MAIN 60000 to 65720
6. Execute the "program to load other programs" starting at address 20000. Execution of this program puts all the security programs and programs to build and remove a relation etc., from core into areas of the disk about which information is contained in "Initial loading program."
7. Load the retrieval programs developed as part of (Ref. 7) into core.
8. Load the paper tape of "program to load other programs" into core with X containing the value of symbol NR and Y containing the value of Symbol NR1 of Initial Loading Program.
9. Execute the "program to load other programs" again starting from 20000. This will load all retrieval programs also into preassigned areas of disk.
10. Repeat Steps 7, 8 and 9 for update programs developed as part of (Ref. 1) with X containing value of symbol NU and Y containing value of symbol NU1 of Initial Loading Programs. This will put update programs also into prefixed areas of the disk.

Now the START program incorporated under KDM with code "DB" shall call appropriate programs from the disk depending upon what the user wants to do, as detailed in Appendix 2.

Appendix 2
UNSERS' MANUAL

To use the system, the following steps must be followed:

Step 1: Press "DB". System prints a message asking the user to specify the first security code.

Step 2: User gives his 13-digit First Security Code as input from the teletype and then line feed. If user is DBA (i.e., the first security code given as input is that of the DBA) then go to Step 5 else go to next step.

Step 3: System asks the user to give his next security code (authority augmenting). User either gives 13-digit security code and then gives "Lf" or gives "N" "Lf". In the former case, the authority vector is enhanced corresponding to the authority augmenting vector associated with that security code and control goes back to Step 3. If input is "N" "Lf", go to next step.

Step 4: System asks the user to specify the job -

- (i) If user gives "R" "Lf" as input from the teletype, system loads Retrieval Programs plus all other utility routines from disk into core and transfers control to retrieval programs. Then user gives retrieval query as detailed in Users' Manual for Retrieval (Ref. 7). After completion of retrieval task control goes back to Step 4.

- (ii) If user gives "U" "Lf" as input from the teletype system loads update programs plus all other/utility routines from disk into core and transfers control to update programs. Then user gives update query as detailed in Users' Manual for Updating (Ref. 1). After completion of update task control goes back to Step 4.
- (iii) If user gives "F" "Lf" then job is finished and system gives "good-bye".
- (iv) If input is none of the above, error message is given and control goes back to Step 4.

Step 5: The DBA is asked to supply the identification number which is a 12-digit number. DBA gives 12-digit number followed by "Lf" or followed by "Y" "Lf". If identification number is correct, then in latter case, DBA is asked to supply new identification number which replaces the old identification number and control goes to next step. In former case control goes directly to next Step.

Step 6: DBA is asked to specify the job. Then DBA specifies one of the jobs as shown on Page 26 of Figure 2-11. If job is "TERMNT", control goes to Step 8 else control goes to Step 7.

Step 7: If job specified is "CCCCCC", then data base is cleared and control goes back to Step 6. If job specified is "INTLZE", "MODIFY" or "REMOVE", DBA follows the procedure mentioned in Chapters 3, 4 and 6 respectively. If job specified is "RDLBLD", DBA follows the instructions given in the Users' Manual of (Ref. 5). After completion of these tasks control goes back to Step 6. If job is "PRINTS" DBA follows steps given in Chapter 5, and control goes to next step.

Step 8: All tables from core are stored back on the disk. Job is finished and system gives "GOOD-BYE".

1 "PROGPAK TO LOAD OTHER PROGRAMS" IS USED TO LOAD SECURITY AND *
 1 OTHER PROGRAMS USED BY THE DRA INTO PRE ASSIGNED AREAS OF THE *
 1 DISK. TO LOAD RETRIEVAL PROGRAMS, REPLACE X=ND BY X=N1 & Y=N11 BY *
 1 Y=N11. TO LOAD UPDATE PROGRAMS, REPLACE X=ND BY X=N1 & Y=N11 BY *
 1 Y=N11. PROCEDURE FOR LOADING CAN BE SEEN FROM APPENDIX1 *

	20H00
R1=	X1
R2=	X2
R3=	X3
R4=	X4
N1=	26146
N11=	26150
N2=	26152
N3=	26154
N4=	26156
N5=	26160
X=	ND
Y=	N11
TSR	X,R2
TSR	Y,R3
L01:	TSR (R3)+,R4
	TSR (R3)+,STR8
	TSR (R3)+,CYL8
L02:	TSR (R3)+,SCT8
	JMS R4, WRITE
CYL8:	WORD 0
SCT8:	WORD 0
STR8:	WORD 0
	ADD #400,STR8
	DEC R4
	BRT L02
	DEC R2
	BRT L01
	STOP
WRITE:	TSR #200,#177456
	BR7S WRITE
	TSR (R4)+,#177452
	TSR (R4)+,#177444
	TSR (R4)+,#177460
	TSR #-128,#177462
	TSR #4113,#177454
	WAIT
SENSE:	TSR #1,#177456
	BR7C SENSE
	TSR #40000,#177456
	BR7S ERR,11
	RTS H4
ERR.11:	STOP
	END

INITIAL LOADING PROGRAM PUTS THE ROUTINES VORHLE, DVORG, GIVSEC *
 AND RD-FILE ON PREASSIGNED AREAS OF THE DISK. IT ALSO INITIALISES *
 VALUES OF VARIOUS VARIABLES AND TABLES AND PUTS THEM ON THE DISK *
 IT'S DATA PORTION CONTAINS ADDRESSES OF THE DISK WHERE PROGRAMS *
 NEEDED BY THE DIA, RETRIEVAL PROGRAMS AND UPDATE PROGRAMS ARE *
 STORED.

	130000
H1	X1
H2	X2
H3	X3
H4	X4
H5	X5
H6	X6
H7	X7
M0	177756
JMS	R4, WRITE
WORD	200., 293., 24400
JMS	R4, WRITE
WORD	200., 294., 25000
JMS	R4, WRITE
WORD	200., 295., 25400
JMS	R4, WRITE
WORD	200., 296., 26000
JMS	R4, WRITE
WORD	200., 297., 26400
JMS	R4, WRITE
WORD	200., 298., 27000
JMS	R4, WRITE
WORD	196., 1, 27400 .
JMS	R4, WRITE
WORD	196., 2, 30000
JMS	R4, WRITE
WORD	196., 3, 30400
STOP	
WHITE:	TSB #200, #177456
	BRZS WRITE
	TSR (R4)+, #177452
	TSR (R4)+, #177444
	TSR (R4)+, #177460
	TSR #-128., #177462
	TSR #4113, #177454
SENSE:	WAIT
	TSB #1, #177456
	BR ZC SENSE
	TSB #40000, #177456
	BR ZS ERR, 11
	RTS R4
ERR, 11:	STOP
	#24400
VORHLE:	TSB (R7)+, VVV, 1
	TSB VVV, 1, (R7)+

	CLR	(R7)+
	RTS	R7
VVV.1:	WORD	0
	LSL	24000
JVJH.LT:	TSR	(R7)+,DDD,1
	TSR	(R7)+,DDD,2
	TSR	DDD,1,(R7)+
	RTS	R7
LLD.1:	WORD	0
DDD.2:	WORD	0
	LSL	25000
GIVSR.L:	TSR	R5,-(R1)
	TSR	R4,-(R1)
	TSR	R3,-(R1)
	TSR	R2,-(R1)
	TSR	#2,R3
	TSR	(R7)+,R5
LG.0:	TSR	#5,R2
	CLR	M0
LG.1:	M0Y	#10..R4
	BTSH	(R5)+,R4
	ADD	R4,M0
	DEC	R2
	BRZC	LG.1
	TSR	M0,(R7)+
	DEC	R3
	BRZC	LG.0
	TSR	(R1)+,R2
	TSR	(R1)+,R3
	TSR	(R1)+,R4
	TSR	(R1)+,R5
	RTS	R7
LG.TT1:	WORD	0
	LSL	25400
WONHLE:	TSR	R4,-(R1)
	TSR	R3,-(R1)
	TSR	R2,-(R1)
	TSR	R5,-(R1)
	TSR	R6,-(R1)
	TSR	(R7)+,R5
	TSR	(R7)+,R6
	TSR	#10..R2
	CLR	R3
	CLR	R4
LPHOR:	BTSH	(R5)+,R3
	BTSH	(R6),R4
	ADD	R4,R3
	CMP	R3,#10,
	BRLT	LPH1
	SUB	#10..R3
LPH1:	BTSH	R3,(R6)+
	DEC	R2
	BRZC	LPHOR
	TSR	(R1)+,R6
	TSR	(R1)+,R9

TSR	(H1)+,H2
TSR	(H1)+,H3
TSR	(H1)+,H4
RTS	H/
:=	26000

; THE FOLLOWING SEGMENT CONTAINS INITIALISING DATA FOR VARIOUS *
; VARIABLES AND INITIAL IDENTIFICATION NO. OF DHA ETC. *

DISC1:	WORD	29.
ISECYL:	WORD	196.
ISLUTH:	WORD	1
ASBCYL:	WORD	196.
ASECT1:	WORD	2
FLAG1:	WORD	0
IONS:	WORD	14.,0,14.
IONU:	BYTE	60,60,60,60,60,60,60,60,60,60,60,60
ACYL:	WORD	200.,199.,198.,197.,196.
SITMH:	WORD	-1,0,-1
ND:	WORD	2
ND1:	WORD	26162
NR:	WORD	1
NR1:	WORD	26602
NU:	WORD	1
NUI:	WORD	27126

; THE FOLLOWING SEGMENT CONTAINS INFORMATION ABOUT THE ADDRESSES *
; WHERE PROGRAMS CALLED BY THE DDA GO ON THE DISK *

WORD	100.,27400,201.
WORD	1.,2.,3.,4.,5.,6.,7.,8.,9.,10.
WORD	33.,34.,35.,36.,37.,38.,39.,40.,41.,42.
WORD	65.,66.,67.,68.,69.,70.,71.,72.,73.,74.
WORD	97.,98.,99.,100.,101.,102.,103.,104.,105.,106.
WORD	129.,130.,131.,132.,133.,134.,135.,136.,137.,138.
WORD	161.,162.,163.,164.,165.,166.,167.,168.,169.,170.
WORD	193.,194.,195.,196.,197.,198.,199.,200.,201.,202.
WORD	225.,226.,227.,228.,229.,230.,231.,232.,233.,234.
WORD	257.,258.,259.,260.,261.,262.,263.,264.,265.,266.
WORD	289.,290.,291.,292.,293.,294.,295.,296.,297.,298.
WORD	10.,11400,202.
WORD	1.,2.,3.,4.,5.,6.,7.,8.,9.,10.
WORD	33.,34.,35.,36.,37.,38.,39.,40.,41.,42.
WORD	65.,66.,67.,68.,69.,70.,71.,72.,73.,74.

; THE FOLLOWING SEGMENT CONTAINS INFORMATION ABOUT ADDRESSES *
; WHERE RETRIEVAL PROGRAMS ARE STORED ON THE DISK *

RR1:	WORD	70.,54000,202.
WORD	97.,98.,99.,100.,101.,102.,103.,104.,105.,106.	
WORD	129.,130.,131.,132.,133.,134.,135.,136.,137.,138.	
WORD	161.,162.,163.,164.,165.,166.,167.,168.,169.,170.	
WORD	193.,194.,195.,196.,197.,198.,199.,200.,201.,202.	
WORD	225.,226.,227.,228.,229.,230.,231.,232.,233.,234.	

W0110 257., 258., 259., 260., 261., 262., 263., 264., 265., 266.
W0110 259., 260., 261., 262., 263., 264., 265., 266., 267., 268.

THE FOLLOWING SEGMENT CONTAINS INFORMATION ABOUT AREAS OF DISK WHERE UTILITY ROUTINES (SCANNER AND IT'S ROUTINES) AND OTHER TABLES, WHICH NEED TO BE LOADED WITH BOTH RETRIEVAL AND UPDATE PROGRAMS ARE STORED.

WURD	30., 27400, 201.
WOLD	1., 2., 3., 4., 5., 6., 7., 8., 9., 10.
WURD	33., 34., 35., 36., 37., 38., 39., 40., 41., 42.
WURD	65., 66., 67., 68., 69., 70., 71., 72., 73., 74.

THE FOLLOWING SEGMENT CONTAINS INFORMATION ABOUT AREAS OF DISK WHERE UPDATE PROGRAMS ARE STORED

UU: WORD 70..54000.195.
WORD 97..98..99..100..101..102..103..104..105..106.
WORD 129..130..131..132..133..134..135..136..137..138.
WORD 161..162..163..164..165..166..167..168..169..170.
WORD 193..194..195..196..197..198..199..200..201..202.
WORD 225..226..227..228..229..230..231..232..233..234.
WORD 257..258..259..260..261..262..263..264..265..266.
WORD 289..290..291..292..293..294..295..296..297..298.

INITIALISED FIRST SECURITY CODES: TABLE: DIRECTORY (FIND)

INITIALISED AUGMENTING SECURITY CODES' TABLE DIRECTORY (ASECDT)

```

      30000
WORD   -1,0,0,0,0,0,0
WORD   -1,0,0,0,0,0,0
WORD   -1,0,0,0,0,0,0
WORD   -1,0,0,0,0,0,0
WORD   -1,0,0,0,0,0,0

```

: INITIALISED SECTOR OF FIRST SECURITY CODES; TABLE RESULTING IN *

: SECURITY CODES 0271234512345 AND 0291234512345 *

HYTI	0,0,0,0,48,
MYTL	0,0,0,0,48,
MYTE	0,0,0,0,48,
YTH	0,0,0,0,48,
BYTE	0,0,0,0,48,
BYTH	0,0,0,0,48,
HYTE	0,0,0,0,48,
BYTE	0,0,0,0,48,
END	0,0,0,0,48,

START PROGRAM WHICH IS INCORPORATED UNDER KEM WITH CODE 'D'! *
 FOR STARTING EXECUTION FROM ADDRESS 20014 WHICH WILL PRINT *
 THE SECURITY CODES AS THEY ARE GIVEN FROM THE TELETYPE) AND *
 WITH CODE "RU" (FOR STARTING EXECUTION FROM ADDRESS 20000 WHICH *
 WILL SUPPRESS PRINTING OF SECURITY CODES AS THEY ARE INPUTTED *
 FROM THE TELETYPE) . *

L =	20000
K1=	0
TTY=	1
K1=	21
K2=	22
K3=	23
K4=	24
RJ=	25
K5=	26
R2=	27
M0=	177736
VORBLE=	24400
DVORBL=	24600
GIVSFC=	25000
WORBLE=	25400
DPSC.1=	26000
ISFCYL=	26002
ISFCTR=	26004
ASFCYL=	26006
ASECTR=	26010
TSR	#200,AA
TSR	#200,SIN,0+2

THE FOLLOWING PROGRAM SEGMENT READS ROUTINES VORBLE,DVORBL, *
 GIVSFC,WORBLE AND OTHER INITIAL INFORMATION FROM DISK INTO CORE. *

ST.1:	TSR	#137740,R1
	JMS	R4,READ
	WORD	200,,293,,24400
	JMS	R4,READ
	WORD	200,,294,,25000
	JMS	R4,READ
	WORD	200,,295,,25400
	JMS	R4,READ
	WORD	200,,296,,26000
	JMS	R4,READ
	WORD	200,,297,,26400
	JMS	R4,READ
	WORD	200,,298,,27000

THE FOLLOWING SEGMENT ASKS FOR FIRST SECURITY CODE INPUT. IF *
 FOUND VALID, IT JUMPS TO DBA'S ROUTINE IF SECURITY CODE IS THAT *
 OF THE DBA. OTHERWISE IT JUMPS TO THE SEGMENT WHICH ASKS FOR *
 AUTHORITY AUGMENTING SECURITY CODE INPUT. *

RF-SET		
	NEP	
	MEP	
	NEP	
	SWRITE	TTY, SME.21
W21:	WAITR	TTY, W21
	SQFAD	R40, SIN.0
S21:	WAITR	R40, SH2
	TSR	ISECYL, CYLO
	TSR	ISECTR, SCTO
	JMS	R4, READ
CYLO:	WORD	0
SCTO:	WORD	0
	0RD	SECSP1
	CLR	M0
	HTSR	SIN.0+6, M0
	SUB	#60, M0
	P	#10., R2
	HTSR	SIN.0+7, R2
	SUB	#60, R2
	ADD	R2, M0
	MPY	#10., R2
	HTSR	SIN.0+8., R2
	SUB	#60, R2
	ADD	M0, R2
	TSR	R2, R6
	TST	R6
	BREF	ST.4
	CMP	R6, #91B.
	BRGT	ST.4
	DEC	R2
	CLR	M0
ST.1.3:	CMP	R2, #51.
	BRET	ST.1.4
	INC	M0
	SUB	#51., R2
	BHN	ST.1.3
	MPY	#14., R3
	TSR	M0, SEC.01
	TSR	M0, R3
	TSR	#10., R5
	CLR	R4
ST.1.4:	HTSR	SIN.0+8.(R5), R4
	SUB	#60, R4
	HTSR	R4, SIN.0+8.(R5)
	DEC	R5
	BRZC	ST.1.1
	TST	SECSP1(R5)

	BHGF	ST.5.1
ST.5.2:	SWRITE	TTY,SME.23
	WAITR	TTY,W23
	SHEAD	KBD,SIN.0
SR3:	WAITR	KBD,SR3
	HCMR	SIN.0+6.#'N
	BR7C	ST.5.3
	CLR	SIN.0+2
	JMP	NDBALD
ST.5.3:	CLR	M0
	BTSR	SIN.0+6,M0
	SUB	#60,M0
	MPY	#10.,R2
	BTSR	SIN.0+7,R2
	SUR	#60,R2
	ADD	R2,M0
	MPY	#10.,R2
	BTSR	SIN.0+8..R2
	SUR	#60,R2
	ADD	M0,R2
	CLR	M0
	TST	R2
	BRIE	ST.9
	CMP	R2,#450.
	BRGT	ST.9
	DFC	R2
ST.5.6:	CMP	R2,#25.
	BRLT	ST.5.7
	INC	M0
	SUB	#25.,R2
	HRN	ST.5.6
ST.5.7:	MPY	#14.,R3
	TSR	M0,R3
	CLR	R4
	TSR	#10.,R5
ST.5.4:	BTSR	SIN.0+8.(R5),R4
	SUB	#60,R4
	BTSR	R4,SIN.0+8.(R5)
	DFC	R5
	BR7C	ST.5.4
	TST	SECSP1(R3)
	BHGF	ST.6
	JMP	ST.9
ST.6:	TSR	SECSP1(R3),CYL3
	TSR	SECSP1+2(R3),SCT3
	JMS	R4,READ
CYL3:	WORD	0
SCT3:	WORD	0
	WORD	SECSP2
	TSR	R2,M0
	MPY	#10.,R2
	TSR	M0,R2
	TST	SECSP2(R2)
	BHGF	ST.7
	JMP	ST.9

```

ST.7:    TSR      #SECSP1+4,SWOR05
          ADD     R3,SWOR05
          TSR      #SIN.0+9.,SWOR06
          JMS     R7,WORBLE
SWOR05: WORD    0
SWOR06: WORD    0
          TSR      SWOR06,SWOR07
          JMS     R7,GIVSEC
SWOR07: WORD    0
SWOR08: WORD    .+
          .+4
          BCMR    SWOR08,SECSP2+6(R2)
          BRZC    ST.9
          BCMR    SWOR08+1,SECSP2+7,(R2)
          BRZC    ST.9
          BCMR    SWOR08+2,SECSP2+8,(R2)
          BRZC    ST.9
          BCMR    SWOR08+3,SECSP2+9,(R2)
          BRZC    ST.9
          TSR      SECSP2(R2),CYL4
          TSR      SECSP2+2(R2),DVO.13
          TSR      SECSP2+4(R2),DVO.14
          JMS     R7,DVORBL
DVO.13: WORD    0
DVO.14: WORD    0
DVO.15: WORD    0
          TSR      DVO.15,SCT4
          JMS     R4,READ
CYL4:    WORD    0
SCT4:    WORD    0
          WORD    SECSP2
          TSR      #255.,R4
          BSTB    SECSP2(R4),SECSP3(R4)
          DFC     R4
          BRGF    ST.8.
          JMP     ST.5.2
***** THE FOLLOWING SEGMENT IS USED TO CANCEL FIRST SECURITY CODE OF ****
***** THE USER WHO ACTS IN AN UNAUTHORISED MANNER ****
*****
ST.9:    TSR      ISECYL,CYL5
          TSR      ISECTR,SCT5
          JMS     R4,READ
CYL5:    WORD    0
SCT5:    WORD    0
          WORD    SECSP1
          TSR      SEC.01,R3
          TSR      SECSP1(R3),CYL6
          TSR      SECSP1+2(R3),SCT6
          JMS     R4,READ
CYL6:    WORD    0
SCT6:    WORD    0
          WORD    SECSP2
          TSR      SEC.02,R2
          BTSR    ERBITS,SECSP2+4(R2)
          TSR     CYL6,CYL7

```

	TSR	SCT6,SCT7
	JMS	R4,WRITE
CYL7:	WORD	0
SCT7:	WORD	0
	WORD	SECSP2
W22.L1:	SWRITE	TTY,SME,22
	WAITR	TTY,W22.1
	STOP	
AA:	WORD	0
DB:	BCI	%DB%
STN.0:	WORD	16.,0.16.
		.+16.
SME.21:	WORD	40.,0.40.
	BYTF	15,12
	BCI	% PLEASE GIVE YOUR FIRST SECURITY
	BCI	%TY CODE %
SME.22:	WORD	28.,0.28.
	BYTF	15,12
	BCI	% SORRY,UNABLE TO SERVE YOU%
SMF.23:	WORD	60.,0,60.
	BYTE	15,12
	BCI	% PLEASE GIVE YOUR NEXT SECURITY %
	BCI	%CODE(AUTHORITY AUGMENTING)%
SEC.01:	WORD	0
SEC.02:	WORD	0
BYTE1:	BYTF	0,0

* THE FOLLOWING SEGMENT CONTAINS ROUTINES FOR READING & WRITING *
* ON THE DISK *

READ:	TSR	#4107,-(R1)
	BRN	RW
WRITE:	TSR	#4113,-(R1)
RW:	TSB	#200,0#177456
	BR7S	RW
	TSR	(R4),0#177452
	CMP	(R4)+,#195.
	BR1T	ERR,12
	TSR	(R4)+,0#177444
	TSR	(R4)+,0#177460
	TSR	#-128,,0#177462
	TSR	(R1)+,0#177454
	WAIT	
SENSE:	TSB	#1,0#177456
	BR7C	SENSE
	TSB	#40000,0#177456
	BR7S	ERR,11
	RTS	R4
ERR.12:	STOP	.
ERBITS:	WORD	12.
ERR.11:	STOP	
SECSP1:	,*	,+256,
SECSP2:	,*	,+256,
SECSP3:	,*	,+256.
	STDRA#	75762

MDRAST=	54000
R7=	26146
RD1=	26150
RD2=	26152
RD1=	26154
R11=	26156
NU1=	26160

***** THIS SEGMENT LOADS SECURITY, BUILD AND REMOVE PROGRAM
NEEDED BY THE DBA AND THEN JUMPS TO DBA'S ROUTINE

***** S ETC. *

TRALD:	TSR	NC,R2
	TSR	(R1),R3
LD1:	TSR	(R3)+,R4
	TSR	(R3)+,STR8
LC2:	TSR	(R3)+,CYL8
	TSR	(R3)+,SCT8
	JMS	R4,READ
CYL6:	WORD	0
SCT8:	WORD	0
STR8:	WORD	0
	ADD	#400,STR8
	DEC	R4
	BRGT	LD2
	DEC	R2
	BRGT	LD1
	JMP	STDRA

THIS SEGMENT LOADS PROGRAMS FOR RETRIEVAL OR UPDATING
UPON USER'S INTEREST & JUMPS TO THE CORRESPONDING

***** DEPENDING *
UTINE *

NDBALD:	SWRITE	TTY,SME,30
WNR1:	WAITR	TTY,WNR1
	SREAD	KBD,SIN,0
WNR2:	WAITR	KBD,WNR2
	RCMP	SIN,0+6,#'R
	RR7S	RETRVL
	RCMP	SIN,0+6,#'U
	RR7S	UPDATE
	RCMP	SIN,0+6,#'F
	RR7C	NDBALD
	SWRITE	TTY,SME,31
WNR3:	WAITR	TTY,WNR3
	STOP	
	NOP	

RETRVL:	TSR	NR,R2
	TSR	NR1,R3
	JMP	LD3
UPDATE:	TSR	NU,R2
	TSR	NU1,R3
	SUB	#66,,R3
LD3:	INC	R2
LD3.1:	TSR	(R3)+,R4
	TSR	(R3)+,STR9
	TSR	(R3)+,CYL9

```

LD4:    TSH      (R3)+,SCT9
        JMS      R4,READ
CYC9:    WORD     0
SCT9:    WORD     0
STR9:    WORD     0
        ADD      #400,STR9
        DEC      R4
        BRGT    LD4
        DEC      R2
        BHGT    LD3.i
        JMP      NDBAST
SME.30:  WORD     18.,0.18.
        BYTE    15,12
        BCI      % SPECIFY THE JOB#
SME.31:  WORD     12.,0.12.
        BYTE    15,12
        BCI      % GOOD BYE %

```

***** THE FOLLOWING ROUTINE IS USED TO CHECK WHETHER USE
 FALLS WITHIN HIS AUTHORITY. IF NOT, HIS FIRST SECU
 CANCELLED

S REQUEST %
 TY CODE IS %

```

CHKSEC:  TSR      R3,-(R1)
        TSH      R2,-(R1)
        CLR      R2
        CLR      R3
SL.1:    BC1B    SECSP3(R2),SECSP1(R2)
        BTSR    SECSP1(R2),R3
        TST      R3
        BRZC    SL.2
        INC      R2
        CMP      R2,#256
        BRI T   SL.1
        TSR      (R1)+,R2
        TSR      (R1)+,R3
        RTS      R7
SL.2:    TSR      #3,ERBITS
        JMP      ST.9
        END

```

; PRINTS PROGRAM WHICH IS USED TO PRINT ALL SECURITY TABLES & ALSO
; AUTHORITY VECTORS AS REQUESTED BY THE GUA

	1=	70000
	R1=	X1
	R2=	X2
	R3=	X3
	R4=	X4
	R5=	X5
	R6=	X6
	R7=	X7
	M0=	177736
	KBD=	0
	TTY=	1
	PN=	6
	ISFCYL=	26002
	ISFCTR=	26004
	ASFCYL=	26006
	ASECTR=	26010
	DVORHL=	24600
	SFCSP1=	22170
	SFCSP2=	22570
	SFCSP3=	23170
	READ=	022060
	WRITE=	022066
	EXIT=	76214
PRINTS:	TSR	R7,-(R1)
	TSR	R6,-(R1)
	TSR	R5,-(R1)
	TSR	R4,-(R1)
	TSR	R3,-(R1)
	TSR	R2,-(R1)
	TSR	R1,SSAVEP
	RESET	
S.71:	SWRITE	TTY,SME,14
W14:	WAITR	TTY,W14
	SREAD	KBD,SIN,4
WRP1:	WAITR	KBD,WRP1
	CMP	SIN,4+6,#"PN
	BRZC	S.72
	TSR	PRNTR,INPUTP
	BRN	S.73
S.72:	CMP	SIN,4+6,#"TT
	BRZC	S.71
	TSR	TELTYR,INPUTP
	TSR	#5015,TBF1
	TSR	#5015,TBF2
	TSR	#5015,TBF3
	TSR	#5015,TBF4
	TSR	#5015,TBF5
	TSR	#5015,TBF6
	TSR	#5015,TBF7

	TSR	#5015.T0FB
	TSR	#5015.T0F9
	TSR	#5015.T0F10
S.73:	INIT	PRN,INPUTP
	SWRITE	TTY,SME.15
	WAITR	TTY,W15
	SRHAD	KBD,SIN.4
	WP2:	WAITR,KBD,WRP2
	BCMP	SIN.4+6,W'Y
	BRZS	S.74
	JMP	S.HB
S.74:	TSR	ISECYL.CYL.41
	TSR	ISECTR.SCT.41
	JMS	R4,READ
CYL.41:	WORD	0
SCT.41:	WORD	0
	WORD	SECSP1
	CIN	R6
	SWRITE	PRN,PHF1
WP1:	WAITR	PRN,WP1
S.P.11:	SWRITE	PRN,PHF10
WP2:	WAITR	PRN,WP2
WP3:	SWRITE	PRN,PHF2
WP3:	WAITR	PRN,WP3
WP4:	SWRITE	PRN,PHF10
	WAITR	PRN,WP4
	TSR	#14..R2
	CLR	PCOUNT
	CLR	R3
S.P.11:	TSR	#20040,PHF3+36,(R3)
	ADD	#2,R3
	CMP	R3,#70.
	BRZC	S.P.11
	INC	PCOUNT
	CMP	PCOUNT,#19.
	BRZS	S.76
	ADD	#14..R2
	TST	SECSP1(R2)
	BRLT	S.75
	JMS	R7,POWR
	WORD	2
	WORD	PCOUNT
	WORD	PHF3+47.
	TSR	SECSP1(R2),PTEMP
	JMS	R7,POWR
	WORD	3
	WORD	PTEMP
	WORD	PHF3+58.
	TSR	SECSP1+2(R2),PTEMP
	JMS	R7,POWROC
	WORD	PTEMP
	WORD	PHF3+70.
	TSR	#SECSP1+4,PTEMP
	ADD	R2,PTEMP
	JMS	R7,CDPC

	WORD	10.
	WORD	PTEMP
	WORD	PHF3+83.
	SWRITE	PRN, PBF3
WP5:	WAITR	PRN, WP5
	BRN	S.75
	TST	R6
	BR7S	S.P.2
	JMP	S.90
S.P.2:	TSR	#141..R2
	TSR	#19..R3
	SWRITE	PRN, PBF4
WP6:	WAITR	PRN, WP6
	SWRITE	PRN, PBF10
WP7:	WAITR	PRN, WP7
	SWRITE	PRN, PBF5
WP8:	WAITR	PRN, WP8
	SWRITE	PRN, PBF10
WP9:	WAITR	PRN, WP9
	CLR	R4
S.P.12:	TSR	#20040, PBF3+36..(R4)
	ADD	#2,R4
	CMP	R4,#70.
	BRZC	S.P.12
	CLR	R5
S.77:	ADD	#14..R2
	DEC	R3
	BRZC	S.P.15
	JMP	S.89
S.P.15:	TST	SECSP1(R2)
	BRLT	S.77
	TSR	SECSP1(R2), CYL.42
	TSR	SECSP1+2(R2), SCT.42
	JMS	R4, READ
CYL.42:	WORD	0
SCT.42:	WORD	0
	WORD	SECSP2
	TSR	#18..PCOUNT
	SUH	R3, PCOUNT
	TSR	PCOUNT, H0
	MPY	#51..PCOUNT
	TSR	H0, PCOUNT
	TSR	#-5,R4
S.7d:	INC	PCOUNT
	ADD	#9,R4
	CMP	R4,#250.
	BRGT	S.77
	BTSR	SECSP2+4(R4), BYT4
	BCLB	BYT1..BYT4
	BCMP	BYT4,#48.
	BR7S	S.78
	BTSR	SECSP2+4(R4), BYT4
	BCLB	BYT2..BYT4
	BCMP	BYT4,#12.
	BRZC	S.79

	TSR	ILLGLF+PBF3+46,(R5)
	TSR	ILLGLF+2,PBF3+48,(R5)
	TSR	ILLGLF+4,PBF3+50,(R5)
	BRN	S.P.13
S.71:	BTSH	SECSP2+4(R2),BYT4
	DC1H	BYT3,BYT4
	BCNP	BYT4,R3
	BRZC	S.80
	TSR	MURASK,PBF3+46,(R5)
	TSR	MURASK+2,PBF3+48,(R5)
	TSR	MURASK+4,PBF3+50,(R5)
	BRN	S.P.13
S.80:	TSR	#20040,PBF3+46,(R5)
	TSR	#20040,PBF3+48,(R5)
	TSR	#20040,PBF3+50,(R5)
S.P.13:	BTSH	SECSP2(R4),PTEMP
	BTSH	SECSP2+1(R4),PTEMP+1
	JMS	R7,POWR
	WORD	3
	WORD	PTCMP
S.81:	WORD	PBF3+57
	BTSH	SECSP2+2(R4),PTEMP
	BTSH	SECSP2+3(R4),PTEMP+1
	JMS	R7,POWR
	WORD	5
	WORD	PTEMP
S.82:	WORD	PBF3+62
	JMS	R7,POWR
	WORD	3
	WORD	PCOUNT
S.83:	WORD	PBF3+54
	INC	R6
	CMP	R6,#1
	BRZC	S.84
	ADD	#29..S.81
	ADD	#29..S.82
	ADD	#29..S.83
	TSR	#28..R5
	BRN	S.78
S.84:	SUB	#29..S.81
	SUB	#29..S.82
	SUB	#29..S.83
	CLR	R5
	CLR	R6
	SWRITE	PRN,PBF3
WP10:	WAITR	PRN,WP10
	BRN	S.78
S.85:	TST	R6
	BRZS	S.88
	TSR	#25..R2
S.86:	DEC	R2
	BRZS	S.87
	BTSH	#40,PBF3+73,(R2)
	BRN	S.86
S.87:	SWRITE	PRN,PBF3

WP11:	WAITR	PRN,WP11
S.89:	SWRITE	TTY,SME.16
W16:	WAITR	TTY,W16
	SREAD	KBD,SIN.4
WRP3:	WAITR	KBD,WRP3
	SCMP	SIN.4+6,#'Y
	BZS	S.89
	JMP	ENDPRI
S.89:	TSR	#1,R6
	TSR	ASECYL,CYL.43
	TSR	ASECTH,SCT.43
	JMS	H4,HEAD
CYL.43:	WORD	0
SCT.43:	WORD	0
	WORD	SRCSPI
	SWRITE	PRN,PBF6
WP12:	WAITR	PRN,WP12
	JMP	S.P.1
S.90:	SWRITE	TTY,SME.17
W17:	WAITR	TTY,W17
	SREAD	KBD,SIN.4
WRP4:	WAITR	KBD,WRP4
	CLR	R6
	RTSR	SIN.4+6,R6
	SUB	#60,R6
	TSR	R6,MD
	MPY	#10.,R6
	RTSR	SIN.4+7,R6
	SUB	#60,R6
	ADD	R6,MD
	MPY	#10.,R6
	RTSR	SIN.4+8.,R6
	SUB	#60,R6
	ADD	MD,R6
	TST	R6
	BRLT	S.90
	CMP	R6,#256.
	BNLE	S.91
	TSR	#256.,R6
S.91:	SWRITE	PRN,PBF7
WP121:	WAITR	PRN,WP121
	SWRITE	PRN,PBF10
WP13:	WAITR	PRN,WP13
	TSR	#-14.,R2
	TSR	#19.,R3
	CLR	R4
S.P.14:	TSR	#20040,PBF3+36.(R4)
	ADD	#2,R4
	CMP	R4,#70.
	BZC	S.P.14
S.92:	ADD	#14.,R2
	DEC	R3
	BZS	ENDPRI
	TST	SECSP1(R2)
	BRLT	S.92

	TSR	SECSP1(H2), CYL.44
	TSR	SECSP1+2(H2), SCT.44
	JMS	R4, READ
CYL.44:	WORD	0
SCT.44:	WORD	0
	WORD	SECSP2
	TSR	#18., PCOUNT
	SUB	R3, PCOUNT
	TSR	PCOUNT, MU
	MPS	#25., PCOUNT
	TSR	MU, PCOUNT
	TSR	#-10., R4
	TSR	#26., R2
S.93:	DEC	R5
	SH2S	S.92
	ADD	#10., R4
	INC	PCOUNT
	TST	SECSP2(R4)
	BAL T	S.93
	TSR	SECSP2(R4), CYL.45
	TSR	SECSP2+2(R4), DVO.10
	TSR	SECSP2+4(R4), DVO.11
	JMS	R7, DVORBL
DVO.10:	WORD	0
DVO.11:	WORD	0
DVO.12:	WORD	0
	TSR	DVO.12, SCT.45
	JMS	R4, READ
CYL.45:	WORD	0
SCT.45:	WORD	0
	WORD	SECSP3
	JMS	R7, PWR
	WORD	3
	WORD	CYL.45
	WORD	PBF8+79.
	JMS	R7, PWRROC
	WORD	SCT.45
	WORD	PBF8+94.
	TSR	SECSP2+6(R4), PTEMP
	JMS	R7, PWR
	WORD	5
	WORD	PTEMP
	WORD	PBF8+55.
	TSR	SECSP2+8.(R4), PTEMP
	JMS	R7, PWR
	WORD	5
	WORD	PTEMP
	WORD	PBF8+60.
	TSR	PCOUNT, PTEMP
	JMS	R7, PWR
	WORD	3
	WORD	PTEMP
	WORD	PBF8+52.
	SWRITE	PRN, PBF8
WP14:	WAITR	PRN, WP14

	SWRITE	PRN, PBFF
WP12:	WORD	HNN, WP15
	SWRITI	PRN, PBFI0
WP14:	WAITR	PRN, WP16
	TSR	R6, NUMBER
	JMS	R7, PRISec
AUG IPR:	WORD	0
	WORD	SECUP3
	WORD	PBF 3+41.
	JMP	S, 93
ENDPRN1:	TSR	SSAVEP, R1
	TSR	(R1)+, R2
	TSR	(R1)+, R3
	TSR	(R1)+, R4
	TSR	(R1)+, R5
	TSR	(R1)+, R6
	TSR	(R1)+, R7
	JMP	TERMNT
PRF1:	WORD	100..0.100.
	BCI	%
TdF1:	BCI	% FIRST %
	BCI	% SECURITY CODES' TABLE %
	BCI	% DIRECTORY %
PRF2:	WORD	100..0.100.
	BCI	%
TBF2:	BCI	% SING CYLINDER %
	BCI	% NO. SECTOR NO. ROTATING %
	BCI	% DIGITS %
PRF3:	WORD	100..0.100.
	BCI	%
TdF3:	BCI	%
	BCI	%
	BCI	%
PRF4:	WORD	100..0.100.
	BCI	%
	BYTE	15, 12, 15, 12
TBF4:	BCI	%
	BCI	% FIRST SECURITY CODES' %
	BCI	% TABLE %
PRF5:	WORD	100..0.100.
	BCI	%
TBF5:	BCI	% STATUS SECURITY %
	BCI	% CODE STATUS %
	BCI	% SECURITY CODE %
PRF6:	WORD	100..0.100.
	BCI	%
	BYTE	15, 12, 15, 12
TBF6:	BCI	% AUGMENTING %
	BCI	% SECURITY CODES' TABLE %
	BCI	% DIRECTORY %
PRF7:	WORD	100..0.100.
	BCI	%
	BYTE	15, 12, 15, 12
TBF7:	BCI	% AUGMENTING SECURITY %
	BCI	% CODES' TABLE WITH AUTHORITY %
	BCI	% VECTORS %

PBF8%	WORD	loc., o, loc-
	BCI	8
TBF%:	BCI	% SECURITY CODES
	BCI	% CYLINDER NO.= %
	BCI	% SECTOR NO.= %
PBF4%:	WORD	100,,0,100.
	BCI	2
TBF%:	BCI	% CORRESPONDING AUTHORITY %
	BCI	% VECTOR IS AS BELOW:
	BCI	2
MHF10%:	WORD	2,0,2
MHF10%:	WORD	2040
SAVER%:	WORD	0
SIN.4%:	WORD	14,,0,14.
	BCI	.+14.
PRNTR%:	WORD	10
INPUT%:	WORD	0
TELTYPE%:	WORD	2
PTEMP%:	WORD	0
PCOUNT%:	WORD	0
BYT1%:	BYTE	-49.
BYT2%:	BYTE	-13.
BYT3%:	BYTE	-4
BYT4%:	BYTE	0.
ILLOLR%:	BCI	%ILLGLE%
MORASK%:	BCI	%MONASK%
SME.14%:	WORD	64,,0,64.
	BYTE	15,12
	BCI	% SPECIFY THE DEVICE WHERE %
	BCI	% SECURITY TABLES ARE TO BE PRINTED %
SME.15%:	WORD	56,,0,56.
	BYTE	15,12
	BCI	% DO YOU WANT TO PRINT %
	BCI	% FIRST SECURITY CODES' TABLES? %
SME.16%:	WORD	60,,0,60.
	BYTE	15,12
	BCI	% DO YOU WANT TO PRINT AUG %
	BCI	% MENTING SECURITY CODES' TABLES? %
SME.17%:	WORD	54,,0,54.
	BYTE	15,12
	BCI	% GIVE NO. OF FIELDS IN AUTHORITY%
	BCI	% VECTORS IN 3 DIGITS%
<hr/> THIS ROUTINE PRINTS A SECTOR OF AUTHORITY VECTOR & IS CALLED BY %		
THE PRINTS ROUTINE .		
PRISEC%:	TSR	R5,-(R1)
	TSR	R4,-(R1)
	TSR	R3,-(R1)
	TSR	R2,-(R1)
	TSR	(R7)+.R2
	TSR	(R7)+.R3
	CLR	
	TSR	PDL1
	TSR	(R7)+.STPB
PHL0%:	TSR	#20,.R4
	TSR	STPB,R5

PRL1:	BTSH	(R5)+,BYT5
	BTSH	BYT5,BYT6
	BCLR	BYT7,BYT5
	BCLB	BYT8,BYT6
	BCMP	BYT5,#9.
	BRCT	PRL2
	BTSH	BYT5,PDL1
	ADD	#60,PDL1
	BRN	PRL3
PRL2:	BTSH	BYT5,PDL1
	ADD	#67,PDL1
PRL3:	BCLR	BYT5
	LSL	BYT6
	LSL	BYT6
	LSL	BYT6
	BTSH	PDL1,BYT6
	BCMP	BYT5,#9.
	BRCT	PRL4
	BTSH	BYT5,PDL1
	ADD	#60,PDL1
	BRN	PRL5
PRL4:	BTSH	BYT5,PDL1
	ADD	#67,PDL1
PRL5:	BTSH	PDL1,(R5)+
	BTSH	BYT6,(R5)+
	BTSH	#40,(R5)+
	DEC	R2
	BRLE	PRL6
	DEC	R4
	BRZC	PRL1
	SWRITE	PRN,PBF3
WP17:	WAITR	PRN,WP17
	BRN	PRL0
PRL6:	DEC	R4
	TSH	R4,M0
	MPY	#3,R4
	TSR	M0,R4
PRL7:	BTSH	#40,(R5)+
	DEC	R4
	BRGT	PRL7
	SWRITE	PRN,PBF3
WP18:	WAITR	PRN,WP18
	TSR	(R1)+,R2
	TSR	(R1)+,R3
	TSR	(R1)+,R4
	TSR	(R1)+,R5
	RTS	R7
PDL1:	WORD	0
BYT8:	BYTE	15.
BYT7:	BYTE	-16.
BYT6:	BYTE	0
BYT5:	BYTE	0
	STPB:	WORD 0

; COPC, POWROC & POWR ARE A FEW UTILITY ROUTINES CALLED BY PRINTS &
; PROGRAM & ARE SELF EXPLANATORY.

```
COPC:    TSR      R5,-(R1)
         TSR      R4,-(R1)
         TSR      R3,-(R1)
         TSR      R2,-(R1)
         CLR      R5
         TSR      (R7)+,R2
         TSR      #0(R7)+,R3
         TSR      (R7)+,R4
LCP1:    BTSH     (H3)+,R5
         ADD      #60,R5
         BTSH     R5,(R4)-
         DRC      R2
         BRZC    LCP1
         TSR      (R1)+,R2
         TSR      (R1)+,R3
         TSR      (R1)+,R4
         TSR      (R1)+,R5
         RTS      R7
         TSR      R5,-(R1)
         TSR      R4,-(R1)
         TSR      R3,-(R1)
         TSR      R2,-(R1)
         TSR      #0(R7)+,R2
         TSR      (R7)+,R3
         TSR      R2,R4
         CLB      PW1,R2
         CLB      PW2,R4
         BTSH     #60,(R3)-
         ASR      R2
         ASR      R2
         ASR      R2
         ASR      R2
         ASR      R2
         ADD      #60,R2
         BTSH     R2,(R3)-
         BTSH     #1,(R3)-
         CMP      R4,#10.
         BRZS    PP1
         ADD      #60,R4
         BTSH     #60,(R3)-
         BTSH     R4,(R3)-
         BRN      PP2
PP1:    BTSH     #61,(R3)-
         BTSH     #60,(R3)-
PP2:    TSR      (R1)+,R2
         TSR      (R1)+,R3
         TSR      (R1)+,R4
         TSR      (R1)+,R5
         RTS      R7
PW1:    WORD     -481.
PW2:    WORD     -32.
POWR:   TSR      R2,-(R1)
```

	TSR	R4,-(R1)
	TSW	R6,-(R1)
	TSR	(R7)+,R2
	TSR	* (R7)+,R4
	TSR	(R7)+,XT
LPP.1:	ADD	R2,X1
	TSI	R2
	BHZS	RTNU
	DEC	R2
	DFC	XT
	TSR	R4,MU
	CLR	R4
	DIV	#10,,R4
	ADD	#60,R4
	BTSR	R4,XT
	TSR	MU,R4
	BRN	LPP.1
RTNU:	TSR	(R1)+,R6
	TSR	(R1)+,R4
	TSR	(R1)+,R2
	RTS	R7
XT:	WORD	0

 ; TERMNT PROGRAM WHICH PUTS VARIABLE CORE SECURITY DATA BACK ON
 ; AND FINISHES THE JOB

TERMNT:	RESET	
	JMS	R4,WRITE
	WORD	200..296,,26000
	SWRITE	1,TMES
WT:	WAITR	1,WT
	STOP	
TMES:	WORD	10.,0,10.
	BYTE	15,12
	HCI	XGOOD BYE%
	END	

 : THE FOLLOWING IS THE DRA'S PROGRAM WHICH ASKS THE DRA FOR HIS *
 : IDENTIFICATION & IF VALID, DOES THE JOBS REQUESTED BY HIM *

```

 1#    75762
 KBD#   0
 TTY#   1
 ID#    177736
 SM#    5
 PWD#   6
 R1#    #41
 R2#    #42
 R3#    #43
 R4#    #44
 R5#    #45
 R6#    #46
 R7#    #47
 TERMNT# 75646
 PRINTS# 70000
 MODIFY# 104100
 INTLZP# 102400
 FLAG1# 26012
 IONS# 26014
 IDNO# 26022
 AA#   21576

```

	RESET	
DMA:	SWHITE	TTY,SMES.1
W1:	WAITR	TTY,W1.
	TSR	AA,SIN.1+2
	SREAD	KBD,SIN.i
WH1:	WAITR	KBD,WR1
S.1:	TSR	#12..R2
	TSR	#IDNO,R3
	TSR	#SIN.1+6,R4
S.2:	RCMP	(R3)+,(R4)+
	RR2C	SError
	DI-C	R2
	RR2C	S.2
	RCMP	(R4).#1Y
	RR2C	S.3
	SWHITE	TTY,SMES.4
W2:	WAITR	TTY,W2
	TSR	AA,IONS+2
	SREAD	KBD,IONS
WR2:	WAITR	KBD,WR2
	CIR	IONS+2
	CIR	SIN.1+2
S.3:	HTST	FLAG1
	RR2S	S.4
	RTSR	FLAG1,SMES.3+41.
	ADD	#30000,SMES.3+40.
	SWRITE	TTY,SMES.3
W3:	WAITR	TTY,W3

	BCMP	FLAG1.#9.
	BR7S	SENROR
	MCLR	FLAG1
S.4:	SWRITE	TTY,SMES.5
w4:	WAITR	TTY,W4
	BRA	S.4.2
S.4.1:	SWRITE	TTY,SMES.6
b4.1:	WAITR	TTY,W4.1
S.4.2:	SREAD	KBD,SIN.1
WK3:	WAITR	KBD,WR3
	TSR	FUNCTN,H2
	TSR	#CUDA_1,R3
L.1:	TSR	#SIN.1+6,R4
	TSR	(R3)+,R6
	TSR	LNGTJ,R5
L.2:	BCMP	(R6)+,(R4)+
	BR7C	NXT
	DEC	R5
	BR7C	L.2
	LSI	R2
	JMP	BRW-2(R2)
NXT:	DEC	R2
	BR7S	S.4
	BRN	L.1
SENROR:	HINC	FLAG1
	BCMP	FLAG1,#9.
	BR1E	S.5
	BTSR	#9.,FLAG1
S.5:	SWRITE	TTY,SMES.2
b5:	WAITR	TTY,WS
	JMP	TERMNT
SIN.1:	WORD	16.,0,16.
	"	*16.
FUNCTN:	WORD	7
LNGTJ:	WORD	6
SMES.1:	WORD	44.,0,44.
	BYTE	15.12
	BCI	% PLEASE GIVE YOUR IDENTIFICATION%
	BCI	% NUMBER %
SMES.2:	WORD	30.,0,30.
	BYTE	15.12
	BCI	% SORRY,UNABLE TO SERVE YOU %
SMES.3:	WORD	42.,0,42.
	BYTE	15.12
	BCI	% ILLEGAL REFERENCE MADE %
	BCI	%EARLIER, TIMES%
SMES.4:	WORD	34.,0,34.
	BYTE	15.12
	BCI	% GIVE NEW IDENTIFICATION%
	BCI	% NUMBER %
SMES.5:	WORD	18.,0,18.
	BYTE	15.12
	BCI	% SPECIFY THE JOB%
SME.6:	WORD	30.,0,30.
	BYTE	15.12

	HCI	% JOB OVER, SPECIFY NEXT JOB %
CODA.1:	WORD	CD1,CD2,CD3,CD4,CD5,CD6,CD7
CD1:	BCI	XINTLZ%
CD2:	BCI	XMODIFY%
CD3:	BCI	%REMOVE%
CD4:	BCI	%TERMMNT%
CD5:	BCI	XPRINTS%
CD6:	BCI	%CCCCCCC%
CD7:	BCI	XRELEASED%
SR:	WORD	RELBLD
	WORD	CCCCCC
	WORD	PRINTS
	WORD	TERMMNT
	WORD	REMOVE
	WORD	MODIFY
	WORD	INTLZE
	HEHLBD:	60040
	CCCCCC:	56630
	CCCCCC:	56630
	REMOVE:	66400

THE FOLLOWING ROUTINE IS USED TO ASSIGN A SECTOR 0
SPACE TO CALLING PROGRAM. CALL IS JMS R7,ALOCTS &
CYL. NO. & SECTOR & SURFACE NO. TO THE NEXT TWO HO

SECURITY *
T RETURNS *
S *

	77106
	ACYL# 26036
	BITMP# 26050
ALOCTS:	TSR R6,-(R1)
	TSR R5,-(R1)
	TSR R4,-(R1)
	TSR R3,-(R1)
	TSR R2,-(R1)
	TSR R1,SSAVE1
	RESET
	JMS R2,RNDY1
R:	WORD 0
	TSR R,MQ
	CLR R
	DIV #62..R
	ASR R
	LSL R
S.11:	TSR R,R2
	TSR #62..R3
	CLR R6
S.12:	TSR #16..R4
	TSR #1,R5
S.13:	TSR R5,BITMP(R2)
	BRZS GOTONE
	LSL R5
	INC R6
	DEC R4
	BRZC S.13
	ADD #2,R2
	CMP R2,#62.
	BRZC S.14

```

S.14:    CLR      R2
        SRS     #2,R3
        BZG      S.12
        SWRITE   TTY,SMES.6
        HLT      TTY,W6
        TSR      EMPTYS,(R7)-
        ADD      #2,R7
        ADD      R,R
        ORN      E5.2
        GUTUM:   STP      H5,HIMP(R2)
        ADD      R,R
        ADD      R,R
        ADD      R,R
        ADD      H6,R
        CMP      H.496.
        BRT      S.15
        SUB     #496.,R
        TSR      H,MQ
        CLR      R
        DIV     #10e.,R
        LSL      MQ
        TSR      MG,R2
        TSR      ACYL(R2),(R7)-
        TSR      R,MQ
        CLR      R
        DIV     #10.,R
        MPY     #32.,R5
        ADD      MQ,R
        INC      R
        ISR      R,(R7)-
E9.2:    TSR      SSAVE1,R1
        TSR      (R1)+,R2
        TSR      (R1)+,R3
        TSR      (R1)+,R4
        TSR      (R1)+,R5
        TSR      (R1)+,R6
        RTS      R7
        WORD    0
        WORD    -1
        WORD    42.,0,42.
        BYTE   .15,12
        RCI      % NO SECTOR AVAILABLE IN %
        HCI      %SECURITY SPACE %

***** FOLLOWING ROUTINE RELEASES A SECTOR OF DISK SECURITY SPACE WITH *
* CALL JMS R7,RELSEC & NEXT TWO WORDS CONTAINING THE ADDRESS OF *
* THE SECTOR TO BE RE EASED .
***** */

RELSEC:  TSR      R6,-(R1)
        TSR      R5,-(R1)
        TSR      R4,-(R1)
        TSR      R3,-(R1)
        TSR      R2,-(R1)
        TSR      R1,SSAVE2
        RESET
        TSR      #0,R2

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S.21:	TSR	(R7)+, CYLSEC
	TSR	(R7)+, SECTRS
	TSR	40, R3
	CMP	CYLSEC, ACYL(R3)
	BRZS	S.22
	ADD	#2, R3
	INC	R2
	CMP	R2, #5
	BRZS	S.24
	BRN	S.21
	TSR	SECTRS, MU
	TSR	#0, R3
	DIV	#32, , R3
	DEC	R3
	MPY	#10, , R4
	ADD	M0, R3
	TSR	R2, MU
	MPY	#100, , R2
	ADD	R3, MU
	CMP	M0, #495,
	BRGT	S.24
	CLR	R3
	DIV	#16, , R3
	INC	R3
	TSR	M0, R4
	ADD	R4, R4
	TSR	#1, SSSS1
S.23:	DFC	R3
	NOP	
	BRZS	S.23.1
	LSL	SSSSS1
	BRN	S.23
S.23.1:	CLD	SSSSS1, BITMP(R4)
	NOP	
	BRN	ES.4
SSSSS1:	WORD	0
S.24:	SWRITE	TTY, SMES.7
W7:	WAITH	TTY, W7
ES.3:	TSR	SSAVE2, R1
	TSR	(R1)+, R2
	TSR	(R1)+, R3
	TSR	(R1)+, R4
	TSR	(R1)+, R5
	TSR	(R1)+, R6
	RTS	H7
SMES.7:	WORD	76, , 0, 76,
	BYTE	15, 12
	BCI	X SECTOR TO BE RELEASED OUT OF X
	BCI	X RANGE, NO SECTOR RELEASED X
	BCI	X, PROGRAM CONTINUES X
SSAVE2:	WORD	0
CYLSEC:	WORD	0
SECTRS:	WORD	0

THIS IS A RANDOM NO. GENERATOR FOR THE ENTIRE SYSTEM

```
RNDY1: TSR      N1,N3
      ADD      N2,N3
      TSR      N2,N1
      TSR      N3,N2
      TSR      N3,(B2)-
      RTS      H2
N1: WORD    0
N2: WORD    1
N3: WORD    0
END
```


	TSR	SSAVEG,R1
	TSR	(R1)+,R4
	RTS	R7
S.38:	SWRITE	TTY,XB
CW:	WAITR	TTY,CW
	SWRITE	TTY,XC
DW:	WAITR	TTY,DW
	SHFAD	KBD,III
WIII:	WAITR	KBD,WIII
	BRN	S.36
III:	WORD	2,0,2,0
SSAVEG:	WORD	0
SAVE1:	WORD	0
XTSUB:	WORD	0
FLAG:	BYTE	0
CH:	BYTE	0
INPUT:	BYTE	1
PRNTR:	BYTE	10
CARDIN:	BYTE	7
PAPIN:	BYTE	5
KYBRD:	BYTE	1
PAPOUT:	BYTE	6
XB:	WORD	2,0,2,3407
XC:	WORD	22,0,22
	BYTE	15,12
	BCT	%INPUT NOT READY%
IN2:	WORD	80,
ST2:	BYTE	0,0
COUNT:	WORD	0
		.+80.
***** THE FOLLOWING ROUTINE READS A 3 DIGIT INPUT & PUTS		
***** FOLLOWING WORD AFTER CONVERTING IT TO IT'S BINARY		
***** JMS R7,GET,3		
***** ***** ***** ***** ***** ***** ***** *****		
GET,3:	TSR	R2,-(R1)
	TSR	R3,-(R1)
	TSR	R4,-(R1)
	TSR	#3,R2
	CLR	MQ
GET,31:	JMS	R7,GET
	BCMP	CH,#40
	BR7S	GET,31
	CLR	R3
	BTSR	CH,R3
	SUB	#60,R3
	MPY	#10..R4
	ADD	R3,MQ
	DEC	R2
	BRZC	GET,31
	TSR	MQ,(R7)+
	TSR	(R1)+,R4
	TSR	(R1)+,R3
	TSR	(R1)+,R2
	RTS	R7

; THE FOLLOWING ROUTINE GENERATES A 10 DIGIT RANDOM
; IT IS A BUFFER WHOSE STARTING ADDRESS FOLLOWS THE
; INPUT LINE.

MBER & PUTS *
LL TO THIS *

```
    TSI      R3,-(R1)
    TSR      R2,-(R1)
    TSR      #10,,R3
    TSI      (R7)+,GEN.1
    JYS      R2,RNDY1
    WORD     0
    CLR      R2
    TSR      RDY,=D
    HTR      #10,,R2
    STSR      R2,@GFR.1
    LDC      GEN.1
    DEC      R3
    BRZC      GEN.S.1
    TSR      (R1)+,R2
    TSR      (R1)+,R3
    RTS      R7
    WORD     0
```

; THIS ROUTINE READS A 10 DIGIT INPUT & GENERATES IT
; BINARY EQUIVALENT.

4 WORD *

```
G&T.10:   TSR      R4,-(R1)
    TSR      R3,-(R1)
    TSR      R2,-(R1)
    TSR      #10,,R2
    JMS      R7,GET
    BCMP     CH,#40
    BR7S     LPGT10
    CLR      R3
    STSR     CH,R3
    SUB      #60,R3
    STSR     R3,(R7)+
    DEC      R2
    BRZC     LPGT10
    TSR      (R1)+,R2
    TSR      (R1)+,R3
    TSR      (R1)+,R4
    RTS      R7
```


; THIS ROUTINE READS THE AUTHORITY VECTOR SPECIFICAT
; MODIFICATION CARD & BUILDS IT'S EQUIVALENT AUTHORI

N. OF ANY *
VECTOR *

```
AUTHSC:  TSR      R4,-(R1)
    TSR      R3,-(R1)
    TSR      R2,-(R1)
    TSR      (R7)+,R4
    TSR      R4,R2
    ADD      #256,,R4
    TSR      #128,,R3
    AU.1:   CLR      (R2)+
```

AU.1:	DEC	R3
	BRZC	AU.1
	TSR	R4,R3
	SUB	#257,.R3
	JMS	R7,GET
	BCMP	CH,#*#
	BRZC	AU.2
AU.2.1:	JMS	R7,AUT.1
AUU1:	WORD	0
	CMP	AU01.#1
	BRET	AU.4
	CMP	AU01.#256.
	BRCT	AU.4
	TSR	AU01.R2
	ADD	R3,R2
AU.3:	JMS	R7,AUT.2
AUU2:	WORD	0
	HTSR	AU02.(R2)+
	HCMP	R2,R4
	BRZS	AUTEND
	BCMP	AU02+1,#*#
	HRZS	AUTEND
	BCMP	AU02+1,#*#
	BRZC	AU.3
	BRN	AU.2.1
AU.4:	SWRITE	TTY,SME.13
W13:	WAITR	TTY,W13
AUTEND:	TSR	(R1)+,R2
	TSR	(R1)+,R3
	TSR	(R1)+,R4
	RTS	R7
SME.13:	WORD	118.,0.118.
	BYTE	15.12
	BCI	% COMMAND ERROR IN AUTHORITY %
	BCI	%VECTOR MODIFICATION,PROGRAM %
	BYTE	15.12
	BCI	% CONTINUES AFTER DELETING PART%
	BCI	% OF COMMAND FOLLOWING ERROR %
AUT.1:	TSR	R2,-(R1)
	CLR	MQ
AUT.12:	JMS	R7,GET
	BCMP	CH,#40
	BRZS	AUT.12
	BCMP	CH,#*,
	BRZS	AUT.13
	MPY	#10.,R2
	BTSR	CH,R2
	SUB	#60,R2
	ADD	R2,MO
	BRN	AUT.12
AUT.13:	TSR	MQ,(R7)+
	TSR	(R1)+,R2
	RTS	R7
AUT.2:	TSR	R2,-(R1)
	CLR	MQ

AUT.21:	JMS	R7,GET
	BCMP	CH,#40
	BR2S	AUT.21
	BCMP	CH,#1.
	BR2S	AUT.22
	BCMP	CH,#1#
	BR2S	AUT.23
	BCMP	CH,#1,
	BR2S	AUT.24
	MPY	#10,R2
	HTSR	CH,R2
	SUB	#60,R2
	ADD	R2,MU
	BRN	AUT.21
	TSR	MU,R2
	HTSR	R2,(R7)+
	HTSR	#1,(R7)+
	BRN	AUT.25
	TSR	MQ,R2
	HTSR	R2,(R7)+
	HTSR	#1#, (R7)+
	BRN	AUT.25
	TSR	MQ,R2
	HTSR	R2,(R7)+
	BCLR	(R7)+
	TSH	(H1)+,R2
	RTS	R7
	END	

INTLZF PROGRAM WHICH INITIALISES THE FIRST SECURITY CODE'S TABLE
TO CONTAIN ONLY DBA'S CODE & MAKES AUGMENTING SECURITY CODES
TABLE EMPTY DEPENDING UPON DBA'S REQUEST.

	102400	
R1=	X1	
R2=	X2	
R3=	X3	
R4=	X4	
R5=	X5	
R6=	X6	
R7=	X7	
KBD=	0	
TIY=	1	
VORBLF=	24400	
DVORBL=	24600	
READ=	022060	
WRITE=	022066	
DRSC.1=	26000	
ISECYL=	26002	
ISECTR=	26004	
ASECYL=	26006	
ASECTR=	26010	
SECSP1=	22170	
SFCSP2=	22570	
EXIT=	76214	
REI SEC=	77550	
ALOCTS=	77106	
INTLZE:		
TSR	R6,-(R1)	
TSR	R5,-(R1)	
TSH	R4,-(R1)	
TSR	R3,-(R1)	
TSR	R2,-(R1)	
TSH	R1,SSAVE3	
RESET		
SWRITE	TTY,SMES.8	
WB:		
WAITR	TTY,WB	
SREAD	KBD,SIN.2	
WR4:		
WAITR	KBD,WR4	
BCMP	SIN.2+6,TTY	
BRZS	FSCINT	
JMP	ASCINT	
FSCINT:		
TSR	ISECYL,CYL.1	
TSR	ISECTR,SCT.1	
JMS	R4,READ	
CYL.1:	WORD	0
SCT.1:	WORD	0
	WORD	SECSP1
	TSR	CYL.1,CYL.2
	TSR	SCT.1,SCT.2
	JMS	R7,RELSEC
CYL.2:	WORD	0

SCT.2:	WORD	0
	JMS	R7,ALOCTS
CYL.3:	WORD	0
SCT.3:	WORD	0
	TSR	CYL.3,1SECYL
	TSR	SCT.3,1SECTR
	TSR	#0,R2
	TSR	#18..R6
S.10:	TST	SECSP1(R2)
	BHBF	S.10
S.17:	ADD	#14..R2
	DEC	R6
	BHZC	S.16
	HRN	S.26
S.18:	TSR	R2,R3
	TST	R2
	BHZS	S.19
	TSR	SECSP1(R2),CYL.4
	ADD	#2,R3
	TSR	SECSP1(R3),SCT.4
	JMS	R7,RELSEC
CYL.4:	WORD	0
SCT.4:	WORD	0
	TSR	#-1,SECSP1(R2)
	BHN	S.17
S.19:	TSR	SECSP1(R2),CYL.5
	ADD	#2,R3
	TSR	SECSP1(R3),SCT.5
	JMS	R4,READ
CYL.5:	WORD	0
SCT.5:	WORD	0
	WORD	SECSP2
	TSR	#4,R4
	TSR	#0,R5
S.20:	INC	R5
	CMP	R5,0BSC.1
	BHZS	S.25
	BTSR	#48..SECSP2(R4)
S.25:	ADD	#5,R4
	CMP	R5,#51.
	BHZC	S.20
	TSR	CYL.5,CYL.6
	TSR	SCT.5,SCT.6
	JMS	R7,RELSEC
CYL.6:	WORD	0
SCT.6:	WORD	0
	JMS	R7,ALOCTS
CYL.7:	WORD	0
SCT.7:	WORD	0
	TSR	CYL.7,CYL.8
	TSR	SCT.7,SCT.8
	JMS	R4,WRITE
CYL.8:	WORD	0
SCT.8:	WORD	0
	WORD	SECSP2

	TSR	CYL.7,SECSP1(R2)
	TSR	SCT.7,SECSP1(R3)
	HRN	S.17
S.26:	TSR	CYL.3,CYL.9
	TSR	SCT.3,SCT.9
	JMS	R4,WHITE
CYL.9:	WORD	0
SCT.9:	WORD	0
	WORD	SECSP1
ASCINT:	SWRITE	TTY,SMBS.9
W9:	WAITR	TTY,W9
	SHEAD	KBD,SIN.2
WRS:	WAITR	KBD,WRS
	BCMP	SIN.2+6,WY
	BRZC	ES.30
	TSR	ASECYL,CYL.10
	TSR	ASECTR,SCT.10
	JMS	R4,HEAD
CYL.10:	WORD	0
SCT.10:	WORD	0
	WORD	SECSP1
	ISR	CYL.10,CYL.11
	TSR	SCT.10,SCT.11
	JMS	R7,RELSEC
CYL.11:	WORD	0
SCT.11:	WORD	0
	JMS	R7,ALOCITS
CYL.12:	WORD	0
SCT.12:	WORD	0
	TSR	CYL.12,ASECYL
	TSR	SCT.12,ASECTR
	TSR	#0,R2
S.27:	TSR	#18,,TEMI.1
	TST	SECSP1(R2)
	BRGF	S.29
S.28:	ADD	#14,,R2
	DEC	TEMI.1
	BRZC	S.27
	HRN	S.33
S.29:	TSR	R2,R3
	TSR	SECSP1(R2),CYL.13
	TSR	#-1,SECSP1(R2)
	ADD	#2,R3
	TSR	SECSP1(R3),SCT.13
	JMS	R4,READ
CYL.13:	WORD	0
SCT.13:	WORD	0
	WORD	SECSP2
	TSR	CYL.13,CYL.14
	TSR	SCT.13,SCT.14
	JMS	R7,RELSEC
CYL.14:	WORD	0
SCT.14:	WORD	0
	TSR	#0,R4
	TSR	#25,,TEMI.2

S.30:	TSR	SECSP2(R4)
	BIGE	S.32
S.31:	ADD	#10,,R4
	DEC	TLM1,2
	HRZC	S.30
	BRN	S.28
S.32:	TSR	R4,R5
	TSR	SECSP2(R4),CYL,15
	AND	#2,R5
	TSR	SECSP2(R5),DVO,1
	ADD	#2,R5
	TSR	SECSP2(R5),DVO,2
	JMS	R7,DVOKDL
DVO,1:	WORD	0
DVO,2:	WORD	0
DVO,3:	WORD	0
	TSR	DVO,3,SCT,15
	JMS	R7,RELSEC
CYL,12:	WORD	0
SCT,15:	WORD	0
	BRN	S.31
S.33:	TSR	CYL,12,CYL,16
	TSR	SCT,12,SCT,16
	JMS	R4,WRITE
CYL,16:	WORD	0
SCT,16:	WORD	0
	WORD	SECSP1
ES,30:	TSR	SSAVE3,R1
	TSR	(R1)+,R2
	TSR	(R1)+,R3
	TSR	(R1)+,R4
	TSR	(R1)+,R5
	TSR	(R1)+,R6
	JMP	EXIT
SSAVE3:	WORD	0
TEMI,1:	WORD	0
TEMI,2:	WORD	0
SMFS,8:	WORD	52,,0,52,
	BYTE	15,12
	BCI	% DO YOU WANT TO INITIALIZE%
	HCI	% FIRST SECURITY CODES %
SMFS,9:	WORD	54,,0,54,
	BYTE	15,12
	BCI	% DO YOU WANT TO INITIALIZE%
	BCI	% AUTHORITY VECTOR CODES %
SIN,2:	WORD	2,,0,2
	/*	+2
	END	

```
*****  
; MODIFY PROGRAM WHICH ACCEPTS DHA'S REQUEST FOR DELETING, ADDING *  
; OR MODIFYING ANY SECURITY CODE OR AUTHORITY VECTOR.  
*****
```

```
.= 104100  
FLAG= 100706  
CARDIN= 100712  
KYBRO= 100714  
PAPIN= 100713  
CH= 100707  
SFCSPI= 22170  
SFCSPI= 22570  
EXIT= 76214  
VORBLE= 24400  
DVORBI= 24600  
READ= 22060  
WRITE= 22066  
ISFCYL= 26002  
ISFCYL= 26004  
ASECYL= 26006  
ASECTR= 26010  
MQ= 177736  
ALOCIS= 77106  
INPUT= 100710  
RELSEC= 77550  
GET.3= 101104  
GET.10= 101262  
GIVSEC= 25000  
WORBLE= 25400  
GFNWOR= 101200  
GFT= 100400  
AUTHSC= 101340
```

```
R1=%1  
R2=%2  
R3=%3  
R4=%4  
R5=%5  
R6=%6  
R7=%7
```

	KBD=	0
	TTY=	1
MODIFY:	RESET	
	SWRITE	TTY,SME'.10
W10:	WAITR	TTY,W10,
	SREAD	KBD,SIN'.3
WR6:	WAITR	KBD,WR6
	BCLR	FLAG
	CMP	SIN'.3+6,#"CA
	BRZC	S'.41
	BTZR	CARDIN,INPUT
	BRN	S'.43
S.41:	CMP	SIN'.3+6,#"KY
	BRZC	S'.42

	BTSR	KYBRD, INPUT
S.42:	BRN	S.43
	CMP	SIN, 3+6, # "PA
	BRZC	MODIFY
	BTSR	PAPIN, INPUT
S.43:	JMS	R7, GET
	BCMP	CH, #40
	BRZS	S.43
S.44:	BCMP	CH, #'D
	BRZC	S.45
	JMP	S.49
S.45:	BCMP	CH, #'E
	BRZC	S.46
	JMP	S.53
S.46:	BCMP	CH, #'F
	BRZC	S.47
	JMP	EXIT
S.47:	SWRITE	TTY, SME.11
W11:	WAITR	TTY, W11
S.48:	JMS	R7, GET
S.AU:	BCMP	CH, #'.
	BRZC	S.48
	BRN	S.43
S.49:	JMS	R7, GET
	BCMP	CH, #40
	BRZS	S.49
	BCMP	CH, #61
	BRZS	S.50
	BCMP	CH, #62
	BRZS	S.51
	BRN	S.47

; THIS SEGMENT DELETES THE FIRST SECURITY CODE AS REQUESTED BY DBA*

S.50:	JMS	R7, GET.3
C.1:	WORD	0
	CMP	C.1, #1
	BRLT	S.47
	CMP	C.1, #918.
	HRGT	S.47
	TSR	ISECYL, CYL, 21
	TSR	ISELECTR, SCT, 21
	JMS	R4, READ
CYL.21:	WORD	0
SCT.21:	WORD	0
	WORD	SECSP1
	SUB	#1, C.1
	TSR	C, 1, MQ
	CLR	C.1
	DIV	#51, .C.1
	MPY	#14, .R2
	TSR	MQ, R2
	TSR	SECSP1(R2), CYL, 22
	TSR	SECSP1+2(R2), SCT, 22
	TST	CYL, 22

	BRGE	NXTMOD
	JMP	S.52
NXTMOD:	JMS	R4,READ
CYL,22:	WORD	0
SCT,22:	WORD	0
	WORD	SECSP2
	TSR	C.1,MQ
	MPY	#5,C.1
	TSR	MQ,R2
	BTSR	#48,,SECSP2+4(R2)
	ISR	CYL,22,CYL,23
	ISR	SCT,22,SCT,23
	JMS	R4,WRITE
CYL,23:	WORD	0
SCT,23:	WORD	0
	WORD	SECSP2
	JMP	S.48

***** THIS SEGMENT DELETES THE AUTHORITY AUGMENTING SECURITY CODE AS *

***** REQUESTED BY THE DBA. *

S,51:	JMS	R7,GET,3
C,2:	WORD	0
	CMP	C,2,#1
	BRLT	S,47
	CMP	C,2,#450
	BRGT	S,47
	ISR	A\$ECYL,CYL,24
	TSR	A\$ECTR,SCT,24
CYL,24:	JMS	R4,READ
SCT,24:	WORD	0
	WORD	0
	WORD	SECSP1
	SUB	#1,C,2
	ISR	C,2,MQ
	CLR	C,2
	DIV	#25,,C,2
	MPY	#14,,R2
	TSB	MQ,R2
	TSB	SECSP1(R2),CYL,25
	TSB	SECSP1+2(R2),SCT,25
	TST	CYL,25
	BRLT	S,52
CYL,25:	JMS	R4,READ
SCT,25:	WORD	0
	WORD	0
	WORD	SECSP2
	TSB	C,2,MQ
	MPY	#10,,C,2
	TSB	MQ,R2
	TST	SECSP2(R2)
	BRLT	S,52
	ISR	SECSP2(R2),CYL,26
	TSB	#-1,SECSP2(R2)
	TSB	SECSP2+2(R2),DVO,4

	TSR	SECSP2+4(R2), DVO.5
DVO.4:	JMS	R7,DVORBL
DVO.5:	WORD	0
DVO.6:	WORD	0
	TSR	DVO.6,SCT.26
CYL.26;	JMS	R7,RELSEC
SCT.26:	WORD	0
	TSR	CYL.25,CYL.27
	TSR	SCT.25,SCT.27
CYL.27;	JMS	R4,WRITE
SCT.27:	WORD	0
	WORD	0
S.52:	JMP	SECSP2
W12:	SWRITE	S.48
	WAITR	TTY,SME.12
	JMP	TTY,W12
S.53:	JMP	S.48
	JMS	R7,GET
	BCMP	CH,#40
	BRZS	S.53
	CLR	R6
	BISB	CH,R6
	BCMP	CH,#61
	BBZS	S.57
	BCMP	CH,#62
	RRZC	S.54
S.54:	JMP	S.60
	BCMP	CH,#63
	BRZC	S.55
S.55:	JMP	S.60
	BCMP	CH,#64
	BRZS	S.56
S.56:	JMP	S.47
	JMP	S.66
<hr/>		
THIS SEGMENT ENTERS A NEW FIRST SECURITY CODE OR M		IFIES AN
EXISTING ONE AS REQUESTED BY THE DBA.		
S.57:	JMS	R7,GET.3
C.3:	WORD	0
	OMP	C.3,#1
	BRGE	S.E.1
S.5.11	JMP	S.47
	OMP	C.3,#918
	BRLE	S.E.2
S.5.21	JMP	S.47
	TSR	ISECYL,CYL.28
	TRP	ISECTR,SCT.28
	JMS	R4,READ
CYL.28:	WORD	0
SCT.28:	WORD	0
	WORD	SECSP1
	SHR	#1,C.3

	TSR	C.3, MO
	CLR	C.3
	DIV	#51..C.3
	MPY	#14..R2
	TSR	MO, R2
	TST	SECSP1(R2)
	BPL T	S.E.5
	TSR	SECSP1(R2), S.E.3
	TSR	SECSP1+2(R2), S.E.4
	JMS	R4, HEAD
S.E.4:	WORD	0
S.E.4:	WORD	0
	WORD	SECSP2
	BRN	S.59
S.E.5:	JMS	R7, ALOCIS
CYL.29:	WORD	0
SCT.29:	WORD	0
	TSR	CYL.29, SECSP1(R2)
	TSR	SCT.29, SECSP1+2(R2)
	TSR	#4, R3
S.58:	TSR	#51..R4
	BTSR	#48..SECSP2(H3)
	ADD	#5, R3
	DEC	R4
	BRZC	S.58
	TSR	#SECSP1+4, WOR.1
	ADD	R2, WOR.1
	JMS	R7, GENWOR
WOR.1:	WORD	0
S.59:	TSR	#SECSP1+4, WOR.2
	ADD	R2, WOR.2
	JMS	R7, GET.10
SE.1:	.*	.+10.
	TSR	#SE.1, WOR.3
	TSR	#SE.1, WOR.4
	JMS	R7, WORBLE
WOR.2:	WORD	0
WOR.3:	WORD	0
	JMS	R7, GIVSEC
WOR.4:	WORD	0
WOR.5:	.*	.+4
	TSR	C.3, MO
	MPY	#5, C.3
	TSR	MO, R3
	BCIR	SECSP2+4(R3)
	BTSR	WOR.5, SECSP2(H3)
	BTSR	WOR.5+1, SECSP2+1(R3)
	BTSR	WOR.5+2, SECSP2+2(R3)
	BTSR	WOR.5+3, SECSP2+3(R3)
	TSR	SECSP1(R2), CYL.30
	TSR	SECSP1+2(R2), SCT.30
	JMS	R4, WRITE
CYL.30:	WORD	0
SCT.30:	WORD	0

```

WORD      SECSP2
TSR       ISECYL.CYL.31
TSR       ISECTR.SCT.31
JMS       R4,WRITE
CYL.31: WORD      0
SCT.31: WORD      0
WORD      SECSP1
JMP       S.48
***** THIS SEGMENT ENTERS A NEW AUTHORITY AUGMENTING SECURITY CODE OR *
* MODIFIES AN EXISTING ONE WITH IT'S ASSOCIATED VECTOR KEPT INTACT*
*****
S.60:   JMS       R7,GET.3
C.4:    WORD      0
        CMP       C.4,#1
        BRGE    S.E.6
        JMP       S.47
S.E.6:  CMP       C.4,#450.
        BRIF    S.E.7
        JMP       S.47
S.E.7:  TSR       ASECYL.CYL.32
        TSR       ASECTR.SCT.32
        JMS       R4,READ
CYL.32: WORD      0
SCT.32: WORD      0
        WORD      SECSP1
        SUB      #1,C.4
        TSR       C.4,MQ
        CLR       C.4
        DIV       #25.,C.4
        MPY       #14.,R2
        TSR       MQ,R2
        TST       SECSP1(R2)
        BRLT    S.E.10
        TSR       SECSP1(R2),S.E.8
        TSR       SECSP1+2(R2),S.E.9
        JMS       R4,READ
S.E.8:  WORD      0
S.E.9:  WORD      0
        WORD      SECSP2
        BRN      S.62
S.E.10: JMS       R7,ALOCTS
CYL.33: WORD      0
SCT.33: WORD      0
        TSR       CYL.33,SECSP1(R2)
        TSR       SCT.33,SECSP1+2(R2)
        CLR       R3
        TSR       #25.,R4
        TSR       #-1,SECSP2(R3)
        ADD      #10.,R3
        DFC
        BRZC    S.61
        TSR       #SECSP1+4,WOR.6
        ADD      R2,WOR.6
        JMS       R7,GENWOR

```

WOR.6:	WORD	0
J.62:	TSR	#SECSP1+4,WOR.7
	ADD	R2,WOR.7
	JMS	R7,GET.10
G.62:	L=	.+10.
	TSR	#SE.2,WOR.8
	TSR	#SE.2,WOR.9
	JMS	R7,WORBLE
WOR.7:	WORD	0
WOR.8:	WORD	0
	JMS	R7,GIVSFC
WOR.9:	WORD	0
WOR.10:	L=	.+4
	TSR	C.4,MQ
	MPY	#10..C.4
	TSR	MQ,R3
	RTSR	WOR.10,SECSP2+6(R3)
	RTSR	WOR.10+1,SECSP2+7(R3)
	RTSR	WOR.10+2,SECSP2+8,(R3)
	RTSR	WOR.10+3,SECSP2+9,(R3)
	TSR	SECSP1(R2),CYL.36
	TSR	SHCSP1+2(R2),SCT.36
	TSR	ASECYL,CYL.34
	TSR	ASECTR,SCT.34
	JMS	R4,WHITE
CYL.34:	WORD	0
SCT.34:	WORD	0
	WORD	SECSP1
	TST	SECSP2(R3)
	BRLT	S.E.13
	TSR	CYL.36,S.E.11
	TSR	SCT.36,S.E.12
	JMS	R4,WHITE
S.E.11:	WORD	0
S.E.12:	WORD	0
	WORD	SECSP2
	BRN	S.64
S.E.13:	JMS	R7,ALOCTS
CYL.35:	WORD	0
SCT.35:	WORD	0
	TSR	CYL.35,SECSP2(R3)
	TSR	SCT.35,VO.1
	JMS	R7,WORBLE
VO.1:	WORD	0
VO.2:	WORD	0
VO.3:	WORD	0
	TSR	VO.2,SECSP2+2(R3)
	TSR	VO.3,SECSP2+4(R3)
????	TSR	VO.3,SECSP2+4(R3)
	JMS	R4,WHITE
CYL.36:	WORD	0
SCT.36:	WORD	0
	WORD	SECSP2
	CLR	R2
S.63:	CLR	SECSP1(R2)

	INC	R2
	INC	R2
	CMP	R2,#256.
	BRZC	S.63
	BCMP	R6,#62
	BH7C	S.65
	TSR	CYL.35,CYL.37
	TSR	SCT.35,SCT.37
	JMS	R4,WRITE
CYL.37:	WORD	0
SCT.37:	WORD	0
	WORD	SECSP1
	JMP	S.48
S.64:	BCMP	R6,#62
	BRZC	S.E.14
	JMP	S.43
S.E.14:	TSR	SECSP2(R3),CYL.35
	TSR	SECSP2+2(R3),DVO.7
	TSR	SECSP2+4(R3),DVO.8
	JMS	R7,DVORBL
DVO.7:	WORD	0
DVO.8:	WORD	0
DVO.9:	WORD	0
	TSR	DVO.9,SCT.35
S.65:	JMS	R7,AUTHSC
	WORD	SECSP1
	TSR	CYL.35,CYL.38
	TSR	SCT.35,SCT.38
	JMS	R4,WRITE
CYL.38:	WORD	0
SCT.38:	WORD	0
	WORD	SECSP1
	JMP	S.AU

***** THIS SEGMENT CHANGES THE AUTHORITY VECTOR ASSOCIATED WITH A GIVIN SECURITY CODE WHILE LEAVING THE CODE INTACT *****

S.66:	JMS	R7,GET.3
C.5:	WORD	0
	CMP	C.5,#1
	BRGE	S.E.15
	JMP	S.47
S.E.15:	CMP	C.5,#450.
	BRLE	S.E.16
	JMP	S.47
S.E.16:	TSR	ASECYL,CYL.39
	TSR	ASECTR,SCT.39
	JMS	R4,READ
CYL.39:	WORD	0
SCT.39:	WORD	0
	WORD	SECSP1
	SUB	#1,C.5
	TSR	C.5,M0
	CLR	C.5
	DIV	#25..C.5

	MPY	#14..R2
	TSR	MQ,R2
	TST	SECSP1(R2)
	HRDF	S.67
	JMP	S.47
S.67:	TSR	SECSP1(R2), CYL.40
	TSR	SECSP1+2(R2), SCT.40
	JMS	R4,READ
CYL.40:	WORD	0
SCT.40:	WORD	0
	WORD	SECSP2
	TSR	C.5,MD.
	MPY	#10..C.5
	TSR	MQ,R3
	TST	SECSP2(R3)
	HRDF	S.E.14
	JMP	S.47
SME.10:	WORD	66..0,66.
	BYTE	15,12
	BCI	% SPECIFY THE DEVICE CONTAINING%
SME.11:	WORD	% SECURITY CODES MODIFICATION CARDS%
	BCI	44..0,44.
	BYTE	15,12
	BCI	% ILLEGAL MODIFICATION INSTRUCTION %
SME.12:	WORD	% OMITTED %.
	BYTE	60..0,60.
	BCI	15,12
SIN.3:	WORD	% DELETING AN EMPTY CODE.%
	BCI	% NO ACTION TAKEN, PROGRAM CONTINUES.%
	WORD	8..0,8.
	END	.+8.

```

; REMOVE PROGRAM WHICH REMOVES A RELATION FROM THE DATA BASE AS *
; REQUESTED BY THE DBA                                         *

;      .E      66400
;      R1=      #1
;      R2=      #2
;      R3=      #3
;      R4=      #4
;      R5=      #5
;      R6=      #6
;      R7=      #7
;      TTY=      1
;      KBD=      0
;      MDE=      177736
;      FLAG=      100706
;      INPUT=      100710
;      NIND=      30626
;      RELID=      30366
;      KEYRP=      31016
;      FMKRP=      27600
;      FILMOT=      46156
;      FLDIIDE=      35516
;      FILDNME=      31516
;      FOMATE=      36502
;      FRIST=      37466
;      FDLISOC=      60006
;      PMINX=      30506
;      PRINXC=      50472
;      PRINDX=      40060
;      BITMAP=      27740

;      BITMAP=      27740
;      GET.3=      101104
;      PUT=      66000
;      EXIT=      76214
;      REMOVE:      CLR      FLAG
;                  BLSR      #1,INPUT
;                  SWRITE    TTY,RME.3
;      WRN2:      WAITR    TTY,WHM2
;                  JMS      R7,GET.3
;      R.1:      WORD      0
;                  CMP      R.1,#39
;                  BRGT     R.021
;                  TST      R.1
;                  BRLE     R.021
;                  TSR      R.1,R2
;                  CLR      R3
;                  BLSR     NIND(R2),R3
;                  LSL      R2
;                  TST      RELID(R2)
;                  BRZS     R.021
;                  CLR      RELID(R2)
;                  CLR      KEYRP(R2)
;                  TSR      FMKRP(R2),R.01

```

	JMS	R5,FLNUT
	WORD	R.01
	WORD	R.02
	ADD	#3,R.02
	TSR	R.02,MQ
	MPS	R3,R.01
	TSR	MQ,R.01
	TSR	#372,R6
	CLR	R3
H.2:	BGMP	R.1,FLD1D(R3)
	BRZC	R.3
	ADD	#2,R3
	DEC	R6
	BRZC	R.2
H.021:	RESET	
	SWRITE	TTY,RME.1
RW1:	WAITR	TTY,RW1
	JMP	EXIT
H.3:	TSR	R3,R5
	TSR	#1,R4
	ADD	#2,R3
H.4:	BGMP	R.1,FLD1D(R3)
	BRZC	R.5
	INC	R4
	ADD	#2,R3
	HRN	R.4
H.5:	TSR	FLD1D(R3),FLD1D(R5)
	TSR	FOMAT(R3),FOMAT(R5)
	LSL	R3
	LSL	R3
	LSL	R5
	LSI	R5
	TSR	FLDNM(R3),FLDNM(R5)
	TSR	FLDNM+2(R3),FLDNM+2(R5)
	TSR	FLDNM+4(R3),FLDNM+4(R5)
	TSR	FLDNM+6(R3),FLDNM+6(R5)
	ASR	R3
	ASR	R3
	ASR	R3
	ASR	R5
	BTSH	FDIST(R3),FDIST(R5)
	INC	R3
	INC	R5
	CMP	R3,FDLISC
	BRGE	R.6
	LSL	R3
	LSL	R5
	BRN	R.5
H.6:	CMP	R5,FDLISC
	BRGE	R.61
	LSL	R5
	CLR	FLD1D(R5)
	ASR	R5

	INC	R5
	BRN	R.6
H.61:	SUB	R4,FDLISC
	TSR	PMINX(R2),R3
	ASR	R2
	ADD	R.02,R3
	SUB	#3,R3
	CIR	R4
	HTSR	NIND(R2),R4
	LSL	R2
H.071:	HTSH	PRINDEX(R3),R.072
	JMS	R7,HELCYL
H.072:	WORD	0
	ADD	R.02,R3
	DEC	R4
	BIGT	R.071
	TSR	PMINX(R2),R3
	TSR	R3,R6
	SUB	R.01,PRINXC
	TSR	R.01,R5
	ADD	R3,R5
H.7:	HTSR	(R5)+,(R3)+
	CMP	R3,PRINXC
	BRGE	R.8
	BRN	R.7
H.8:	TSR	#2,R2
H.9:	TST	RELID(R2)
	BRZS	R.10
	CMP	R6,PMINX(R2)
	BRGT	R.10
	SUB	R.01,PMINX(R2)
H.10:	ADD	#2,R2
	CMP	R2,#80.
	BRLT	R.9
	JMP	PUT
H.01:	WORD	0
H.02:	WORD	0
HME.1:	WORD	52.,0,52.
	BYTE	15,12
	BCI	% DELETING A NON EXISTING%
	BCI	% RELATION,COMMAND IGNORED %
HME.3:	WORD	66.,0,66.
	BYTE	15,12
	BCI	% SPECIFY THE RELATION NO. OF%
	BCI	% RELATION TO BE REMOVED %
	BCI	XIN 3 DIGITS %

THIS ROUTINE IS USED TO RELEASE THE CYLINDERS VACATED BY THE		
RELATION WHICH IS REMOVED FROM THE DATA BASE		

HELCYL:	TSR	R2,-(R1)
	TSR	R3,-(R1)
	TSR	R4,-(R1)
	TSR	(R7)+,R2
	TST	R2

R.15:	BRI T	R.20
	TSR	#24,,R3
	CMP	R2,#16.
	BRI T	R.17
	SUB	#16,,R2
	SUB	#2,R3
	TST	R3
	BRI T	R.20
	BIN	R.15
R.16:	WORD	0
R.17:	TSR	#1,R.16
R.18:	TST	R2
	BRZS	R.19
	DEC	R2
	LSI	R.16
	HBN	R.18
R.19:	CLR	R.16,BITMAP(R3)
	TSR	(R1)+,R4
	TSR	(R1)+,R3
	TSR	(R1)+,R2
	RTS	R7
R.20:	SWRITE	TTY,RME.2
RW2:	WAITR	TTY,RW2
	STOP	
	RME.2:	WORD 40,,0,40.
	BYTE	15,12
	BCI	% CYLINDER TO BE RELEASED OUT%
	BCI	% OF RANGE %
	FND	

```

; PUT ROUTINE WHICH PUTS RELATION DIRECTORY, FIELD LIST & PRIMARY *
; INDEX TABLE & OTHER VARIABLE DATA FROM CORE INTO DISK.          *

        .EQU    66000
        R1 = %1
        R2 = %2
        R3 = %3
        R4 = %4
        R5 = %5
        R6 = %6
        R7 = %7
        NDS = 26146
        ND1 = 26150
        WRITE = 22066
        EXIT = 76214
        TSR = ND,R2
        JMS = R4,WHITE
        WORD = 201..101..50400
        JMS = R4,WHITE
        WORD = 201..138..60000
        TSR = ND1,R3
        TSR = #25..R5
        TSH = (R3)+,R4
        TSR = (R3)+,STR10
        TSH = (R3)+,CYL10
        P2:   TSR = (RJ)+,SCT10
        JMS = R4,WHITE
CYL10:  WORD = 0
SCT10:  WORD = 0
STR10:  WORD = 0
        ADD = #400,STR10
        DEC = R5
        BRZC = P3
        JMP = EXIT
P3:    DEC = R4
        BNCT = P2
        DEC = R2
        BNCT = P1
        JMP = EXIT
END

```